

QuickTime™ and a
decompressor
are needed to see this picture.

LISP: An Architectural Solution to Multi-homing, Traffic Engineering, and Internet Route Scaling

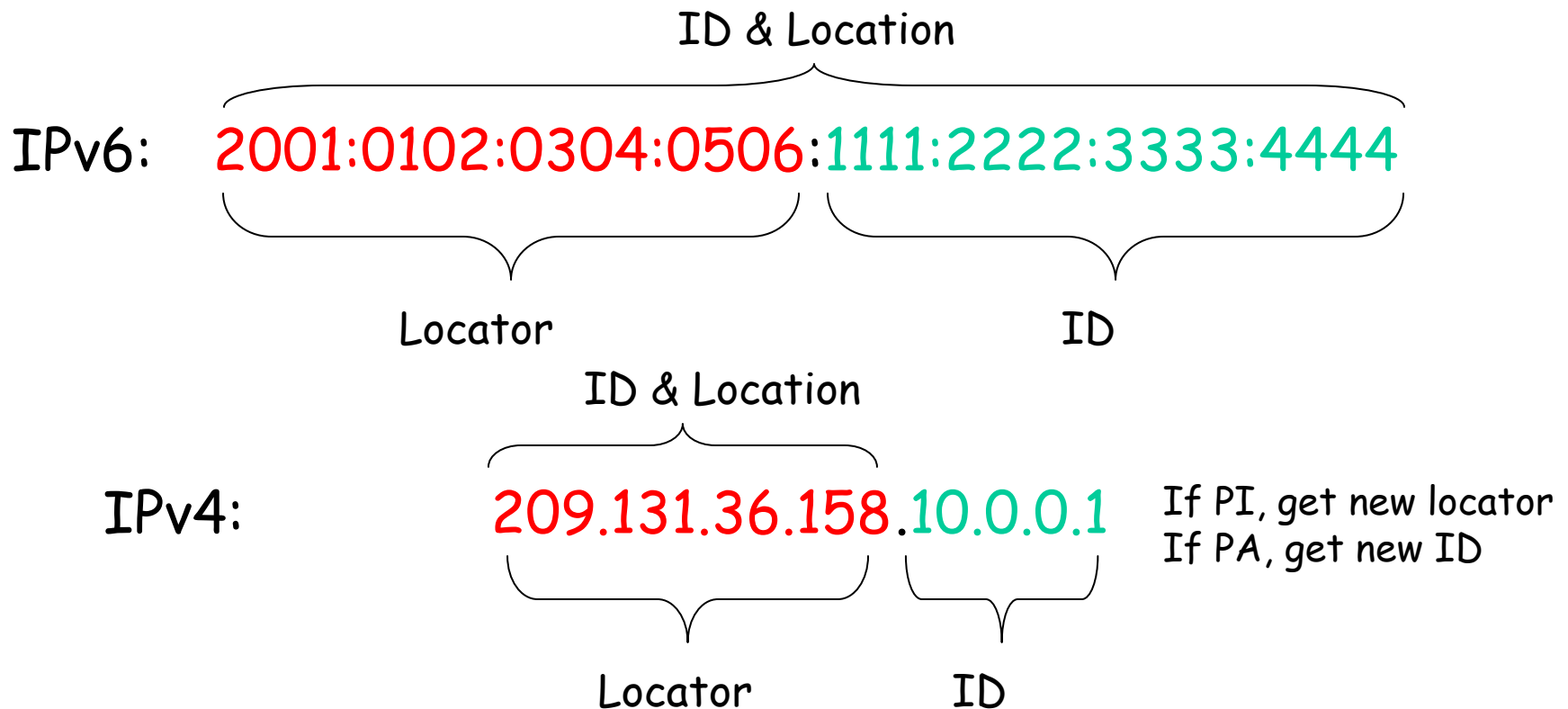
Dave Meyer & Dino Farinacci

LISP Designers:

*Dave Meyer, Vince Fuller, Darrel Lewis, Andrew Partan,
John Zwiebel, Scott Brim, Noel Chiappa & Dino Farinacci*

Separating (or adding) an Address

Changing the semantics of the IP address



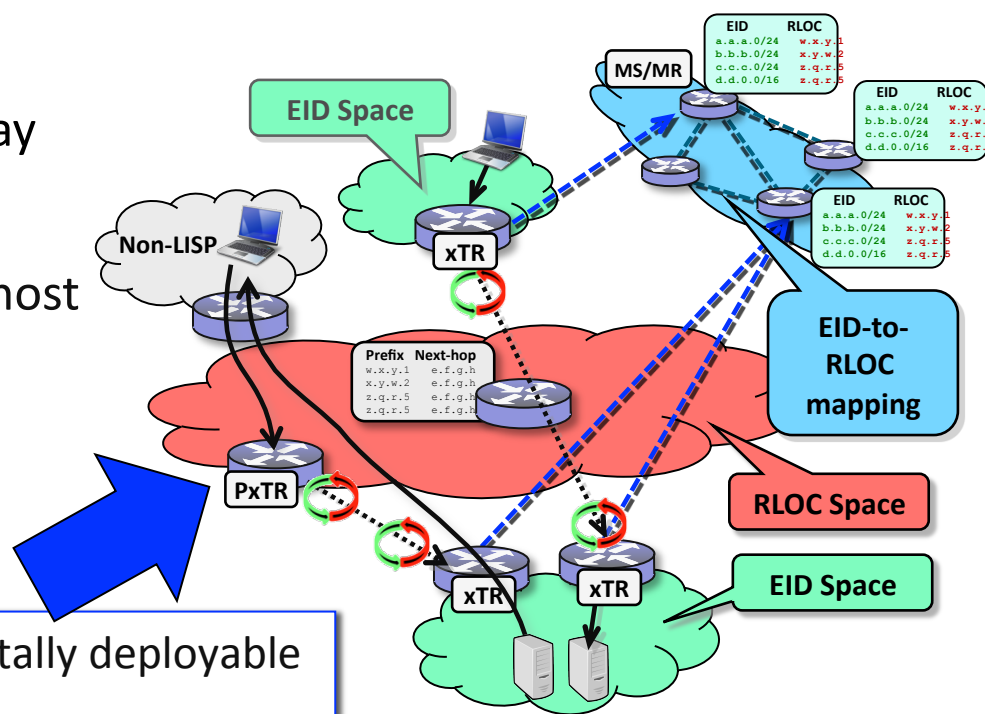
Why the Separation?

