

# Retractable Contracts



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# Map of the talk

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- Why retractable contracts?
- What is a retractable contract?
- Results
- Conclusion



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# Contracts

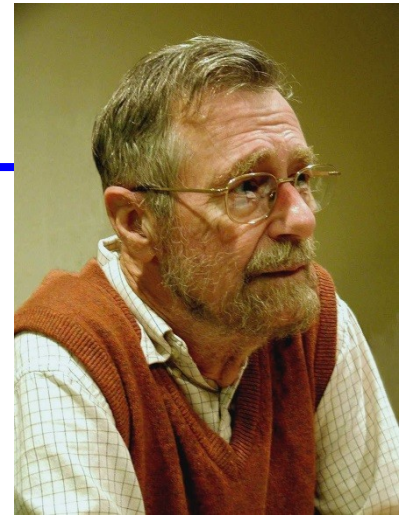
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- A **contract** is the abstract description of the behavior of either a client or a server
- A client **complies** with a server if all her requirements are fulfilled
  - by reaching a distinguished satisfaction state or
  - by running an infinite interaction without ever getting stuck
- A client that does not comply with its server may get stuck
- Compliance is statically decidable

# Undoing things considered harmful

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- Undo operations are useful and widespread
  - Undo command in your favorite editor
  - Restore a past backup
  - Back button in your favorite browser
- In interacting systems (unilateral) undo may lead to unpredictable or undesired results
  - What happens if you press the back button when reserving a flight?
  - You don't want a client to undo her payment after a purchase
- Undo activities must be disciplined
- Retractable contracts are a way to discipline activities including undo operations

# Retractable contracts: approach

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- Getting stuck may depend on wrong choices taken during the interaction
- Going back to past choices and trying different paths may solve the problem
- This will “facilitate” compliance
- In this work we explore a notion of contracts where past decisions are stored and can be undone

# Map of the talk

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# Retractable contracts: syntax

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$\sigma ::=$	$1$	success	} Standard contracts
	$\bigoplus_{i \in I} \bar{a}_i . \sigma_i$	internal output choice	
	$\sum_{i \in I} a_i . \sigma_i$	external input choice	
	$X$	variable	
	$\text{rec } X . \sigma$	recursion	
	$\sum_{i \in I} \bar{a}_i . \sigma_i$	retractable output choice	
	$\bigoplus_{i \in I} a_i . \sigma_i$	internal input choice	



# Retractable contracts: main idea

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- The peculiar operator is retractable output choice:

$$\sum_{i \in I} \bar{a}_i \cdot \sigma_i$$

- It behaves as follows:
  - it performs an output, but other options are stored
  - if the computation gets stuck, undo is performed and another option is tried

# Retractable contracts: history information

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- To give semantics to contracts we need history information
- We add  $\circ$  (empty contract) to contracts  $\sigma$
- Histories are stacks of contracts  $h ::= [] \mid h:\sigma$
- Contracts with history:  $h \prec \sigma$

# Motivating problem

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- A buyer wants to buy either a bag or a belt
- She will decide whether to pay by card or cash after knowing the price
- Buyer =  
 $\overline{\text{bag.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})} \oplus \overline{\text{belt.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})}$
- The seller accepts cards only for bags, not for belts
- Seller =  
 $\overline{\text{bag.price.}(\overline{\text{card}} + \overline{\text{cash}})} + \overline{\text{belt.price.cash}}$
- Buyer and seller are not compliant

# Reversibility to the rescue

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- Buyer =  
 $\overline{\text{bag.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})} \oplus \overline{\text{belt.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})}$
- Seller =  
 $\text{bag.}\overline{\text{price.}(\text{card} + \text{cash})} + \text{belt.}\overline{\text{price.cash}}$
- They become compliant if we make the buyer choice between bag and belt retractable
  - Or the one between card and cash (for belt)
- The buyer is still able to pay a belt with card if interacting with a seller allowing this

# Reversibility to the rescue

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- They become compliant if we make the buyer choice between bag and belt retractable
  - Or the one between card and cash (for belt)
- The buyer is still able to pay a belt with card if interacting with a seller allowing this

