

# Network Science: Aggregation

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## Processes on networks

- Networks serve as the underlying transport mechanism for processes that are being carried out *on top* of them
- Interested in studying the effects of the underlying network on the dynamics of the higher-level process
- Processes already studied
  - Gossiping
  - Heartbeat synchronization
  - Formation creation
- Today we will study
  - Aggregation

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## Aggregation (collective computation)

- Initially, each node has a (numeric) local state
- Want to compute in a *decentralized manner* a (global) aggregate function over the initial values
- In the end, the aggregate value must be known (locally) at each node
- Examples of aggregate functions:
  - Average
  - Min-max
  - Geometric mean
  - Variance
  - Network size

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## Aggregation Gossip framework instantiation

- Style of interaction: push-pull
- Local state  $s$ : Current estimate of global aggregate
- Method `selectPeer()`: Single random neighbor
- Method `update()`: Numerical function defined according to desired global aggregate (arithmetic/geometric mean, min, max, etc.)

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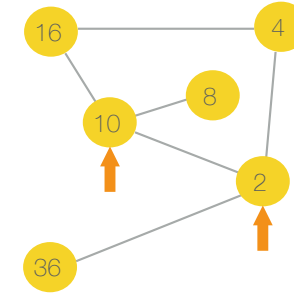
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## Aggregation Gossip framework instantiation

- Local variable  $S_p$  contains current estimate of the aggregate
- Need to give implementations for
  - $\text{SelectPeer}()$
  - $\text{Update}(S_p, S_q)$
- $\text{SelectPeer}()$  picks a random neighbor
  - $\text{Update}(S_p, S_q) = \frac{(S_p + S_q)}{2}$  (average)
  - $\text{Update}(S_p, S_q) = \sqrt{S_p S_q}$  (geometric mean)
  - $\text{Update}(S_p, S_q) = \max(S_p, S_q)$  (maximum)
- More complex functions built by combining the above

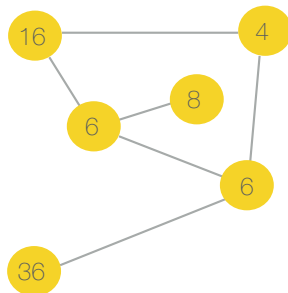
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## Aggregation Averaging example



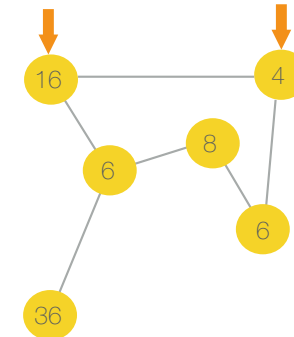
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## Aggregation Averaging example



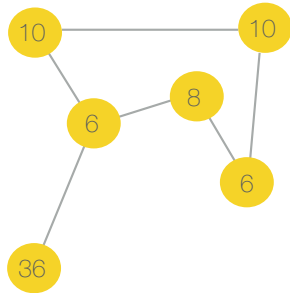
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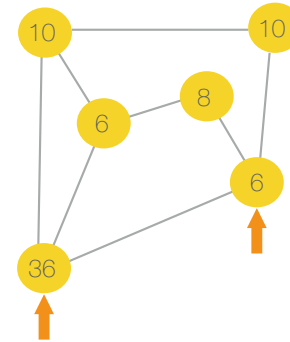
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## Aggregation Averaging example



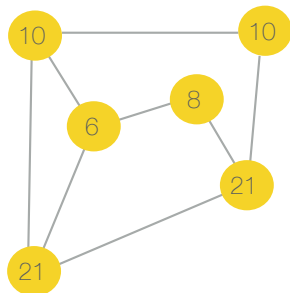
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## Aggregation Averaging example



