



**Croll- π , a reversible calculus
with compensations**

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Roadmap

- The approach
- Croll- π
- Applications
- Conclusions



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Uncontrolled perfect reversibility

- RCCS, CCSk, $\rho\pi$, reversible structures all feature uncontrolled perfect reversibility
 - Uncontrolled since the process can go nondeterministically backward or forward
 - Perfect since a backward step perfectly undoes a forward step
- Some approaches feature some form of control of reversibility
 - Irreversible actions in RCCS
 - Roll- π roll operator
 - Controller process in CCSk
- Reversibility is still perfect in all these approaches

Drawbacks of perfect reversibility

- With perfect reversibility one goes back to a past state
- From this state the same computation can restart
 - Probably leading to the same rollback
 - Again and again
- The programmer has no tool to avoid this
- We want to provide such a tool
 - Using a compensation mechanism
- We build on top of roll- π

Do you remember roll- π ?



- HO π extended with an explicit roll operator
- The roll has a parameter γ that refers to a past trigger
- The roll undoes all the actions depending on the communication done by the trigger
- Idea: in case of error I undo the computation that caused the error

Compensations

- From database theory and business processes
 - Piece of code to recover from an error
- For us, a piece of code allowing to move
 - From the past state reached by the rollback
 - To a slightly different state
 - Keeping into account the past try
 - Trying to avoid to repeat the same error
- We call $\text{croll-}\pi$ the language integrating compensations into $\text{roll-}\pi$

Messages with compensations

- Instead of roll- π messages $a\langle P \rangle$ we use messages with compensations
 - $a\langle P \rangle \% 0$: try $a\langle P \rangle$ then stop trying
 - $a\langle P \rangle \% b\langle Q \rangle \% 0$: try $a\langle P \rangle$ then $b\langle Q \rangle$ then stop trying
- If the message with compensation is the target of a rollback, it is replaced by its compensation
- Minimal change to roll- π syntax and semantics
- The expressive power increases considerably

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Croll- π syntax

$$P, Q ::= \mathbf{0} \mid X \mid \nu a. P \mid (P \mid Q) \mid a(X) \triangleright_{\gamma} P \mid a\langle P \rangle \div C \mid \text{roll } k \mid \text{roll } \gamma$$
$$M, N ::= \mathbf{0} \mid \nu u. M \mid (M \mid N) \mid k : P \mid [\mu; k] \mid k \prec (k_1, k_2)$$
$$C ::= a\langle P \rangle \div \mathbf{0} \mid \mathbf{0}$$
$$\mu ::= (k_1 : a\langle P \rangle \div C) \mid (k_2 : a(X) \triangleright_{\gamma} Q)$$

- HO π processes + **roll** + γ + **compensations**
- Configurations featuring memories and causal dependencies

Croll- π semantics

$$\frac{\mu = (k_1 : a \langle P \rangle \div C) \mid (k_2 : a(X) \triangleright_{\gamma} Q_2)}{(k_1 : a \langle P \rangle \div C) \mid (k_2 : a(X) \triangleright_{\gamma} Q_2) \longrightarrow \nu k. (k : Q_2 \{^{P,k} / X, \gamma \}) \mid [\mu; k]}$$

$$\frac{k \prec : N \text{ complete}(N \mid [\mu; k] \mid (k_r : \text{roll } k)) \quad \mu' = \text{xtr}(\mu)}{N \mid [\mu; k] \mid (k_r : \text{roll } k) \longrightarrow \mu' \mid N \not\downarrow k}$$

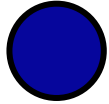
$$\frac{}{k : P \mid Q \longrightarrow \nu k_1, k_2. k \prec (k_1, k_2) \mid k_1 : P \mid k_2 : Q}$$

$$\frac{}{k : \nu a. P \longrightarrow \nu a. k : P}$$

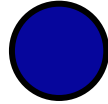
- Instead of restoring the initial configuration we replace the message with its compensation
- $\text{xtr}(a \langle P \rangle \% C) = C$