# Sample computation

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● Buyer' =

$$[] \prec \overline{\text{bag.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})} + \overline{\text{belt.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})}$$

● Seller =

$$[] \prec \overline{\text{bag.price.}(\text{card} + \text{cash})} + \overline{\text{belt.price.cash}}$$

# Sample computation

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- Buyer' =

$$[] \prec \overline{\text{bag.price}}.(\overline{\text{card}} \oplus \overline{\text{cash}}) + \overline{\text{belt.price}}.(\overline{\text{card}} \oplus \overline{\text{cash}})$$

- ▶  $\overline{\text{bag.price}}.(\overline{\text{card}} \oplus \overline{\text{cash}}) \prec \overline{\text{price}}.(\overline{\text{card}} \oplus \overline{\text{cash}})$

- Seller =

$$[] \prec \overline{\text{bag.price}}.(\text{card} + \text{cash}) + \overline{\text{belt.price}}.\text{cash}$$

- ▶  $\overline{\text{bag.price}}.(\text{card} + \text{cash}) \prec \overline{\text{price}}.\text{cash}$

# Sample computation

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- ▶  $\overline{\text{bag.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})} : \circ \prec \overline{\text{card}} \oplus \overline{\text{cash}}$

- Seller =

$$[] \prec \overline{\text{bag.price.}(\text{card} + \text{cash})} + \overline{\text{belt.price.cash}}$$

- ▶  $\overline{\text{bag.price.}(\text{card} + \text{cash})} \prec \overline{\text{price.cash}}$

- ▶  $\overline{\text{bag.price.}(\text{card} + \text{cash})} : \circ \prec \text{cash}$



# Sample computation

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- ▶  $\overline{\text{bag.price.}(\text{card} + \text{cash})} : \circ \prec \text{cash}$

# Sample computation

---

- Buyer' =  
 $\overline{\text{bag.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})} : \circ \prec \overline{\text{card}}$
  
- Seller =  
 $\overline{\text{bag.price.}(\text{card} + \text{cash})} : \circ \prec \text{cash}$

# Sample computation

---

- Buyer' =  
 $\overline{\text{bag.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})} : \circ \prec \overline{\text{card}}$
- ▶  $\overline{\text{bag.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})} \prec \circ$

- Seller =  
 $\overline{\text{bag.price.}(\text{card} + \text{cash})} : \circ \prec \text{cash}$
- ▶  $\overline{\text{bag.price.}(\text{card} + \text{cash})} \prec \circ$

# Sample computation

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- Buyer' =
  - $\overline{\text{bag.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})} : \circ \prec \overline{\text{card}}$
  - ▶  $\overline{\text{bag.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})} \prec \circ$
  - ▶  $[] \prec \overline{\text{bag.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})}$
  
- Seller =
  - $\overline{\text{bag.price.}(\text{card} + \text{cash})} : \circ \prec \text{cash}$
  - ▶  $\overline{\text{bag.price.}(\text{card} + \text{cash})} \prec \circ$
  - ▶  $[] \prec \overline{\text{bag.price.}(\text{card} + \text{cash})}$

# Sample computation

---

- Buyer' =  
 $[\ ] \prec \overline{\text{bag.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})}$
  
- Seller =  
 $[\ ] \prec \overline{\text{bag.price.}(\text{card} + \text{cash})}$

# Sample computation

---

- Buyer' =

- []  $\prec$   $\overline{\text{bag.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})}$

- ▶ ○  $\prec$   $\overline{\text{price.}(\overline{\text{card}} \oplus \overline{\text{cash}})}$

- Seller =

- []  $\prec$   $\overline{\text{bag.price.}(\text{card} + \text{cash})}$

- ▶ ○  $\prec$   $\overline{\text{price.}(\text{card} + \text{cash})}$

# Sample computation

---

- Buyer' =

- $\prec \overline{\text{bag.price.}(\overline{\text{card}} \oplus \overline{\text{cash}})}$

- ▶ ○  $\prec \overline{\text{price.}(\overline{\text{card}} \oplus \overline{\text{cash}})}$