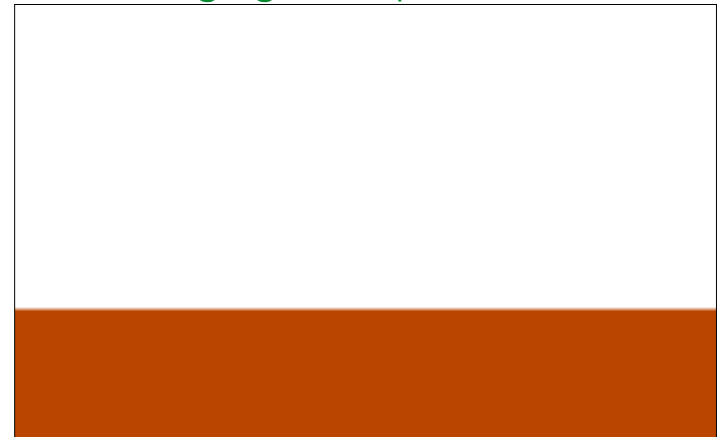
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## Aggregation Averaging example



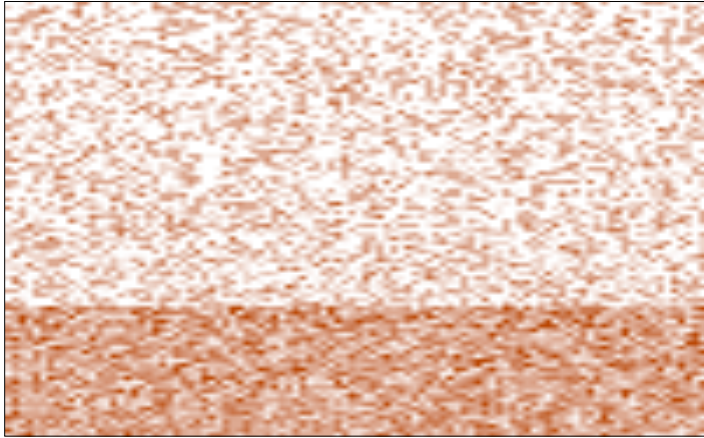
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## Aggregation Averaging example: Round 0



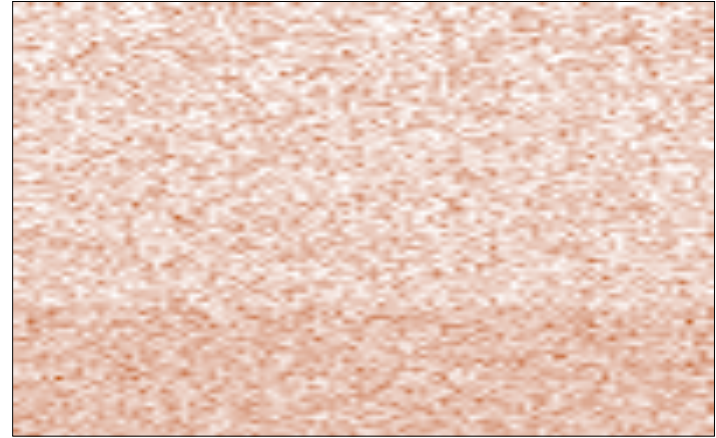
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Aggregation  
Averaging example: Round 1



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Aggregation  
Averaging example: Round 2



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Aggregation  
Averaging example: Round 3



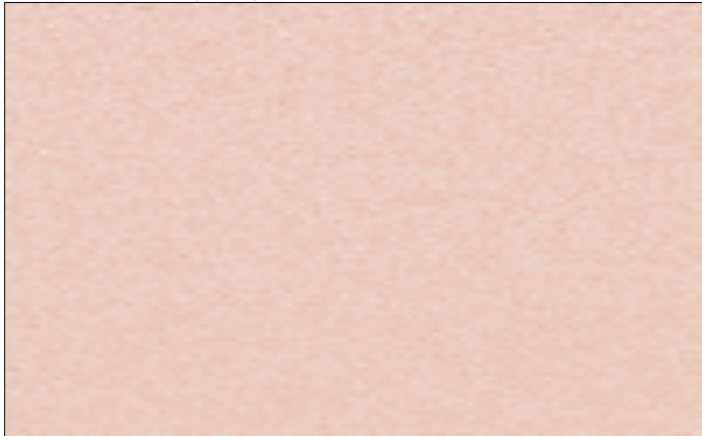
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Aggregation  
Averaging example: Round 4

