

What is it about?

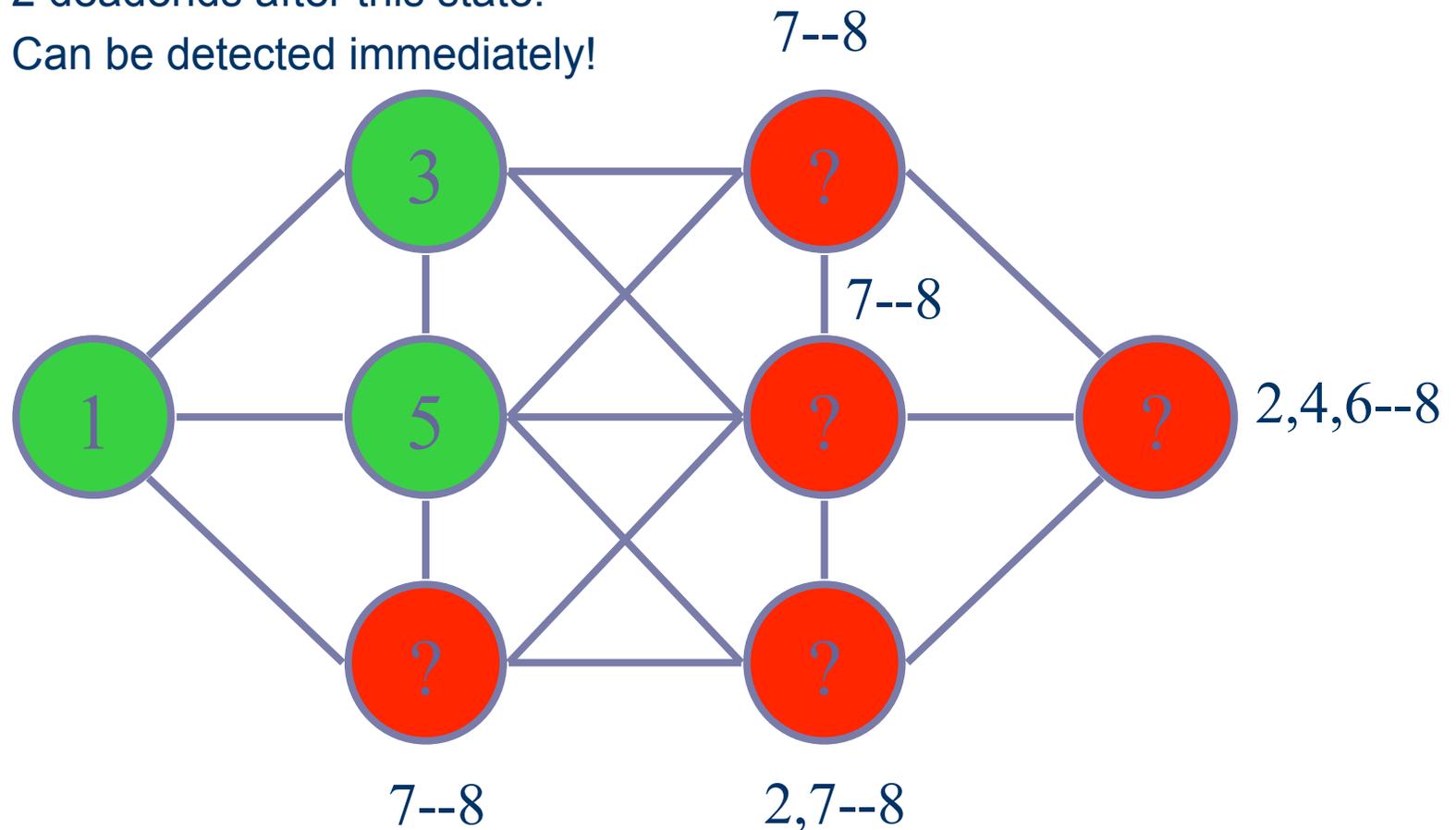
- Theory
 - The core of solving constraint problems using Constraint Programming (CP), with emphasis on:
 - Modeling
 - Solving:
 - Local consistency and propagation;
 - Backtracking search + heuristics.
 - Brief overview of advanced topics.
- Practice
 - Programming examples.
 - Project work.

