WG4 Case Study
Debugging of Concurrent and Distributed Systems

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Roadmap

- Reversible debugging
- State of the art: sequential debugging
- Debugging concurrent/distributed systems: causal-consistent reversible debugging
- Future directions
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Why debugging?

- Developers spend 50% of their programming time finding and fixing bugs
- The global cost of debugging has been estimated in $312 billions annually
- The cost of debugging is bound to increase with the increasing complexity of software
  - Size
  - Concurrency, distribution
- Surprisingly, a very little amount of research concentrates on debugging
Standard debugging strategy

- When a failure occurs, one has to re-execute the program with a breakpoint before the expected bug.
- Then one executes step-by-step forward from the breakpoint, till the bug is found.
- Limitations:
  - High cost of replaying
    » Time, use of the actual execution environment
  - Difficult to precisely replay the execution
    » Concurrency or non-determinism
  - Difficult to find the exact point where to put the breakpoint
    » If the breakpoint is too late, the execution needs to be redone
    » Frequently many attempts are needed.
Reversibility for debugging

- Reversible debuggers extend standard debuggers with the ability to execute the program under analysis also backward.
- Avoids the common “Damn, I put the breakpoint too late” error
  - Just execute backward from where the program stopped or where a wrong result appeared till the desired point is reached.
- The overhead due to storing history information is a main limitation for reversible debuggers.
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State of the art: sequential debugging

- Reversible debuggers exist
  - GDB, UndoDB
- Many reversible debuggers deal only with sequential programs
- Some of them allow one to debug concurrent programs
  - They register scheduler events
  - The same scheduling is used when the program is replayed
Sequential reversible debugging strategy

- Take an execution containing a failure and move backward and forward along it looking for the bug
- The exact same execution can be explored many times forward and backward
  - Non-determinism is no more a problem
A reversible debugger: GDB

- GDB supports reversible debugging since version 7.0 (2009)
- Uses record and replay
  - One activates the recording modality
  - Executes the program forward
  - Can explore the recorded execution backward and forward
  - When exploring, instructions are not re-executed
GDB reverse commands

- Like the forward commands (step, next, continue), but in the backward direction
- Reverse-step: goes back to the last instruction
- Reverse-next: goes back to the last instruction, does not go inside functions
- Reverse-continue: runs back till the last stop event
- ...
- Breakpoints and watchpoints can be used also in the backward direction
A commercial reversible debugger: UndoDB

- Already presented at the Grenoble meeting
- From UndoSoftware, Cambridge, UK
  http://undo-software.com/
- Built as an extension of GDB
- Available for Linux and Android
- Allows reversible debugging for programs in C/C++
UndoDB commands

- Close to GDB commands
- Some more high-level commands and configuration commands
- Commands to write a recorded execution to file, and reload it
  - Useful to record on client premises and explore at company premises
UndoDB winning feature

**Performance**

- Comparison with GDB, on recording gzipping a 16MB file

<table>
<thead>
<tr>
<th></th>
<th>Native</th>
<th>UndoDB</th>
<th>GDB</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Time</strong></td>
<td>1.49 s</td>
<td>2.16 s (1.75 x)</td>
<td>21 h (50000 x)</td>
</tr>
<tr>
<td><strong>Space</strong></td>
<td>-</td>
<td>17.8 MB</td>
<td>63 GB</td>
</tr>
</tbody>
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Causal-consistent reversible debugging

- We are interested in debugging concurrent/distributed programs
- Since [Danos&Krivine, CONCUR 2004] the notion of reversibility for concurrent systems is causal-consistent reversibility
  - Any action can be undone, provided that its consequences (if any) have been undone
  - Concurrent actions can be undone in any order, but causal-dependent actions are undone in reverse order
- At any point, many actions can be undone
Debugging and causality

- Causal-consistency relates backward computations with causality
- Debugging amounts to find the bug that caused a given misbehavior
- We propose the following debugging strategy: follow causality links backward from misbehavior to bug
- Which primitives do we need to enable such a strategy?
A proposal: the **roll** primitive

- The main primitive we propose is **roll t n**
- Undoes the last **n** actions of thread **t**...
- ... in a causal-consistent way
  - Before undoing an action one has to undo all (and only) the actions depending on it
- A single **roll** may cause undoing actions in many threads
Different interfaces for **roll**

- One interface for each possible misbehavior
  - This depends on the language
- Examples are:
  - **Wrong value in a variable**: `rollvariable id` goes to the state just before the last change of variable `id`
  - **Unexpected thread**: `rollthread t` undoes the creation of thread `t`
Using roll-like primitives

- The programmer can follow causality links backward
- No need for the programmer to know which thread or instruction originated the misbehavior
  - The primitives find them automatically
- The procedure can be iterated till the bug is found
- Only relevant actions are undone
  - Thanks to causal consistency
- Looking at which threads are involved gives useful information
  - If an unexpected thread is involved, an interference between the two threads has happened
CaReDeb: a causal-consistent debugger

- Only a prototype to test our ideas
- Debugs programs in the μOz language
  - Toy functional language with threads and asynchronous communication via ports
- Written in Java
- Available at [http://www.cs.unibo.it/caredeb](http://www.cs.unibo.it/caredeb)
- Description and underlying theory in [Giachino, Lanese & Mezzina, FASE 2014]
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Future directions

- Enable causal-consistent debugging of real languages
  - Many constructs
  - One needs to understand their causal semantics
  - Large implementation effort

- Is the **roll** primitive good?
  - Which is the impact on actual debugging?
  - It would be interesting to setup an experiment

- Are there other useful primitives?
UndoSoftware interests

• Debugging concurrent and distributed systems is challenging and requires a huge effort
  − They are not interested in solutions which are not efficient
  − They will not tackle it in the near future
• A main interest is parallel recording
• Debugging of MPI programs is also of interest
Finally

Thanks!

Questions?
Our target language: μOz

- A kernel language of Oz
- Oz is at the base of the Mozart language
- Thread-based concurrency
- Asynchronous communication via ports
- Shared memory
  - Variable names are sent, not their content
- Variables are always created fresh and never modified
- Higher-order language
  - Procedures can be communicated