- *Level of Indirection* allows us to:
 - Keep either ID or Location fixed while changing the other
 - Create separate namespaces which can have different allocation properties
- By keeping IDs fixed
 - Assign fixed addresses that never change to hosts and routers at a site
- You can change Locators
 - Now the sites can change providers
 - Now the hosts can move

LISP Overview

LISP creates a **Level of indirection** with two namespaces: **EID** and **RLOC**

- **EID (Endpoint Identifier)** is the IP address of a host – just as it is today
- **RLOC (Routing Locator)** is the IP address of the LISP router for the host
- **EID-to-RLOC mapping** is the distributed architecture that maps **EIDs** to **RLOCs**

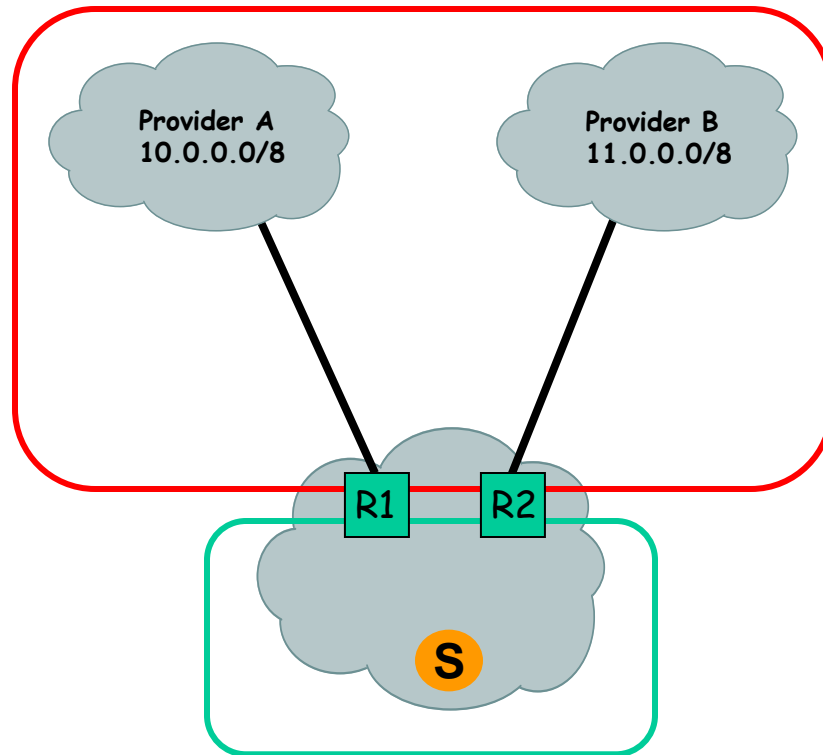


- Network-based solution
- No host changes
- Minimal configuration
- Incrementally deployable
- Support for mobility
- Address Family agnostic

Some Brief Definitions

- IDs or EIDs
 - End-site addresses for hosts and routers at the site
 - They go in DNS records
 - Generally not globally routed on underlying infrastructure
 - New namespace
- RLOCs or Locators
 - Infrastructure addresses for LISP routers and ISP routers
 - Hosts do not know about them
 - They are globally routed and aggregated along the Internet connectivity topology
 - Existing namespace

Multi-Level Addressing



RLOCs used in the core

EIDs are inside of sites

What is LISP?

- Locator/ID Separation Protocol
- Ground rules for LISP
 - Network-based solution
 - No changes to hosts whatsoever
 - No new addressing changes to site devices
 - Very few configuration file changes
 - Imperative to be incrementally deployable
 - Address family agnostic

LISP Overview

IP encapsulation scheme

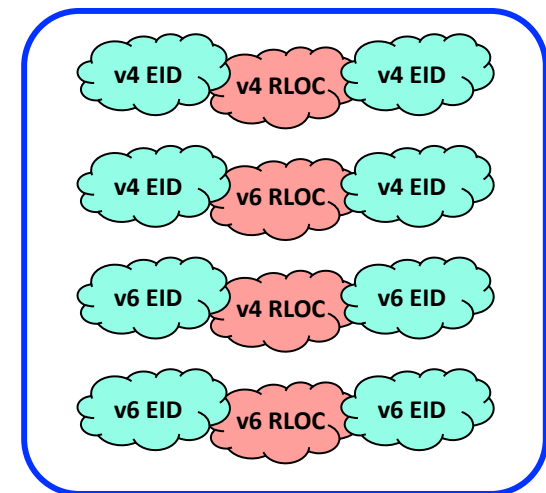
- Decouples host **IDENTITY** and **LOCATION**
- Dynamic **IDENTITY**-to-**LOCATION** mapping resolution
- Address Family agnostic day-one

Minimal Deployment Impact

- No changes to end systems or core
- Minimal changes to edge devices

Incrementally deployable

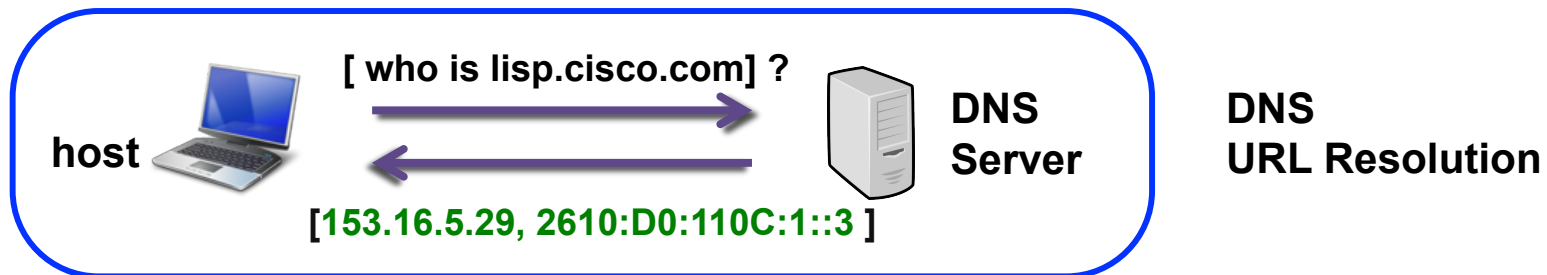
- LISP/LISP and non-LISP/LISP considered day-one