Efficiency

- Bad choice of nodes, bad assignment of values
=> **Good heuristic choice** is very important!
- Good heuristics are always possible?
 - No!
- What can we do then?
 - **Good consistency and propagation!**
- Didn't we do that already?
 - Not as strong as we could!

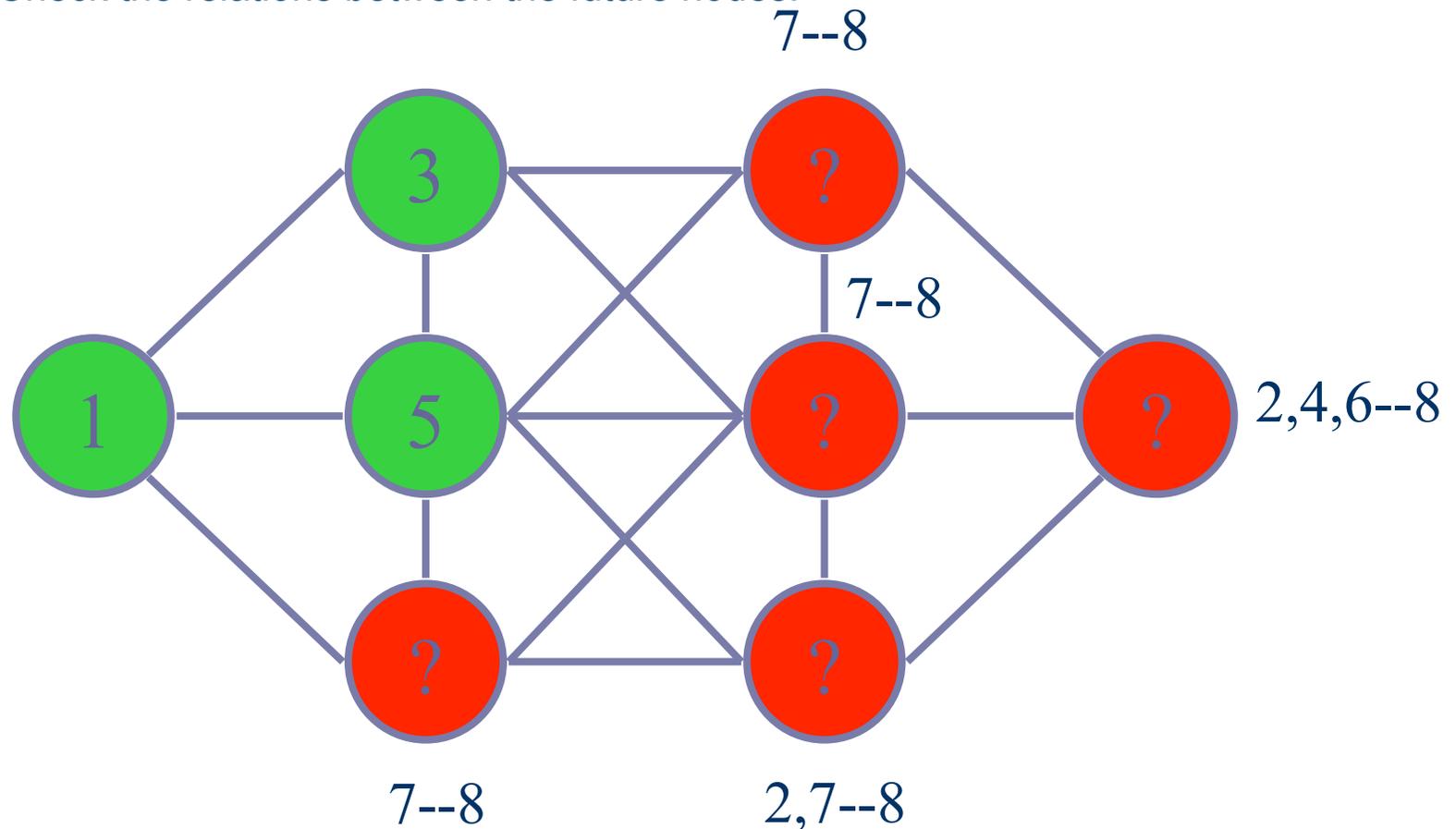
A State During Search

- 2 deadends after this state.
- Can be detected immediately!



A State During Search