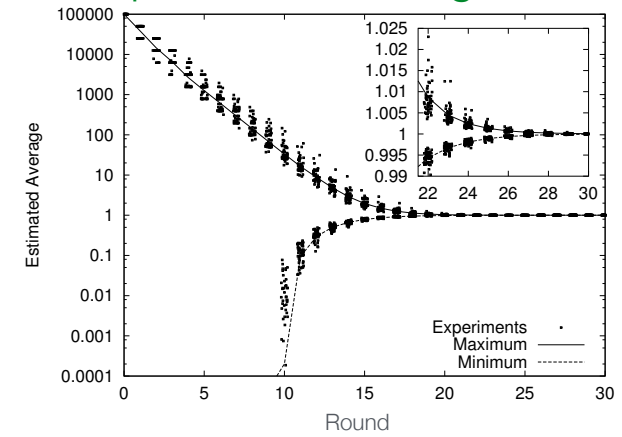


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## Aggregation Averaging example: Round 5



## Aggregation Exponential convergence

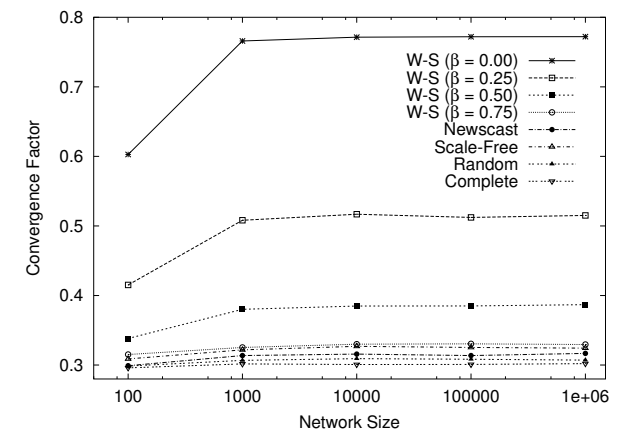


## Aggregation Exponential convergence

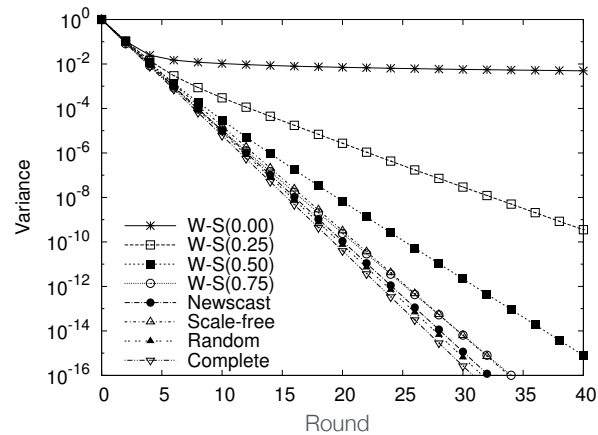
- In gossip-based averaging, if the selected peer is a *globally random* sample, then the expected variance among the estimates decreases exponentially
- Convergence factor:

$$\rho = \frac{E(\sigma_{i+1}^2)}{E(\sigma_i^2)} \approx \frac{1}{2\sqrt{e}} \approx 0.303$$