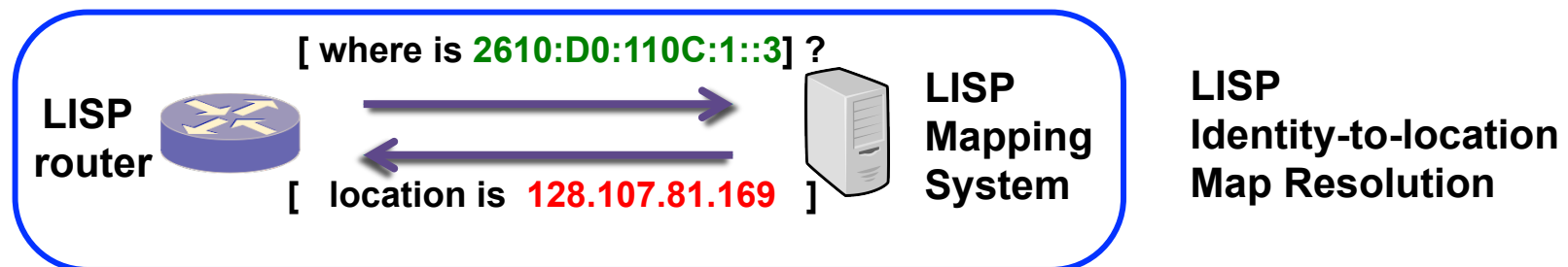
LISP Overview

LISP Map Lookup is analogous to a DNS lookup

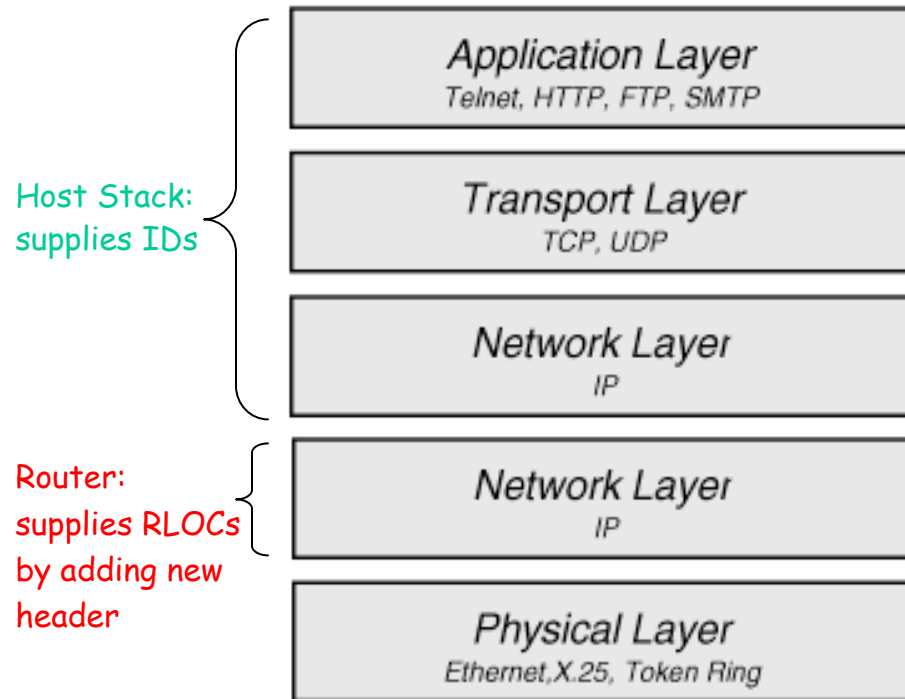
- DNS resolves IP addresses for URLs



- LISP resolves locators for queried identities

