On Formal Choreographic Modelling: a Case Study in EU Business Processes

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Take-away message

Choregraphies & Correctness

Formal Choreographic Methods

aim to correctness-by-construction

by means of

syntactic/semantic restrictions on (global) specifications

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Our question

How do such restriction impact on "usability"?

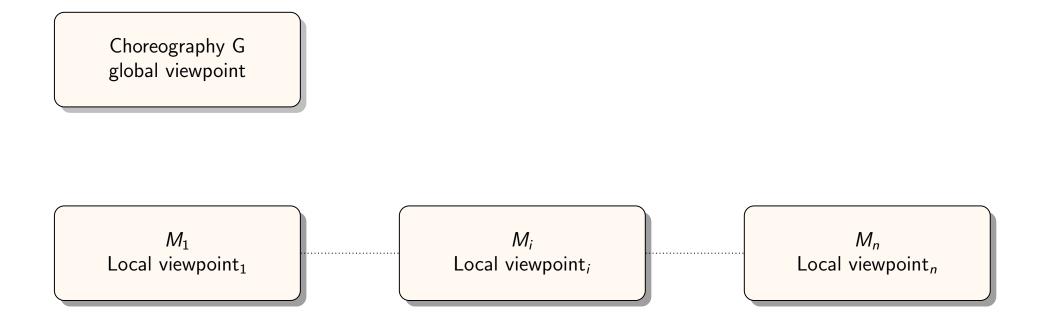
We address such question through

the application of a formal choreographic setting

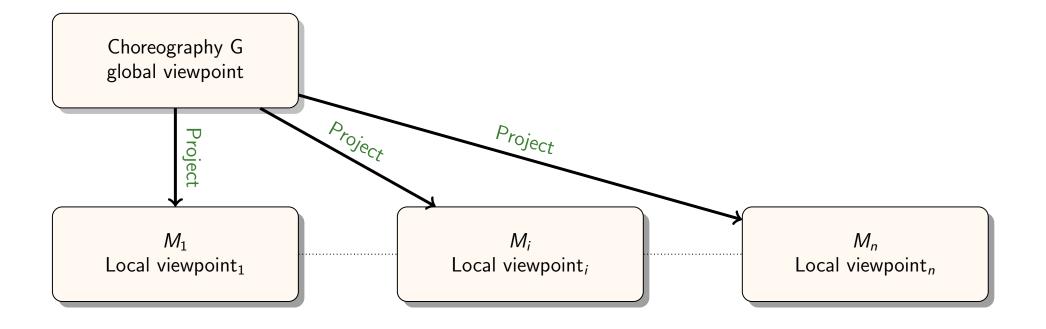
to the EU Custom business process models