## Aggregation Convergence factor vs topology



## Aggregation Convergence factor vs topology



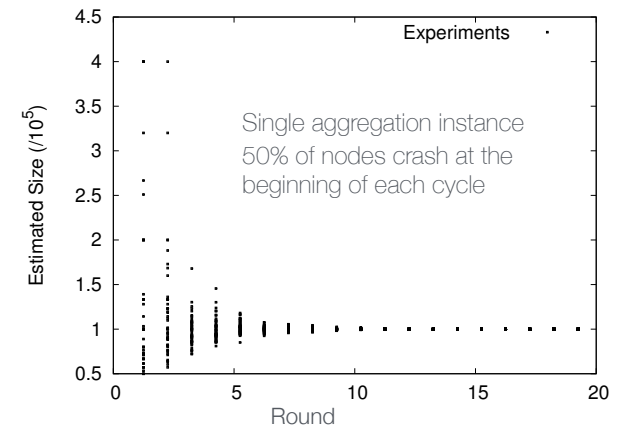
## Aggregation Network size estimation

- How can nodes within a network obtain an estimate for its current size?
- Just “freezing” the network and “counting” does not work
  - Not decentralized
  - Cannot deal with churn
  - Not scalable (think of the Internet)
- Idea: Base the size estimate on an aggregate value that can be computed through a decentralized algorithm
- Compute the arithmetic mean starting from zeroes at all nodes except one that holds 1
- The mean value is  $1/n$  and the network size  $n$  can be obtained simply by taking the inverse of the mean

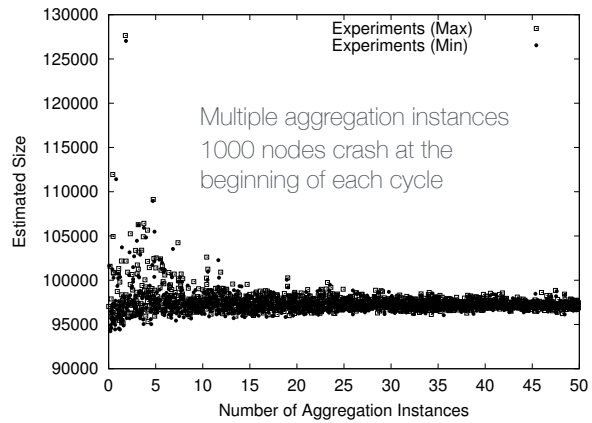
## Aggregation Network size estimation

- Idea works if there are no failures — no nodes fail and no messages are lost
- The estimate will be poor if failures occur during the early phases of the algorithm when the variance is greater
- Failures become less disruptive in later phases of the algorithm
- Worst-case failure scenario: the node with local value 1 fails immediately before exchanging local value with any node
- Idea: start multiple instances of the algorithm with different nodes holding the initial 1 value in each instance and average the results of the different instances to obtain the final estimate

## Aggregation Robustness — node failures

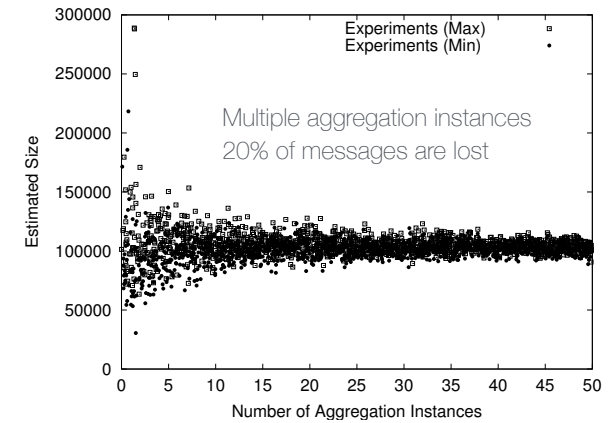


## Aggregation Robustness — node failures



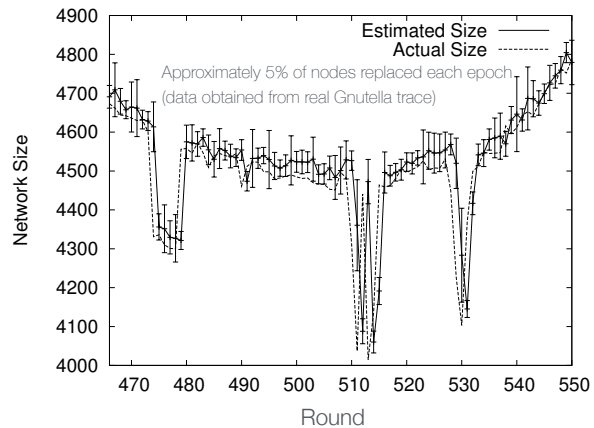
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## Aggregation Robustness — communication



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## Aggregation Robustness — churn



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