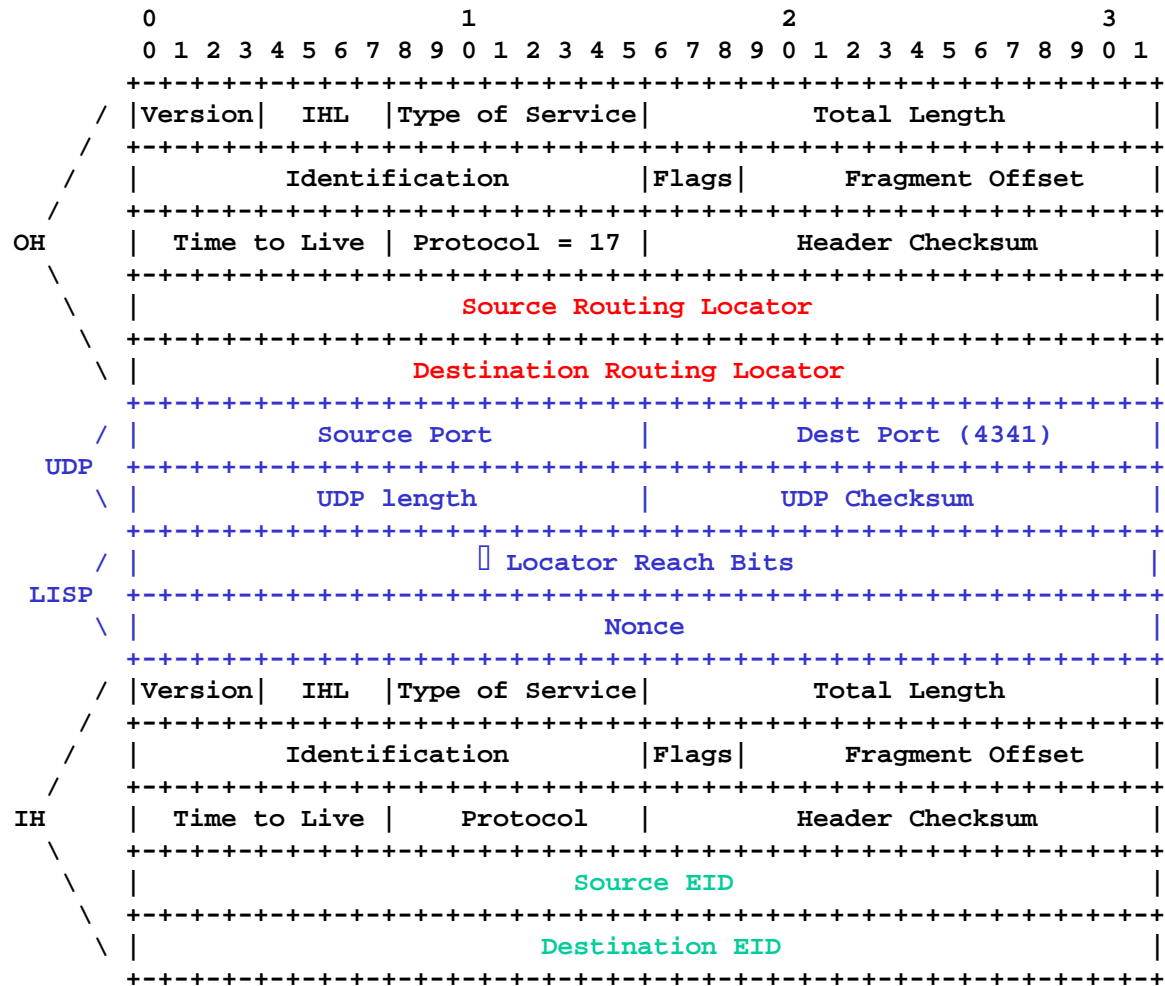


What is LISP?

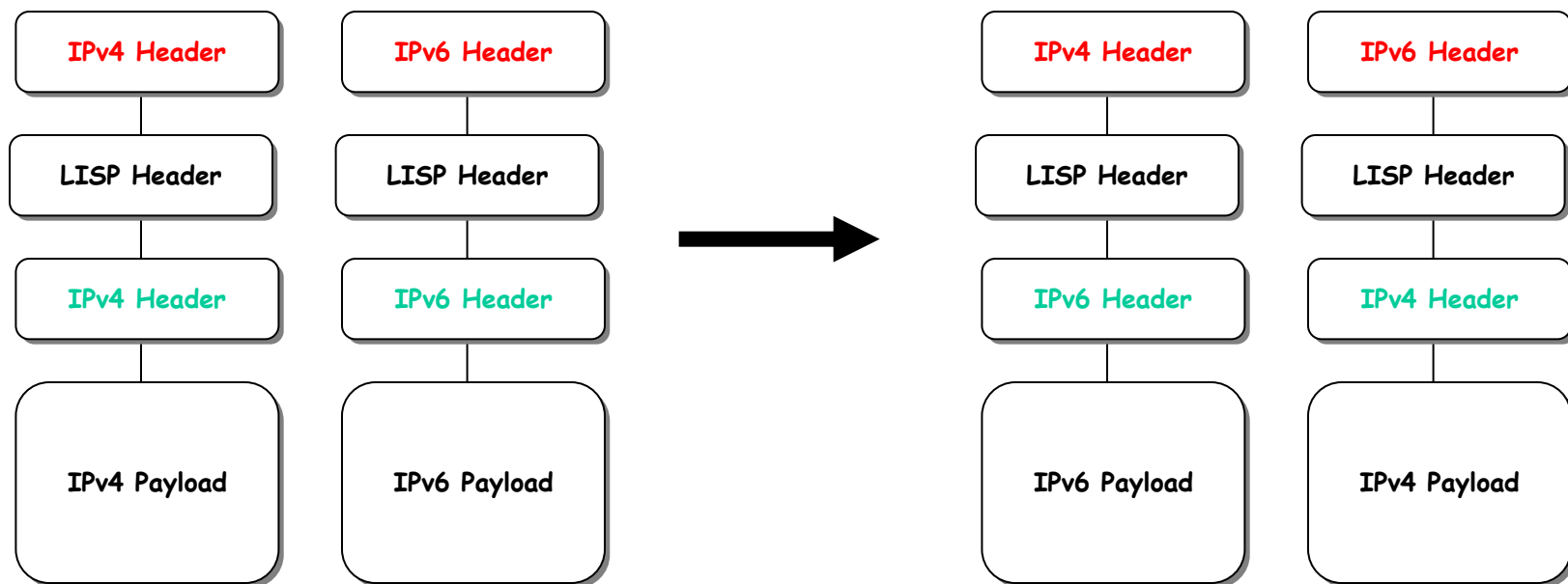


“Jack-Up” or “Map-n-Encap”

draft-farinacci-lisp-11.txt



LISP for IPv6 Transition



Legend:

EIDs -> Green

Locators -> Red

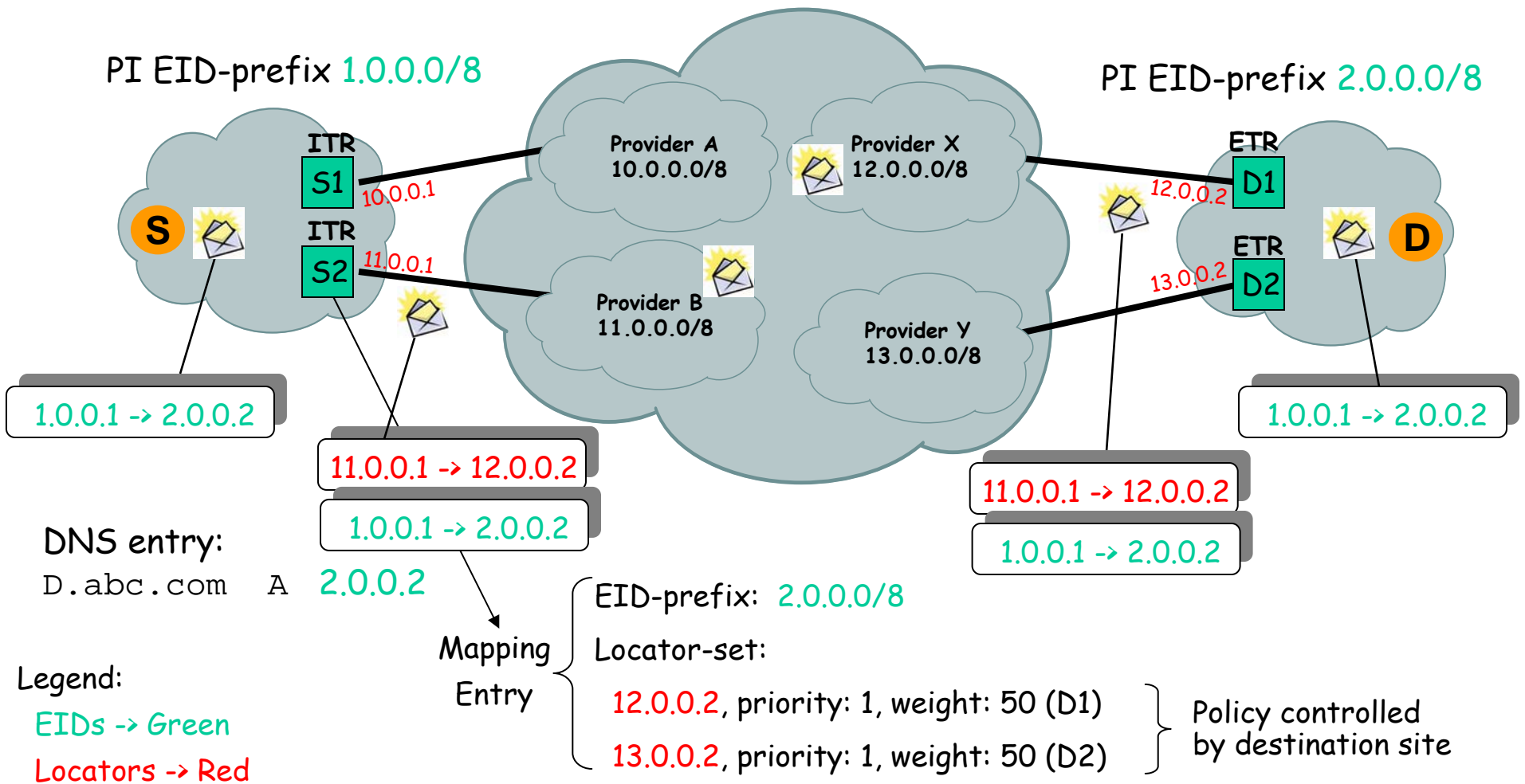
What is LISP?

- Data plane
 - Design for encapsulation and tunnel router placement
 - Design for locator reachability
 - Data-triggered mapping service
- Control plane
 - Design for a scalable mapping service
 - Examples are: CONS, NERD, ALT, EMACS

LISP Network Elements

- Ingress Tunnel Router (ITR)
 - Finds EID to RLOC mapping
 - Encapsulates to Locators at source site
- Egress Tunnel Router (ETR)
 - Owns EID to RLOC mapping
 - Decapsulates at destination site
- xTR
 - Term used when not referring to directionality
 - Basically a LISP router

Unicast Packet Forwarding



When the xTR has no Mapping

- Need a scalable EID to Locator mapping lookup mechanism
- Network based solutions
 - Have query/reply latency
 - Can have packet loss characteristics
 - Or, have a full table like BGP does
- How does one design a scalable Mapping Service?

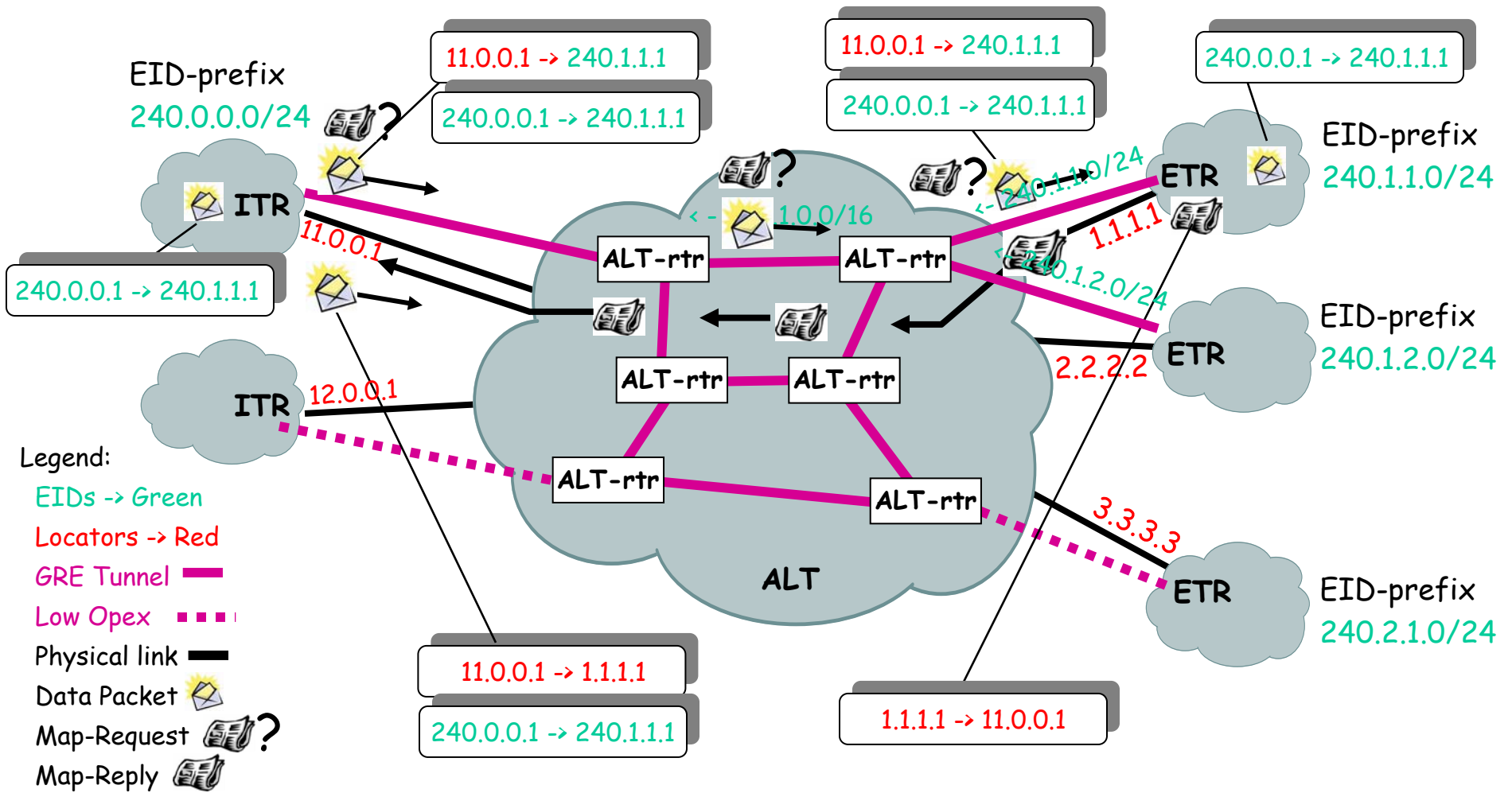
Mapping Database Designs

- You need a “map” before you can “encap”
- We have designed several mapping database protocols
 - CONS, NERD, EMACS, ALT
 - Tradeoff push versus pull benefit/cost
 - Needs to be scalable to 10^{10} entries
- ALT has the most promise
 - We are deploying ALT

