



Reversible concurrent systems

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GNCS project
Sistemi Reversibili Concorrenti:
dai Modelli ai Linguaggi

Roadmap

- Reversible computing
- Debugging Erlang programs
- Petri nets vs event structures
- Conclusion



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What is reversibility?

The possibility of executing a computation both in the standard, forward direction, and in the backward direction, going back to a past state

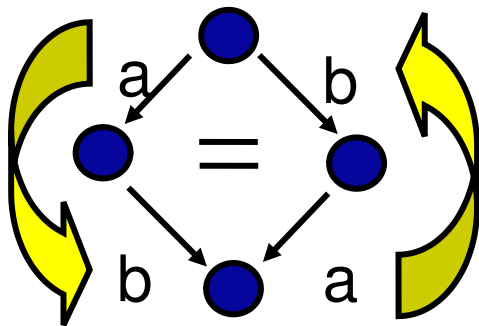
- In some areas systems are naturally reversible: biology, quantum computing, ...
- In other areas making systems reversible can be useful
 - Undo option in text editors
 - Debugging, robotics, ...

Reversibility in concurrent systems

- Reversibility in a sequential setting:
recursively undo the last action
- In concurrent systems execution of different actions
may overlap in time
 - No uniquely defined last action
 - Actions form a partial order, not necessarily a total
order
- Different approaches to reversibility exist
- We follow **causal-consistent reversibility**
[Danos & Krivine, CONCUR 2004]

Causal-consistent reversibility

- Based on causality instead of time
- Causal dependencies must be respected
 - First reverse the consequences, then the causes
- Independent actions are reversed independently



Reversibility in our project

- We considered many aspects of reversibility, both foundational and applicative
- In the talk I will focus on two contributions
 - How to apply causal-consistent reversibility to the debugging of Erlang programs?
 - Ongoing effort since many years
 - By myself, C. Sacerdoti Coen, G. Vidal, ...
 - How to extend the relation between event structures and Petri nets to the reversible setting?
 - By C. A. Mezzina, G. M. Pinna, E. Melgratti, I. Ulidowski

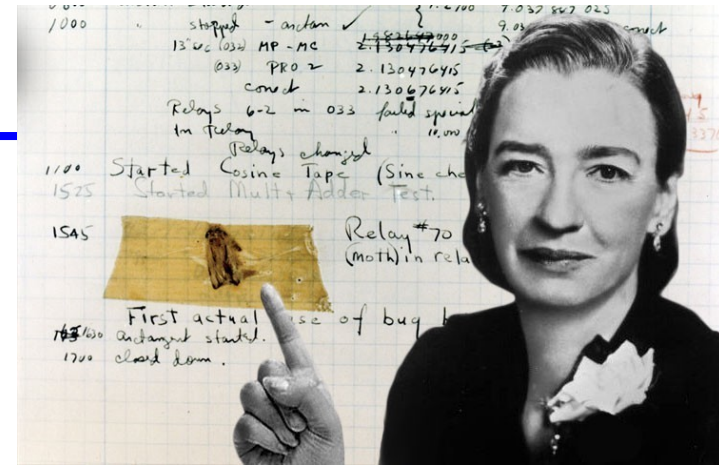
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Debugging

- Debugging amounts to find the wrong line of code (bug) causing a visible misbehavior
- The bug **precedes** and **causes** the misbehavior
 - Quite natural to use reversibility to go back from the misbehavior to the bug



Reversible debugging

- Sequential reversible debugging is well understood
 - Gdb (since 2009), Microsoft time-travel debugger, ...
- Concurrent reversible debugging not so developed
 - Most approaches just linearize the execution
 - Causal information is lost
- Can we use causal-consistent reversibility?

Causal-consistent debugging

- Introduced in [Giachino, Lanese & Mezzina, FASE 2014]
- Allows one to explore a concurrent computation back and forward
 - Any action can be undone provided its consequences have been undone beforehand
- Which action to undo can be selected by the user or by a scheduler
- But we can do better

Debugging and causality

- Standard debugging procedure:
 1. Observing an unexpected behavior
 2. Finding in the code the instruction that **caused** it
 3. Correcting the instruction
- Causal-consistent reversibility naturally tracks lot of causal information
- This information can be used to drive step 2 above
- Debugging strategy: follow causality links backward from the misbehavior to the bug
- We can use the roll operator to this end