Croll- π example

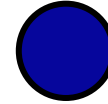
$$k_1 : a\langle \mathbf{0} \rangle \% b\langle \mathbf{0} \rangle$$



$$k_2 : a(X) \triangleright_{\gamma} b\langle \text{roll } \gamma \rangle$$



$$k_3 : b(X) \triangleright c\langle \mathbf{0} \rangle \mid X$$

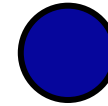


Croll- π example

$$k_1 : a\langle \mathbf{0} \rangle \% b\langle \mathbf{0} \rangle$$

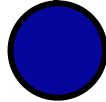
$$k_2 : a(X) \triangleright_{\gamma} b\langle \text{roll } \gamma \rangle$$

$$k_3 : b(X) \triangleright c\langle \mathbf{0} \rangle \mid X$$



$$[k_1 : M \mid k_2 : N; k]$$

$$k : b\langle \text{roll } k \rangle$$

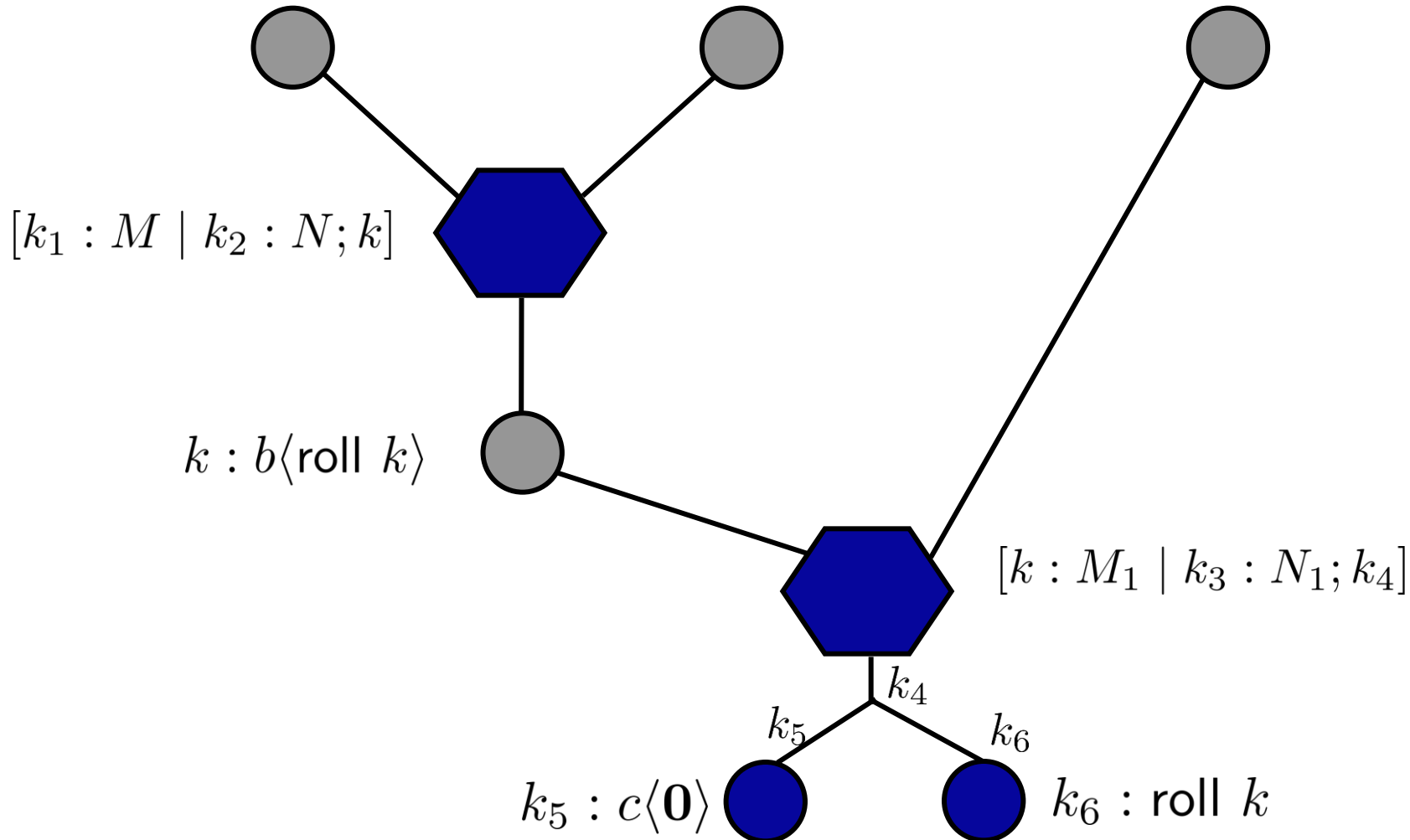


Croll- π example

$$k_1 : a\langle \mathbf{0} \rangle \% b\langle \mathbf{0} \rangle$$