- Check the relations between the future nodes.

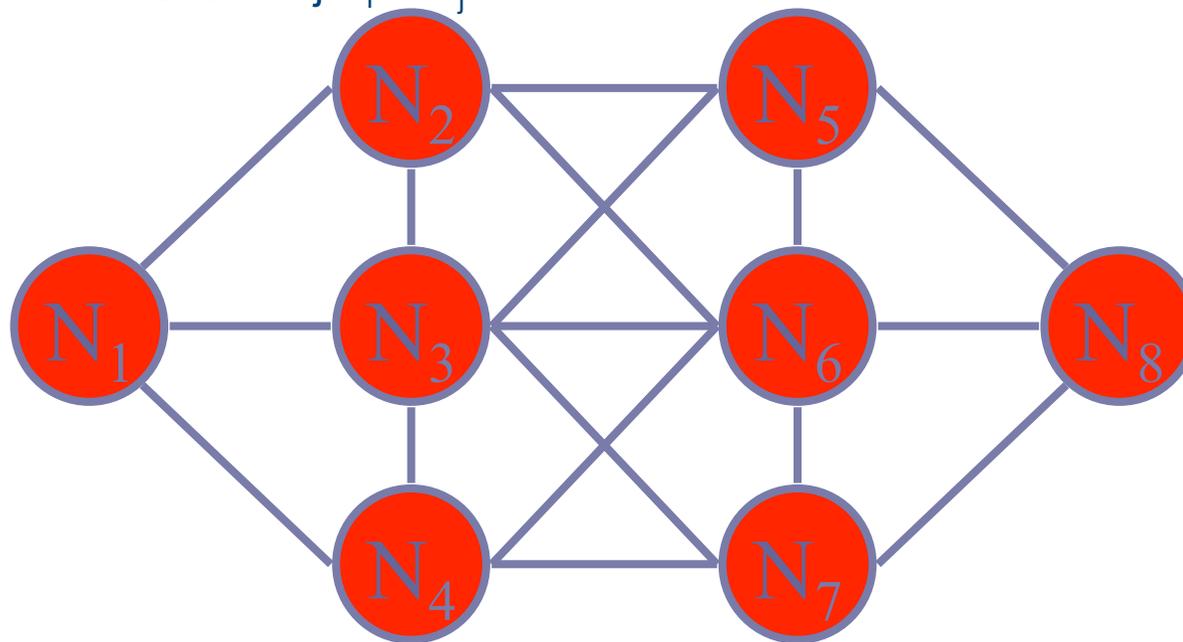


Efficiency

- Bad choice of nodes, bad assignment of values
 - => **Good heuristic choice** is very important!
- Good heuristics are always possible?
 - No!
- What can we do then?
 - **Good consistency and propagation!**
- Didn't we do that already?
 - Not as strong as we could!
- Is that all?
 - Nope!
 - **Good modeling** can result in even stronger consistency & propagation