Causal-consistent debugging: roll

- The roll operator allows one to undo a selected past action, including all and only its consequences
- Minimal set of undos needed to undo the selected action in a causal-consistent way
- Many interfaces for it:
 - N actions in a given process
 - Last assignment to a given variable
 - Send of a given message

Causal-consistent roll at work

- The programmer executes the program and finds some unexpected behavior
- The roll allows him to find automatically the instruction that immediately caused the misbehavior
- Two possibilities:
 - The found instruction is wrong: **bug found**
 - The found instruction gets wrong data from previous instructions: **iterate**
- One can explore the tree of causes, navigating from one process to the other

- Causal-consistent Debugger for Erlang
- Applies the approach outlined above to a fragment of Erlang
 - Functional and concurrent language
 - Based on message passing
 - Used in mainstream applications such as some versions of Facebook chat
- <https://github.com/mistupv/cauder-v2>
- Support for further (non trivial) constructs in Erlang has been added during the project

CauDEr interface

The screenshot displays the CauDEr interface with the following components:

- Code:** A code editor showing Erlang code. Line 26, `XPid!{read,self()},` is highlighted in green. The code includes `varManager(Val).`, `incrementer(MePid,XPid) ->`, `MePid!{request,self()},`, `receive answer ->`, `XPid!{read,self()},`, `receive X ->`, `XPid!{write,X+1},`, and `MePid!{release} end end.`
- Process Info:** A section containing:
 - Bindings:** A table with columns 'Name' and 'Value'.

Name	Value
MePid	97
XPid	98
 - Stack:** A list of stack frames: `receive` and `meViolation:incrementer/2`.
- Log:** A list of log entries: `send(3)`, `rec(4)`, `send(5)`, and `send(6)`.
- History:** A list of history entries: `rec(answer,1)` and `send({request,99},0)`.

Actions: A control panel for process `PID: 99 - meViolation:incrementer/2`. It features tabs for `Manual`, `Automatic`, `Replay`, and `Rollback`. The `Replay` tab is active, showing:

- Steps:** A dropdown menu set to `1` and a `Roll steps` button.
- PID:** An input field and a `Roll spawn` button.
- Msg. Uid:** An input field set to `3` and a `Roll send` button.
- Msg. Uid:** An empty input field and a `Roll receive` button.
- Var.:** An empty input field and a `Roll variable` button.

System Info: A section with a `Mail` table:

Dest.	Value	UID
97	{request,100}	2

Trace: A section with a `Roll Log` tab showing a scrollable list of events:

- `Roll send from Proc. 99 of {release} to Proc. 97 (6)`
- `Roll send from Proc. 97 of answer to Proc. 100 (7)`

Status Bar: At the bottom, it displays `Rolled back sending of message with UID: 3` on the left and `Ln 26, Col 1 Alive 4, Dead 1` on the right.

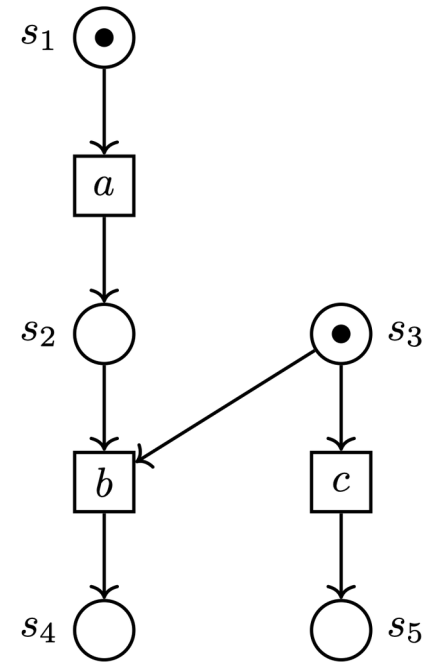
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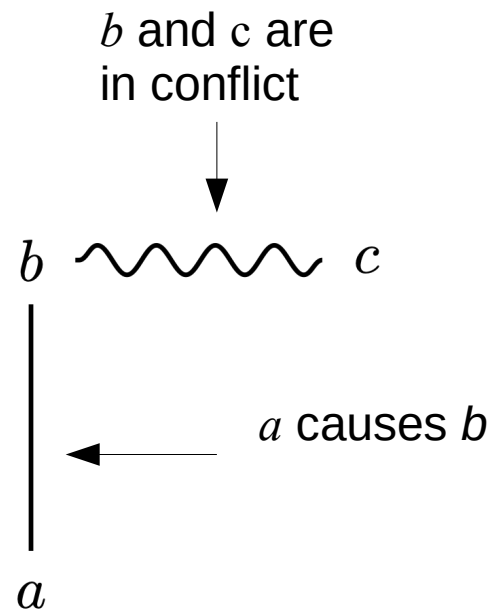
Petri nets