and we draw some conclusions from this exercise

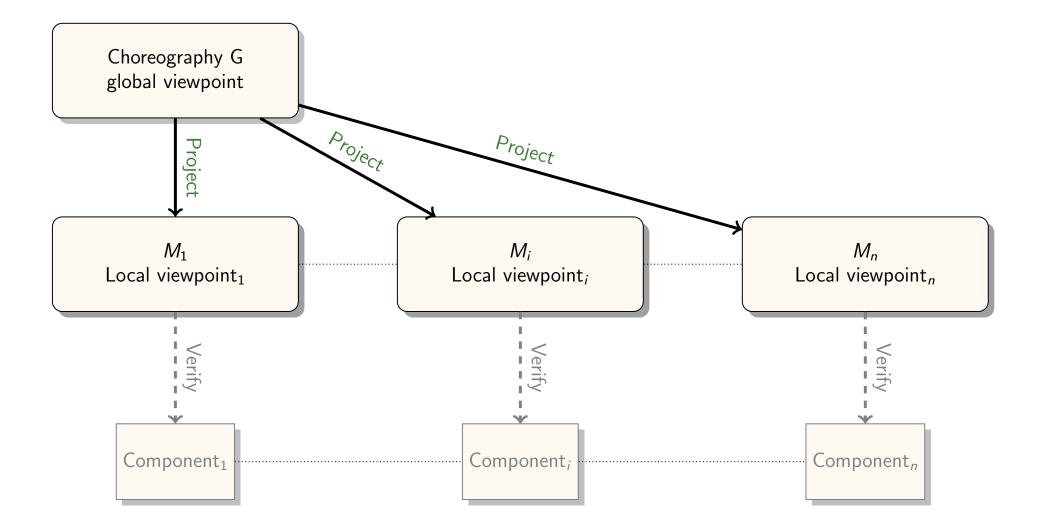
CbC with choreographies



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– Prelude –

Global specs, formally (I)

Global specs as regular expressions

Global choreographies (G-choreographies for short)

• kind of regular expressions

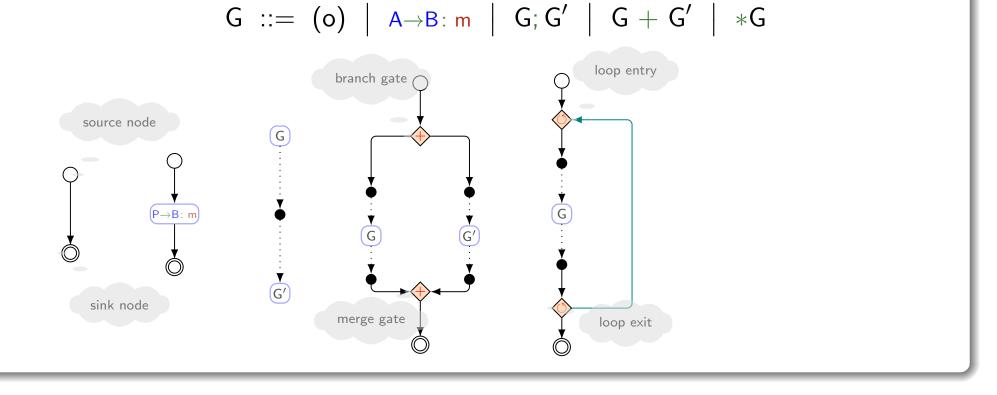
$$\mathsf{G} \ ::= \ (\mathsf{o}) \ \left| \begin{array}{c} \mathsf{A} {\rightarrow} \mathsf{B} \colon \mathsf{m} \end{array} \right| \ \mathsf{G} ; \mathsf{G}' \ \left| \begin{array}{c} \mathsf{G} + \mathsf{G}' \end{array} \right| \ *\mathsf{G}$$

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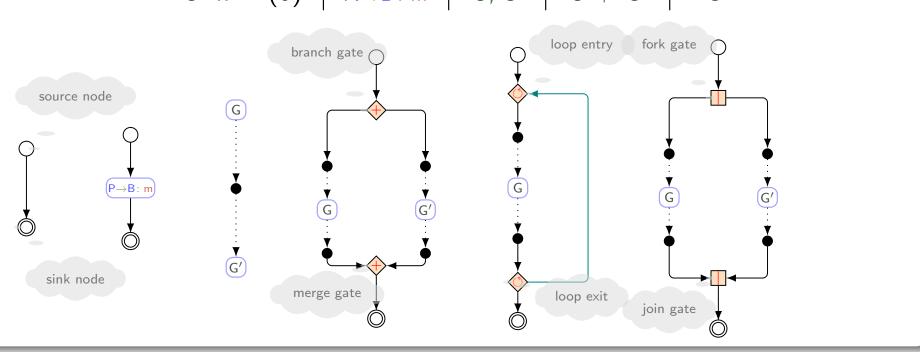


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Global choreographies (G-choreographies for short)

• kind of regular expressions



 $\mathsf{G} ::= (\mathsf{o}) \mid \mathsf{A} \rightarrow \mathsf{B} : \mathsf{m} \mid \mathsf{G}; \mathsf{G}' \mid \mathsf{G} + \mathsf{G}' \mid *\mathsf{G}$

On distributed choices

Assume asynchronous communication. In a branch $G_1 + G_2$

- there should be one active participant
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B is passive when

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- or B "unambiguously understands" which branch A opted for through the information received on each branch

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Well-branchedness

When the above holds true for each choice, the g-choreography is well-branched. This enables correctness-by-design.