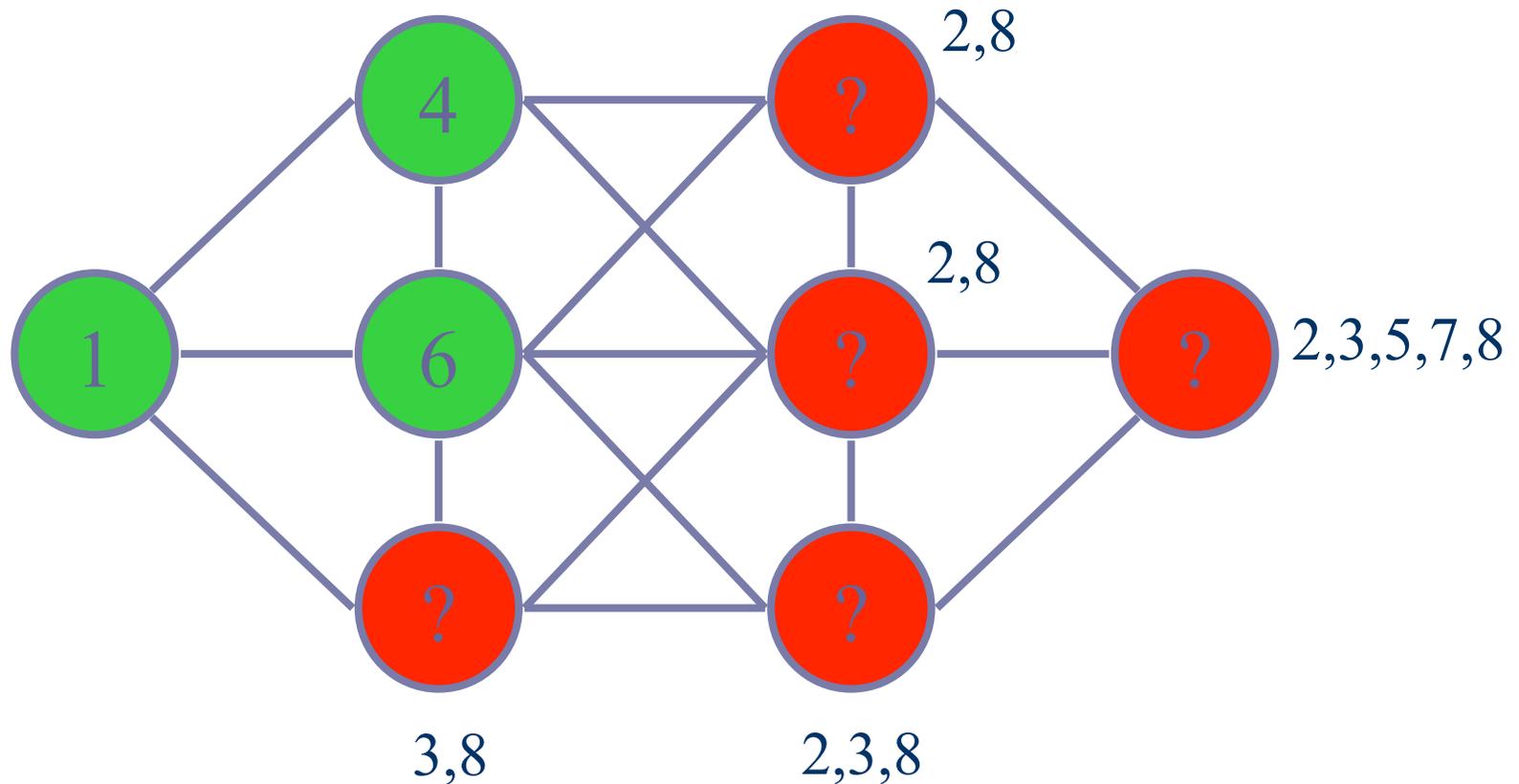
Modeling

- Unknowns: $N_1..N_8$ represent the nodes
- Values that $N_1..N_8$ can take: $\{1,2,3,4,5,6,7,8\}$
- Relations: for all $i < j$ s.t. N_i and N_j are adjacent. $|N_i - N_j| > 1$
for all $i < j$ $N_i \neq N_j$