- Operational model for concurrency
- Based on tokens that enable transitions
- The one in figure is actually an occurrence net (ON)
 - No cycles and max 1 token per place
 - It is always clear where a token comes from
 - Plays well with reversibility
 - Can represent any net via unfolding



Event structures

- Denotational model for concurrency
- Based on events and relations among them
- These are actually prime event structures (PES)
 - Only causation and conflict relations
 - Variants have other relations
- Prime event structures correspond to occurrence nets
 - Classical result by Winskel

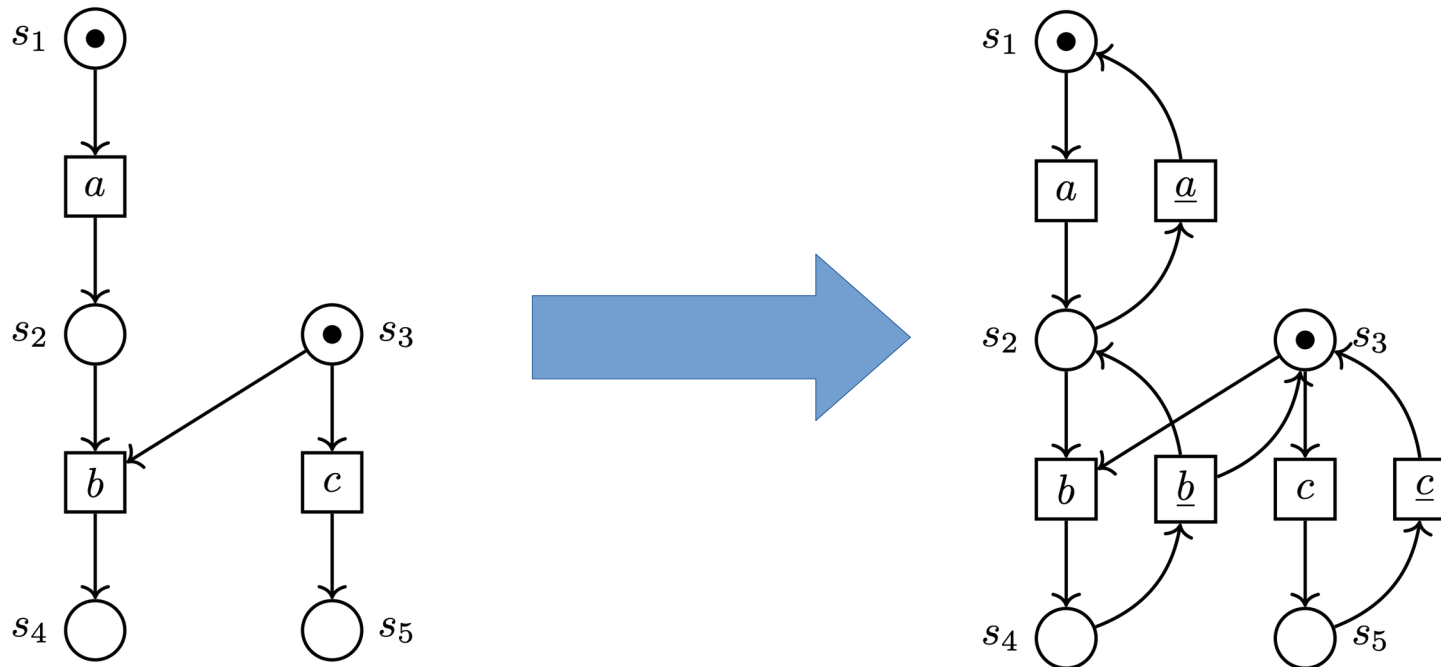


Reversible PES

- Extend PES with reversibility
 - Not only causal-consistent
- Not all the events need to be reversible
- Introduce:
 - Reverse causality \triangleleft : a forward event is a cause of a backward one
 - Prevention \triangleright : a forward event forbids a backward one