(See [Tuosto & Guanciale JLAMP 2018] for the "greek symbols")

Class test

Which of the following global graphs is well-branched?

•
$$G_1 = A \rightarrow B$$
: int $+ A \rightarrow B$: str
• $G_2 = A \rightarrow B$: int $+ (o)$
• $G_2 = A \rightarrow B$: int $+ A \rightarrow C$: str

•
$$G_4 = \begin{pmatrix} A \rightarrow C: int; A \rightarrow B: bool \\ + \\ A \rightarrow B: bool; A \rightarrow C: bool \end{pmatrix}$$

<u>...</u>

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$$G_1 = A \rightarrow B$$
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$$G_3 = A \rightarrow B: int + A \rightarrow C: str$$

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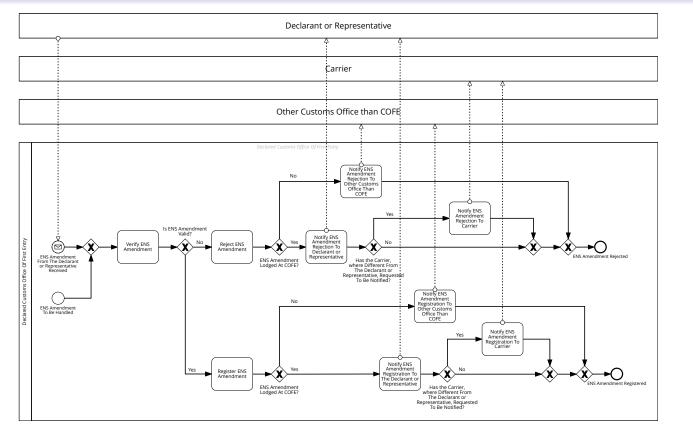
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Global specs as BPMN diagrams

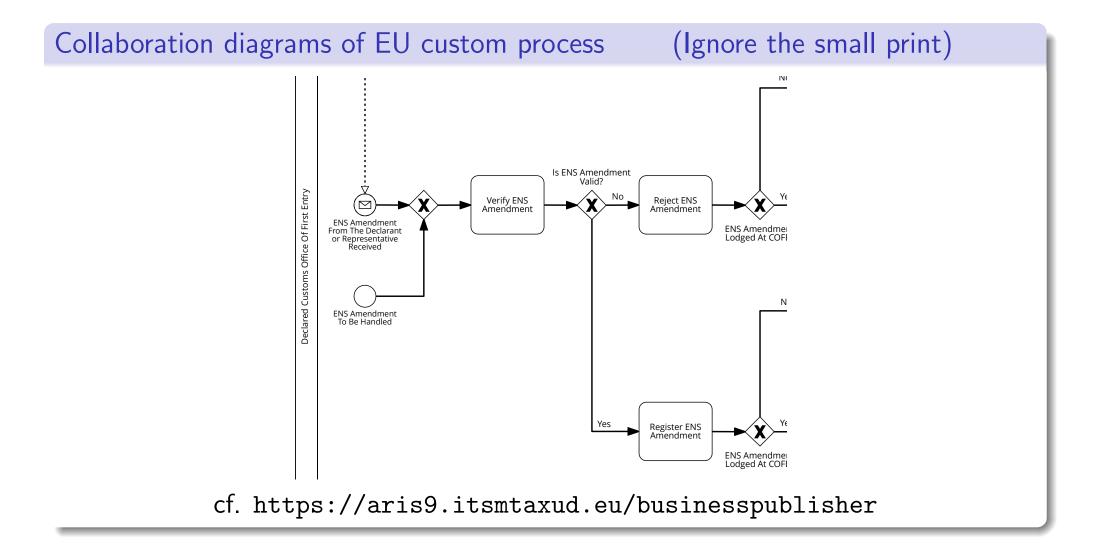
Collaboration diagrams of EU custom process