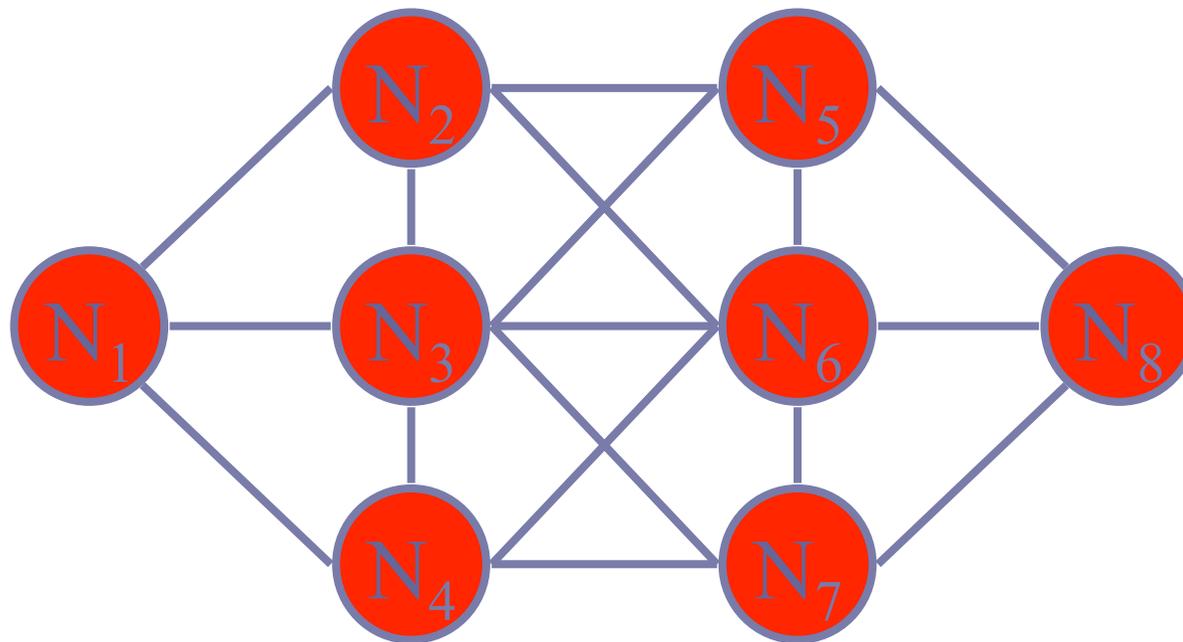
Another State

- Cannot detect the inconsistency of $N_3 = 6$.
 - Future nodes are fine wrt the relations.