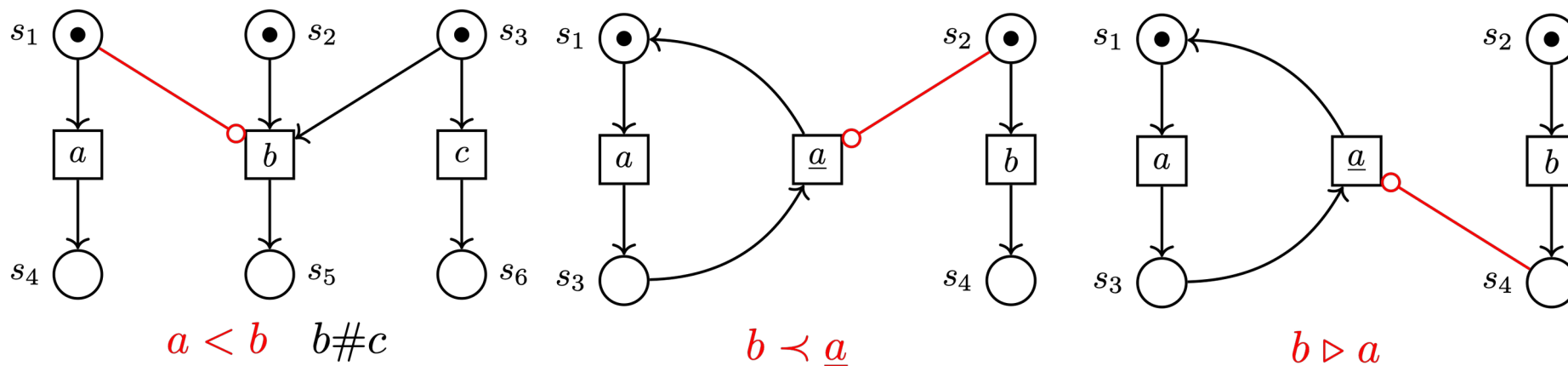
From ON to reversible ON

- Adding reverse transitions in ON gives rise to causal-consistent reversibility

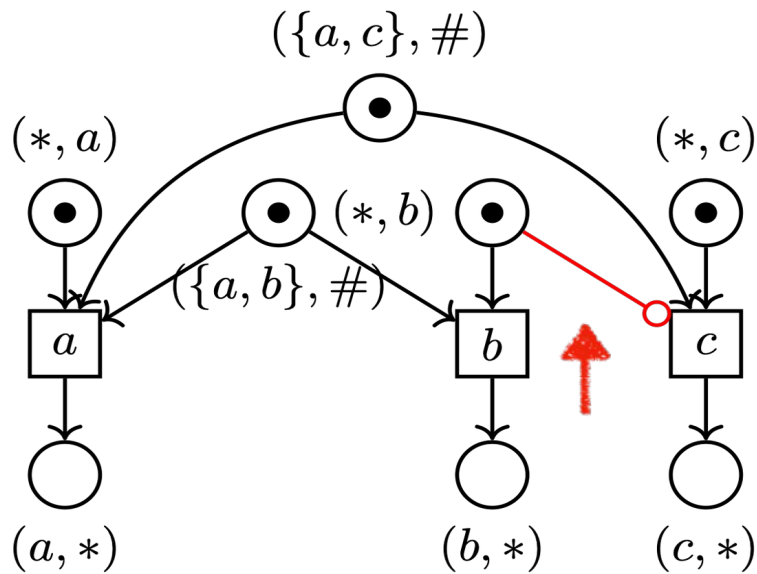
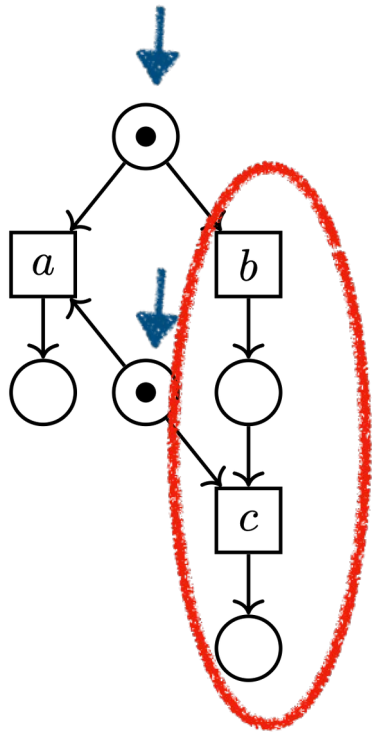