- ▶ ○ : ○  $\prec \overline{\text{card}} \oplus \overline{\text{cash}}$

- Seller =

- $\prec \overline{\text{bag.price.}(\text{card} + \text{cash})}$

- ▶ ○  $\prec \overline{\text{price.}(\text{card} + \text{cash})}$

- ▶ ○ : ○  $\prec \text{card} + \text{cash}$

# Sample computation

---

● Buyer' =

$$\circ : \circ \prec \overline{\text{card}} \oplus \overline{\text{cash}}$$

● Seller =

$$\circ : \circ \prec \text{card} + \text{cash}$$



# Sample computation

---

● Buyer' =

○ : ○  $\prec$   $\overline{\text{card}} \oplus \overline{\text{cash}}$

▶ ○ : ○  $\prec$   $\overline{\text{card}}$

● Seller =

○ : ○  $\prec$   $\text{card} + \text{cash}$

# Sample computation

---

- Buyer' =

- : ○  $\prec \overline{\text{card}} \oplus \overline{\text{cash}}$

- ▶ ○ : ○  $\prec \overline{\text{card}}$

- ▶ ○ : ○ : ○  $\prec 1$

- Seller =

- : ○  $\prec \text{card} + \text{cash}$

- ▶ ○ : ○ :  $\text{cash} \prec 1$

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# Compliance

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- The compliance relation  $h \prec \sigma \dashv\| k \prec \rho$  holds iff  $h \prec \sigma \parallel k \prec \rho \rightarrow^* h' \prec \sigma' \parallel k' \prec \rho' \nrightarrow$  implies  $\sigma' = 1$ 
  - If the computation stops then the client is satisfied
- The compliance relation on contracts is obtained by executing them with an empty history

# Compliance: results

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- Compliance is decidable even for contracts with recursion
- The complexity is  $O(n^5)$ 
  - Straightforward algorithm is exponential
- The algorithm extends in a non trivial way the one for subtyping of recursive arrow and product types from Pierce

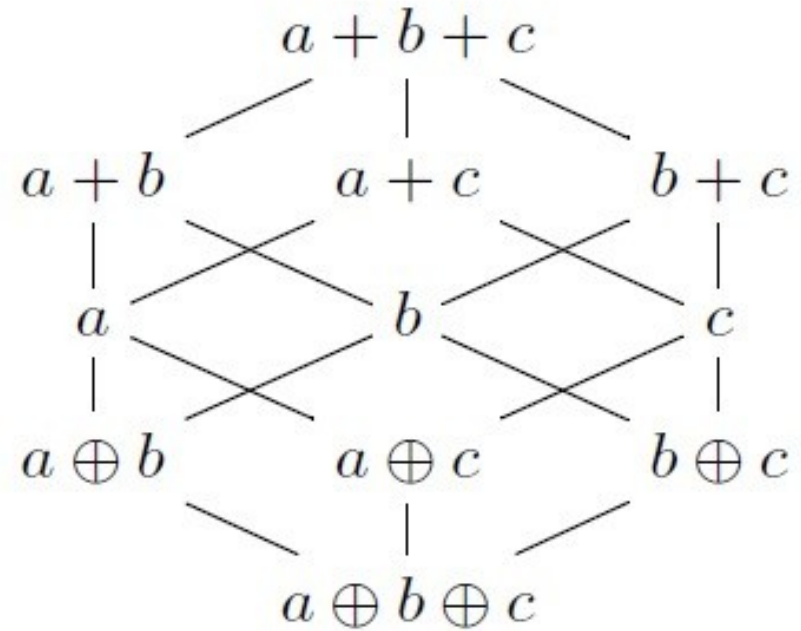
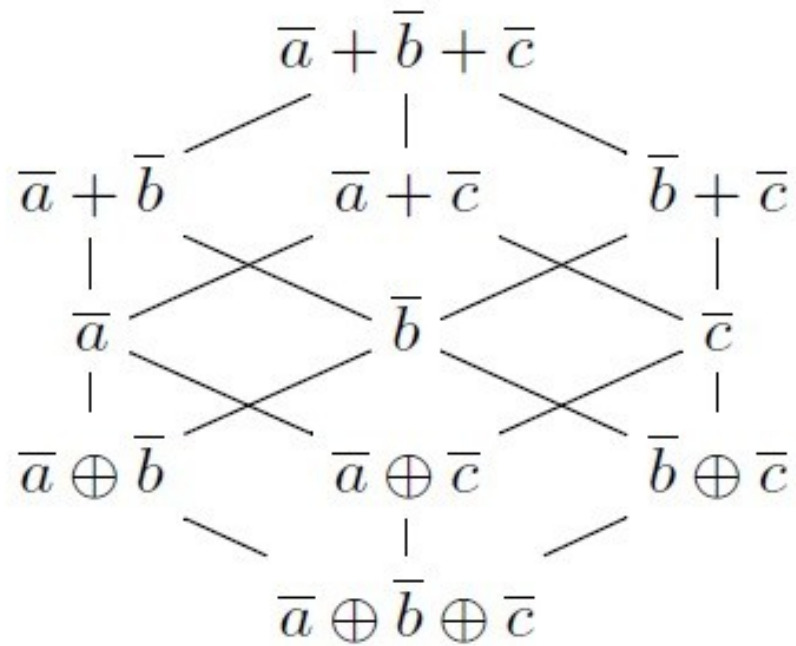
# Subcontract relation