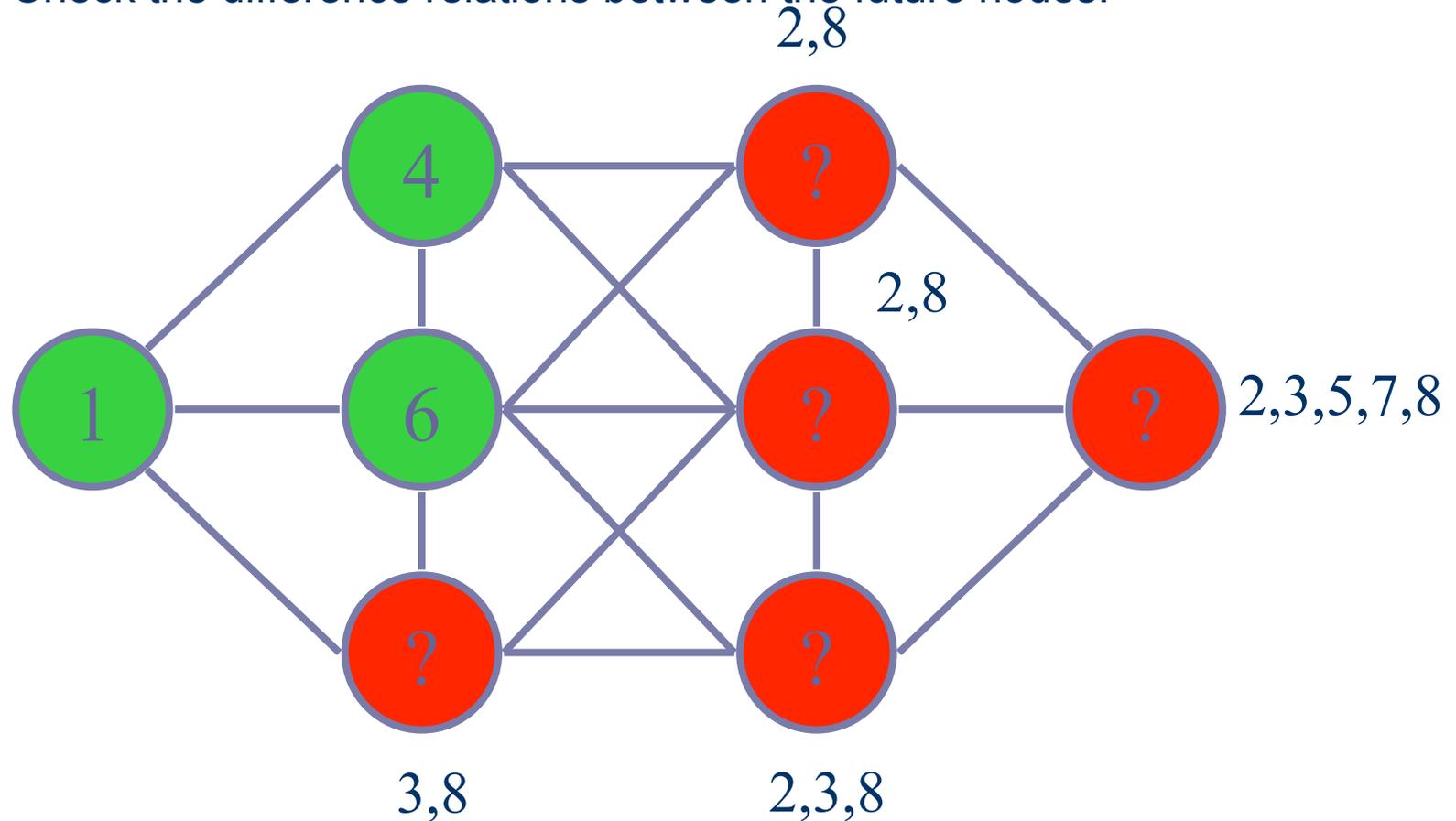
Stronger Model

- Relations:
 - for all $i < j$ s.t. N_i and N_j are adjacent $|N_i - N_j| > 1$
 - for all $i < j$ $N_i \neq N_j$ alldifferent($[N_1, N_2, N_3, N_4, N_5, N_6, N_7, N_8]$)