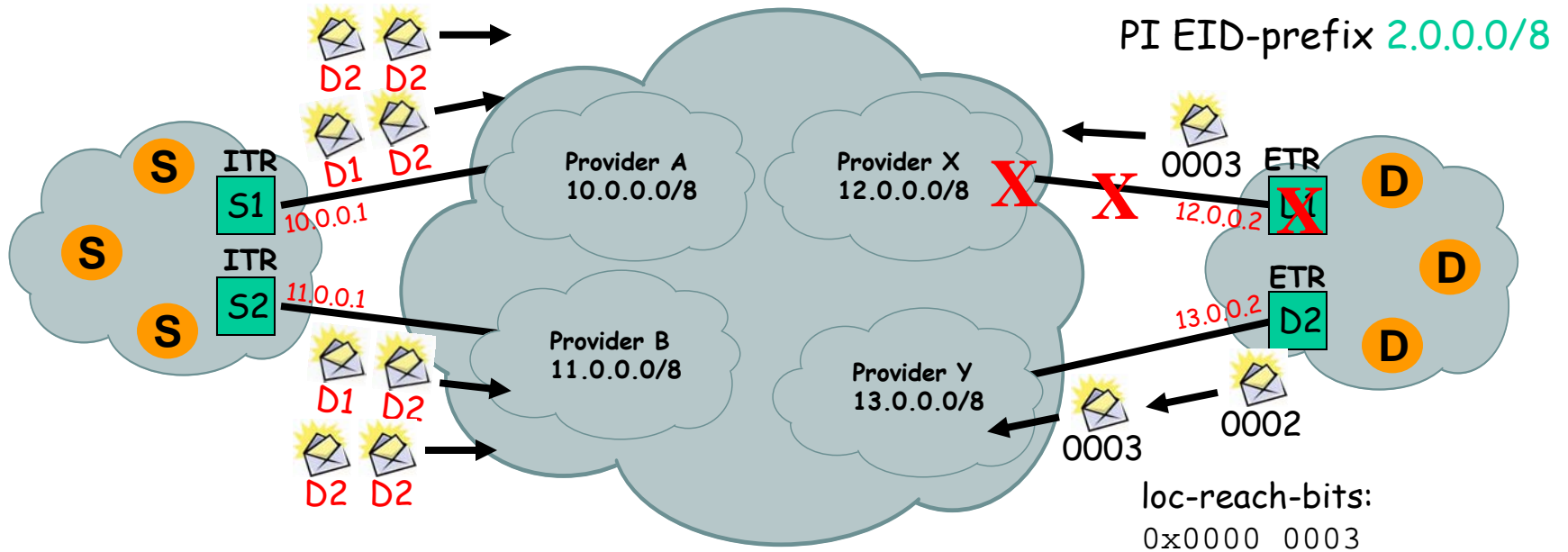
What is LISP+ALT?

- EID namespace is used at the site
- RLOC namespace is used in the Internet core
- Mappings need to be authoritative and reside at site ETRs
- Advertise EID-prefixes in BGP on an alternate topology of GRE tunnels
- ITRs get mappings by routing Map-Requests on ALT topology
- ETRs respond with Map-Replies

How LISP+ALT Works



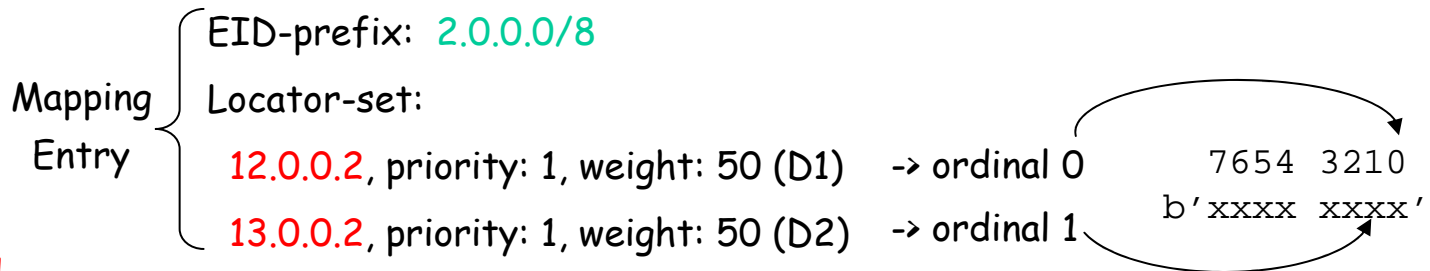
Locator Reachability



Legend:

EIDs -> Green

Locators -> Red



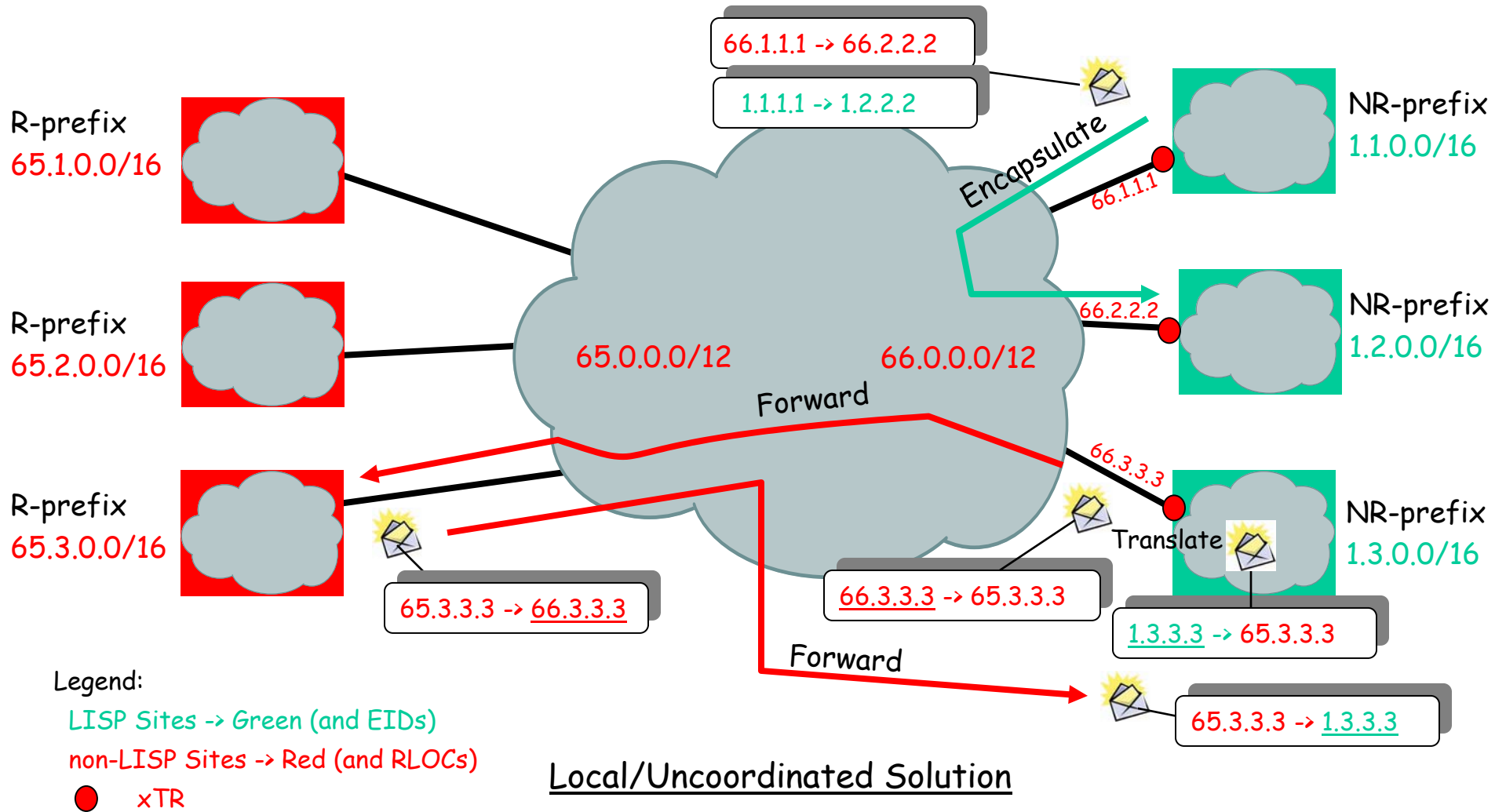
LISP Interworking

- LISP will not be widely deployed day-1
- Need a way for LISP-capable sites to communicate with rest of Internet
- Two basic Techniques
 - LISP Network Address Translators (LISP-NAT)
 - Proxy Tunnel Routers (PTRs)