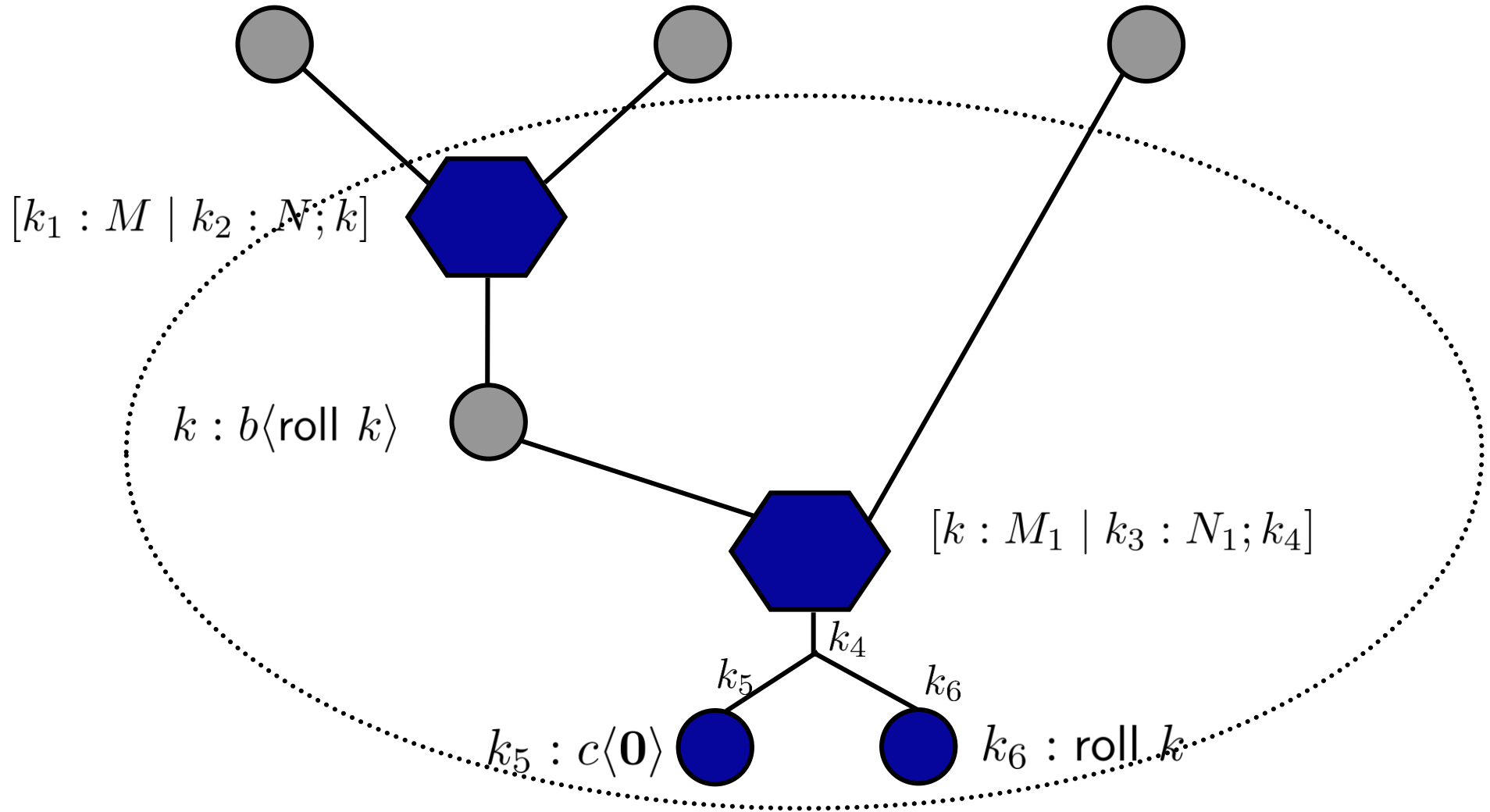
$$k_2 : a(X) \triangleright_{\gamma} b\langle \text{roll } \gamma \rangle$$

$$k_3 : b(X) \triangleright c\langle \mathbf{0} \rangle \mid X$$



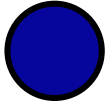
Croll- π example

$k_1 : a\langle \mathbf{0} \rangle \% b\langle \mathbf{0} \rangle$
 $k_2 : a(X) \triangleright_{\gamma} b\langle \text{roll } \gamma \rangle$
 $k_3 : b(X) \triangleright c\langle \mathbf{0} \rangle \mid X$

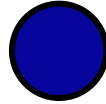


Croll- π example

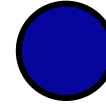
$k_1 : b\langle \mathbf{0} \rangle$



$k_2 : a(X) \triangleright_{\gamma} b\langle \text{roll } \gamma \rangle$



$k_3 : b(X) \triangleright c\langle \mathbf{0} \rangle \mid X$



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And now?