Modeling rPES with Petri nets

- We need to introduce inhibitor arcs
- Inhibitor arcs can model causality as well
- We can reduce to “flat” nets
- We call them Causal Nets (CN)
- Reversible Causal Nets (rCN) extend causal nets with reverse actions and additional inhibitor arcs

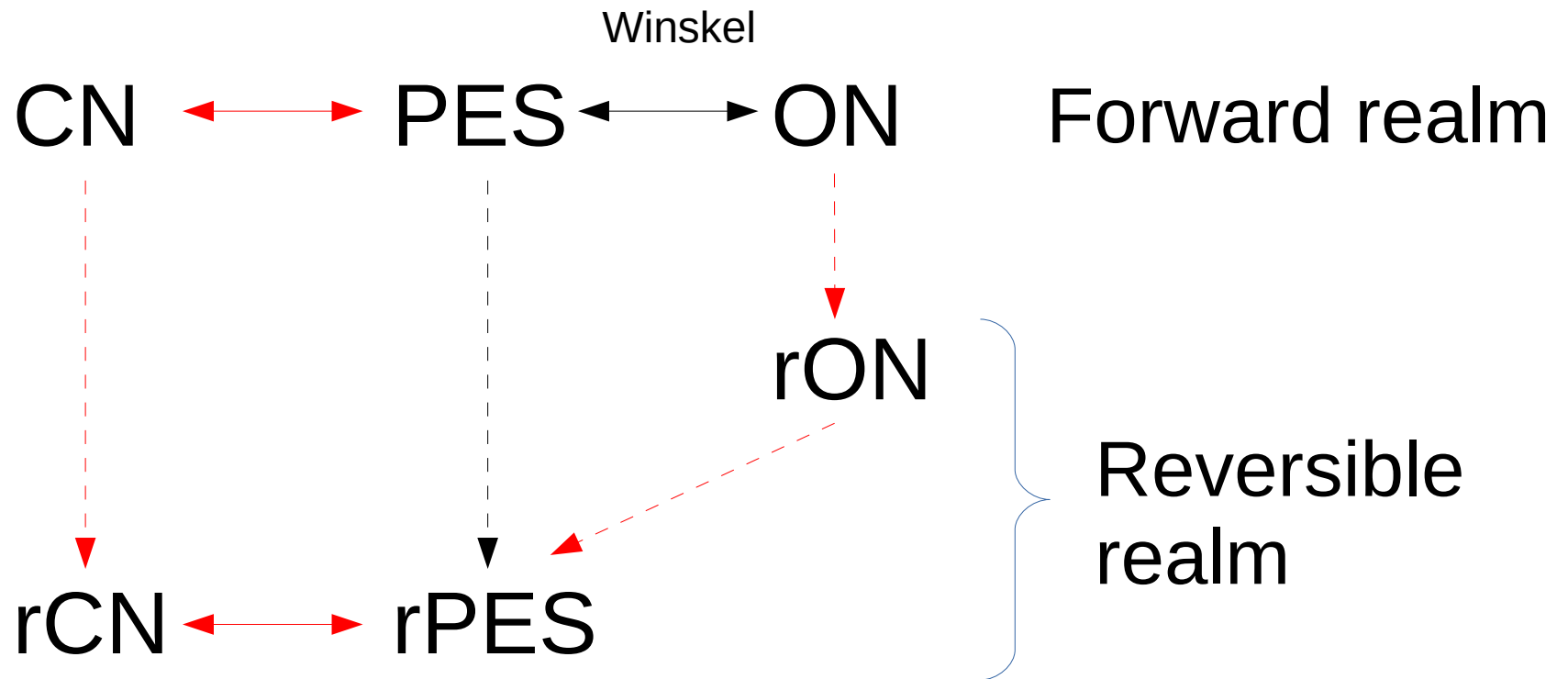


From ON to CN



$a \# b$ $a < c$
 $a \# c$

Graphical summary of results



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Other contributions

- Understanding the interplay between reversibility and time
- Developing techniques for axiomatic reasoning on reversible processes
- Understanding reversibility in Markov chains

Conclusion

- Reversibility is a niche area with many possible applications and open questions
- Studied by both computer scientists and mathematicians (and not only)
- If curious, Reversible Computation conference will be held in Urbino, on July 5-6
 - Online participation will also be allowed

Thanks!

Questions?