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- Subcontract relation for servers:  
 $\rho \preceq_s \rho'$  iff for each client  $\sigma$ .  $\sigma \dashv\| \rho$  implies  $\sigma \dashv\| \rho'$ 
  - $\rho$  has more clients than  $\rho'$
- Subcontract relation for clients is dual:  
 $\sigma \preceq_c \sigma'$  iff for each server  $\rho$ .  $\sigma \dashv\| \rho$  implies  $\sigma' \dashv\| \rho$
- The two subcontract relations are partial orders
- The dual  $\bar{\sigma}$  of a client contract  $\sigma$  is the minimum server compliant with  $\sigma$

# Subcontract relation: example

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# Duality has a simple syntactic characterization

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$1$	$\leftrightarrow$	$1$
$\sum_{i \in I} a_i \cdot \sigma_i$	$\leftrightarrow$	$\bigoplus_{i \in I} \bar{a}_i \cdot \sigma_i$
$\sum_{i \in I} \bar{a}_i \cdot \sigma_i$	$\leftrightarrow$	$\bigoplus_{i \in I} a_i \cdot \sigma_i$
$X$	$\leftrightarrow$	$X$
$\text{rec } X \cdot \sigma$	$\leftrightarrow$	$\text{rec } X \cdot \sigma$



# Subcontract relation: results

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- Subcontract relation for servers and for clients are related:

$$\rho \preceq_s \rho' \text{ iff } \overline{\rho'} \preceq_c \overline{\rho}$$

- Subcontract relation and compliance are related:

$$\rho \preceq_s \rho' \text{ iff } \overline{\rho} \dashv\!\! \dashv \rho'$$

- Also the subcontract relation can be decided in  $O(n^5)$

# Retractable contracts vs reversible computation

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- Take retractable contracts without retraction
- Apply to it the technique to make a calculus reversible from Phillips and Ulidowski
- Retraction corresponds to a sequence of backward steps in the resulting reversible calculus
- Hence, retractable contracts are a form of reversible computation with internal/semantic control
- If you drop these forms of control then compliance becomes trivial

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# Summary

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- We presented a model of contracts with retractable choice
- Using retractable choice instead of normal choice ensures compliance with a larger set of partners
- Retractable contracts have most of the good properties of contracts:
  - decidability of compliance and subcontract relation
  - efficient decidability algorithm
  - easy syntactic characterization of duality

# Future work

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- Explore the notion of retractable contracts in multiparty sessions
- How can we extract a contract from a reversible application?
- Are there other meaningful ways to exploit contracts/behavioural types to control reversibility?



End of talk

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Thanks!

Questions?

# Most related work

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- Franco Barbanera, Mariangiola Dezani-Ciancaglini, Ugo de'Liguoro: Compliance for reversible client/server interactions. BEAT 2014  
also considered contracts with rollback

BEAT 2014

vs

PLACES 2015

- Free rollback

vs

rollback only when stuck

- Explicit checkpoint

vs

implicit checkpoint

- One checkpoint

vs

stack of checkpoints

- **Compliance harder**

vs

**compliance easier**