(Ignore the small print)

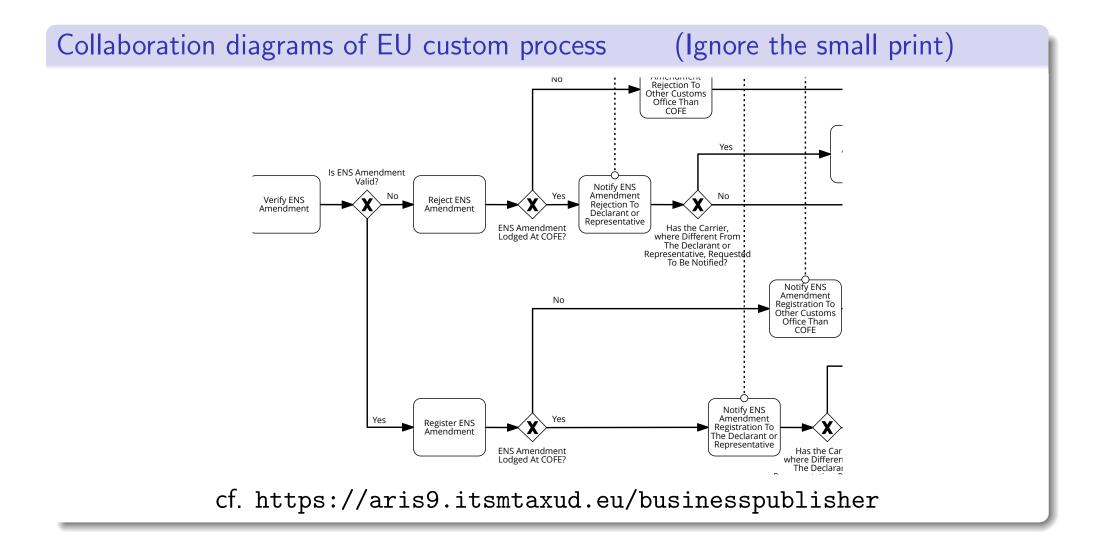


cf. https://aris9.itsmtaxud.eu/businesspublisher

Global specs as BPMN diagrams



Global specs as BPMN diagrams





A modelling exercise

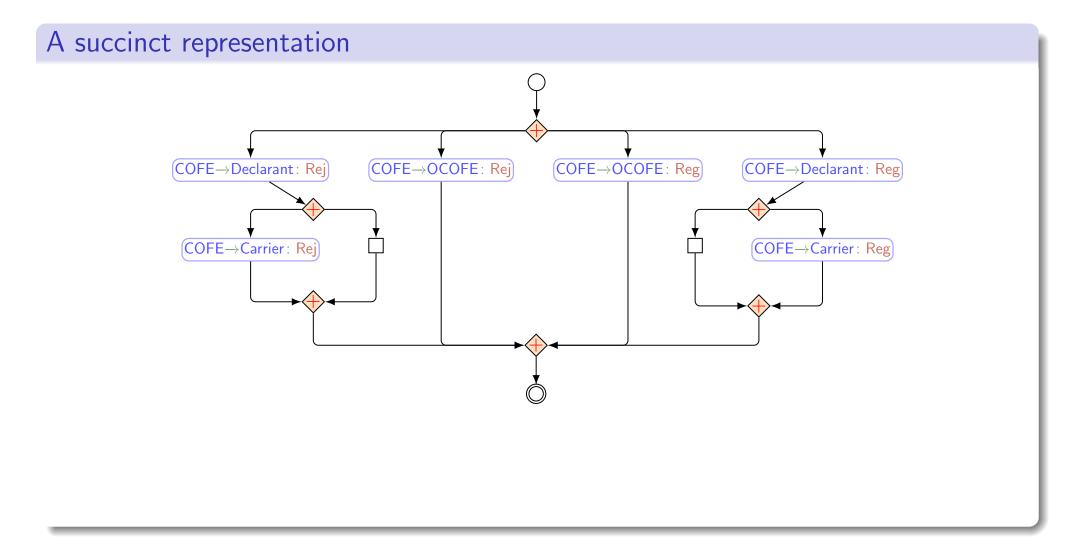
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Legal provisions