- We have to test the expressive power of $\text{croll-}\pi$
- Can we programme interesting applications exploiting rollback and compensations?
- Compensations allow to specify what to do after a rollback
- $\text{Croll-}\pi$ is strictly more expressive than $\text{roll-}\pi$

Messages with compensations are robust

- We can encode different idioms:
 - General compensations: not only messages
 - Finite retry: try n times
 - Endless retry: try forever (same as roll- π)
 - Triggers with compensations: we can attach compensations to triggers instead of to messages

Triggers with compensations

$$a(X) \triangleright_{\gamma} Q \div C$$

- Compensation replaces trigger instead of message

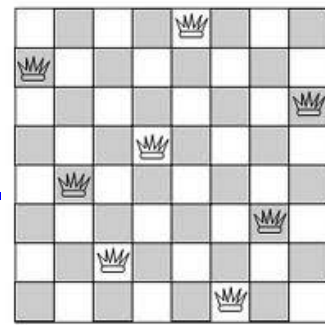
$$\llbracket a(X) \triangleright_{\gamma} Q \div C \rrbracket_{tc} = \nu c, d. \bar{c} \div \bar{d} \div \mathbf{0} \mid (c \triangleright_{\gamma} a(X) \triangleright \llbracket Q \rrbracket_{tc}) \mid (d \triangleright \llbracket C \rrbracket_{tc})$$

- Messages with compensations and triggers with compensations can coexist
- Makes the framework more symmetric

What can we model?

- Interesting problems
 - State space exploration with backtracking: 8 queens problem
 - Error handling scenario: Automotive case study from Sensoria project
- Can we recover/improve existing techniques?
 - Interacting transactions from Hennessy et al. [CONCUR 2010]
 - Software transactional memories
 - » Different approaches
 - » We started from the process calculus account by Acciai et al. [ESOP'07]

Eight queens problem



- Queens are interacting processes
- Queen i tries the 8 positions on column i then fails by notifying the previous queen
- If it succeeds it notifies the next queen, and waits for failures from her

$$Q_i \triangleq c_1(\mathbf{x}_1) \triangleright \dots c_{i-1}(\mathbf{x}_{i-1}) \triangleright p_i \langle i, 1 \rangle \div \dots \div p_i \langle i, 8 \rangle \div f_i \langle 0 \rangle \div 0$$

$$\begin{aligned} & | p_i(\mathbf{x}) \triangleright_{\gamma_i} \text{ if } err(\mathbf{x}_1, \mathbf{x}) \text{ then roll } \gamma_i \text{ else } \dots \text{ if } err(\mathbf{x}_{i-1}, \mathbf{x}) \text{ then roll } \gamma_i \\ & \text{ else } !c_i \langle \mathbf{x} \rangle \div 0 | f_{i+1}(y) \triangleright \text{ roll } \gamma_i \end{aligned}$$

$$err((x_1, x_2), (y_1, y_2)) \triangleq (x_1 = y_1 \vee x_2 = y_2 \vee |x_1 - y_1| = |x_2 - y_2|)$$

- We also have a more concurrent solution

Interacting transactions

- Allow a transaction to interact with the environment
- In case of rollback the effects on the environment are undone

$$\overline{\bar{a} | a.P \longrightarrow P} \qquad \frac{k \notin \text{fn}(R)}{\overline{\llbracket P \triangleright_k Q \rrbracket | R \longrightarrow \llbracket P | R \triangleright_k Q | R \rrbracket}}$$

$$\overline{\llbracket P | \text{co } k \triangleright_k Q \rrbracket \longrightarrow P} \qquad \overline{\llbracket P \triangleright_k Q \rrbracket \longrightarrow Q}$$

- Parts of the environment moved inside the transaction to interact with it
 - Also saved in the compensation

Modeling interacting transactions

- A transaction is a computation originated by a trigger labeled by γ
- Roll γ used as an abort
- Commit by disabling roll γ

$$\llbracket P \triangleright_l Q \rrbracket_t = \nu a, c. \bar{a} \div \bar{c} \div \mathbf{0} \mid a \triangleright_\gamma (\nu l. \llbracket P \rrbracket_t \mid l \langle \text{roll } \gamma \rangle \mid l(X) \triangleright X) \mid c \triangleright \llbracket Q \rrbracket_t$$

- More precise than Hennessy
- Only parts of the environment that interacted with the transaction are actually rolled back

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Summary

- A framework supporting controlled rollback and compensations
- Some interesting applications
 - 8 queens problem
 - Interacting transactions
 - Software transactional memories

Future work



- Is this enough for programming recoverable systems?
 - Some good hints, but more work is needed
- A known limitation
 - We have no information on what caused the roll
 - The compensation does not depend on the reason of the failure
 - A mechanism to keep this information would be helpful
- Studying the behavioral theory
 - To reason about the encodings
 - To have an axiomatic theory
 - History information makes things more complex

Finally

Thanks!

Questions?