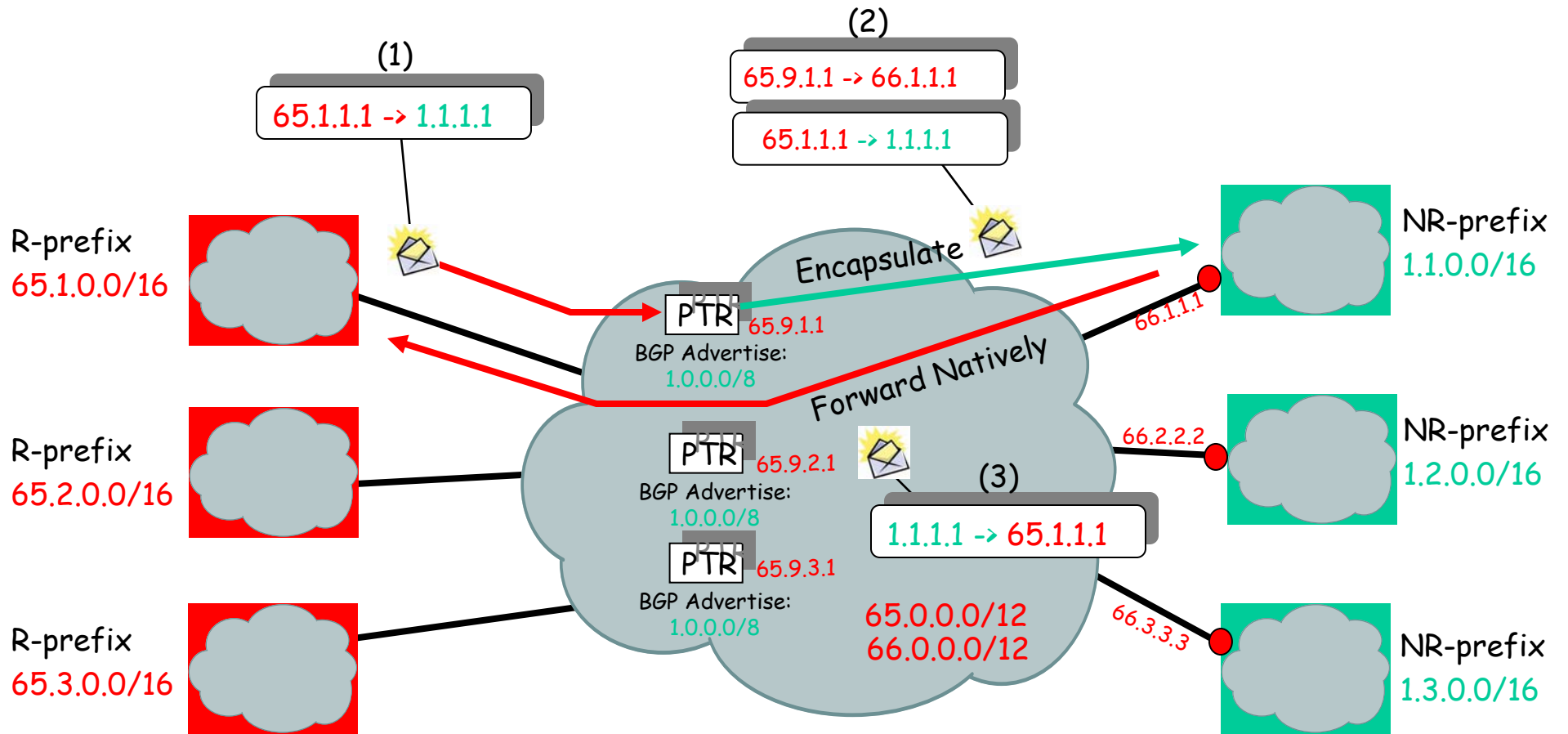
LISP Interworking

- These combinations must be supported
 - Non-LISP site to non-LISP site
 - Today's Internet
 - LISP site to LISP site
 - Encapsulation over IPv4 makes this work
 - IPv4-over-IPv4 or IPv6-over-IPv4
 - LISP-R site to non-LISP site
 - When LISP site has PI or PA routable addresses
 - LISP-NR site to non-LISP site
 - When LISP site has PI or PA non-routable addresses

Interworking using LISP-NAT



Interworking using PTRs



Legend:

LISP Sites -> Green (and EIDs)
 non-LISP Sites -> Red (and RLOCs)

● xTR

Slide 31

LISP Use Cases

VM-Mobility

Needs:

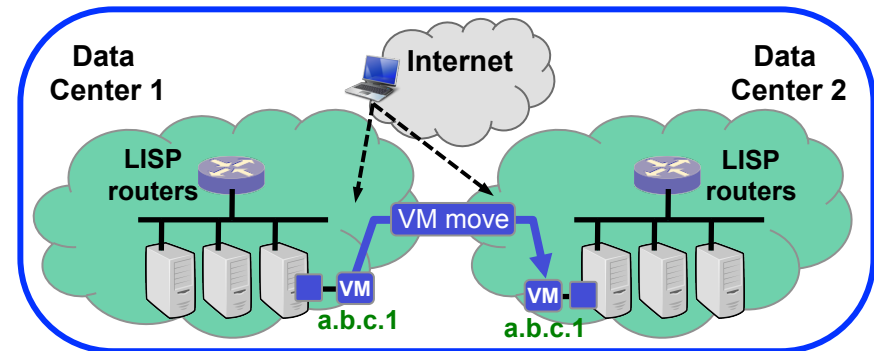
- VM-Mobility **across subnets**
- Move detection, dynamic EID-to-RLOC mappings, traffic redirection

LISP Solution:

- OTV + LISP to extend subnets
- LISP for VM-moves across subnets

Benefits:

- Integrated Mobility
- Direct Path (no triangulation)
- Connections maintained across moves
- No routing re-convergence
- No DNS updates required
- Global Scalability (cloud bursting)
- IPv4/IPv6 Support
- ARP elimination



Applicability:

- VM OS agnostic
- Services Creation (disaster recovery, cloud burst, etc.)

LISP Use Cases

LISP-MN

Needs:

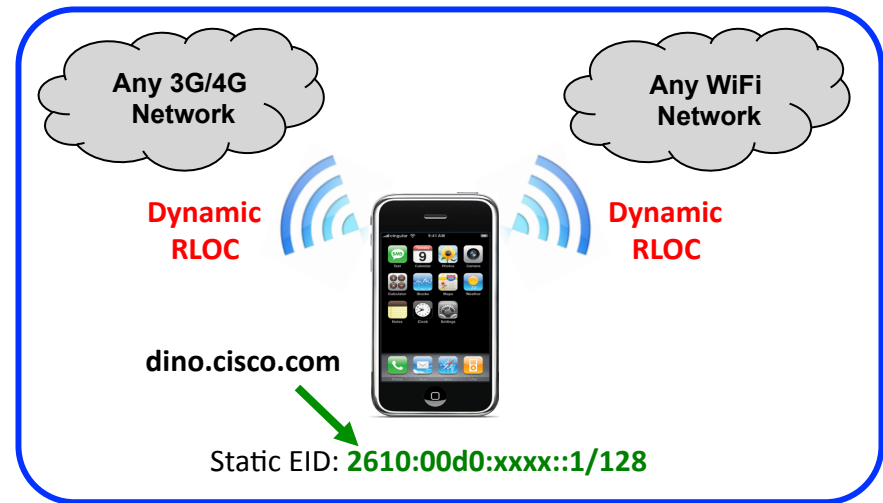
- Mobile devices roaming across any access media without connection reset
- Mobile device keeps the same IP address forever

LISP Solution:

- LISP level or indirection separates endpoints and locators
- Network-based; no host changes, minimal network changes
- Scalable, host-level registration (10^{10})

Benefits:

- MNs can roam and stay connected
- MNs can be servers
- MNs roam without DNS changes
- MNs can use multiple interfaces
- Packets have “stretch-1” reducing latency



Applicability:

- IPv4 and IPv6
- Android and Linux
- Open

Prototype Implementation

- cisco has a LISP prototype implementation
- Supports:
 - `draft-farinacci-lisp-11.txt`
 - `draft-fuller-lisp-alt-03.txt`
 - `draft-lewis-lisp-interworking-02.txt`
- Software switching only
- Supports LISP for both IPv4 and IPv6
 - ITR, ETR, and PTR
 - LISP-NAT for IPv4 only

Internet Drafts

draft-farinacci-lisp-11.txt
draft-farinacci-lisp-multicast-01.txt
draft-fuller-lisp-alt-03.txt
draft-lewis-lisp-interworking-02.txt
draft-meyer-lisp-eid-block-01.txt
draft-meyer-loc-id-implications-01.txt

draft-mathy-lisp-dht-00.txt
draft-iannone-openlisp-implementation-02.txt
draft-brim-lisp-analysis-00.txt

draft-meyer-lisp-cons-04.txt
draft-lear-lisp-nerd-04.txt
draft-curran-lisp-emacs-00.txt

References

- **Public mailing list:**
`lisp@ietf.org`
- **More info at:**
`http://www.lisp4.net`
`http://www.lisp6.net`