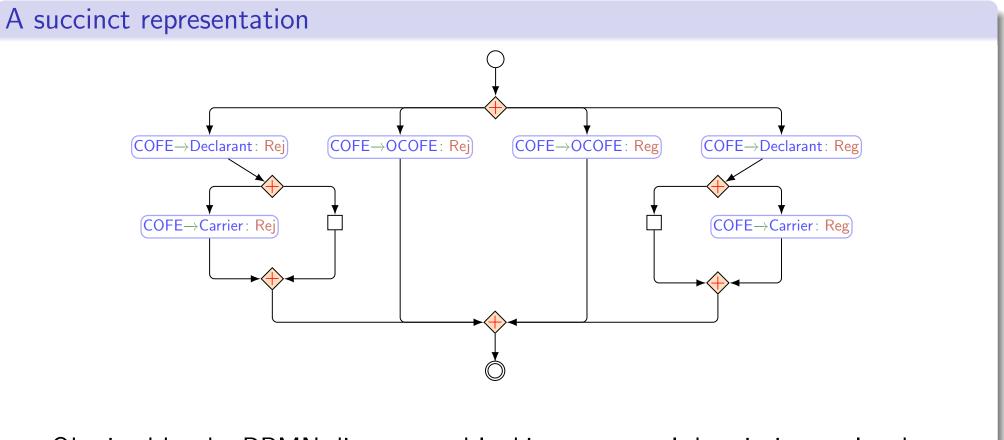
A workflow for importing goods into EU

- Declarants submit an *entry summary declaration* (ENS) (i.e., information for the import such as nature of the goods, carrier, etc.)
- The submission triggers several tasks which depend on multiple factors:
 - where was the ENS lodged?
 - is the good imported by road&rail, sea, air,...?
 - is the ENS valid?
 - register, if it is
 - request amendments & notify involved parties, otherwise

G-choreographies



G-choreographies



Obtained by the BPMN diagram and looking at textual descriptions only when something wasn't clear

Adopting formal models akin BPMN helps

Some advantages of g-choreographies

- G-choreographies are more compact and their visusal representation can be understood by lay stakeholders
- G-choreographies seem clearer than the BPMN spec...
 (I'm biased, of course)
- Other formalisms may be less clear (e.g., those relying on process algebras)

Its great to start from BPMN diagrams, but...

Coping with underspec

- maintaining the correspondence could be problematic E.g., the discussion on loops in the paper
- often sharing of relevant information implicit in value passing E.g., notification to carrier is presumably required in the ENS after carrier & declarant struck a deal
- sticking to the informal specification increases non-determinism (like in our model of the amendment process)

– Epilogue –

Conclusions

Expressivness vs. Correctnes-by-construction

- specs usually violate the conditions for correctness required by formal models E.g., well-branchedness is broken by *optional communications* such as the notifications from COFE to Carrier
- Mitigations:
 - extra interactions instead of relying on value passing
 - formalisation drifts away from BPMN specs
 - but help identifying ambiguities / lack of precision
 - use more general conditions
 - models different from BPMN
 - verification of conditions is more expensive

Thank you!

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If you are looking for job, get in touch