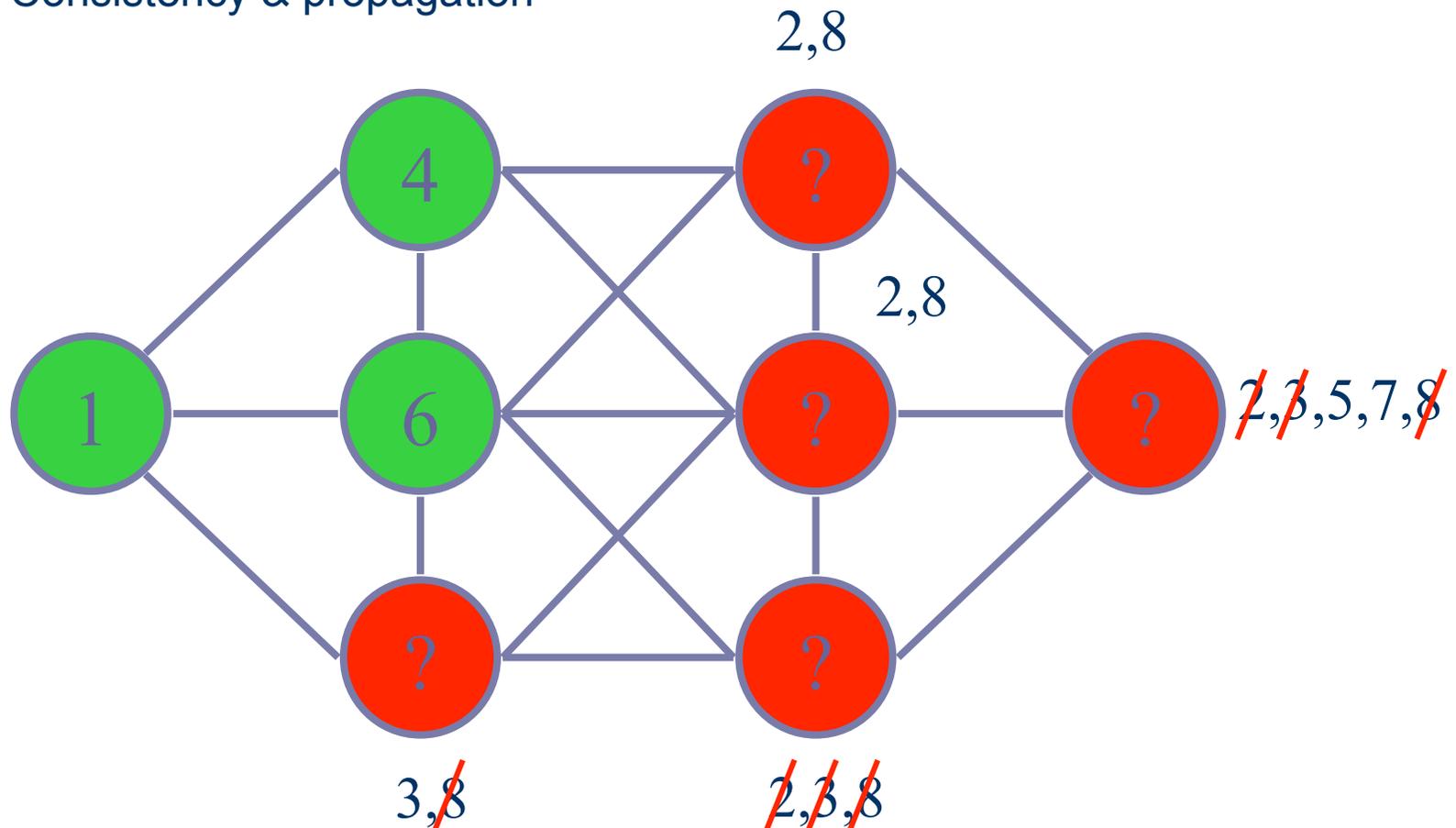
Another State

- Check the difference relations between the future nodes.



Another State

- Consistency & propagation



PART I: Modeling



Modeling

- Representation of a decision problem as a CSP.
- Formalize:
 - unknowns of the decision → variables
 - possible values for unknowns → domains
 - relations between the unknowns → constraints

CSPs: More formally

- A CSP is a triple $\langle X, D, C \rangle$ where:
 - X is a set of decision variables $\{X_1, \dots, X_n\}$;
 - D is a set of domains $\{D_1, \dots, D_n\}$ for X :
 - D_i is a set of possible values for X_i ;
 - usually assume finite domain.
 - C is a set of constraints $\{C_1, \dots, C_m\}$:
 - C_i is a relation over X_j, \dots, X_k , giving the set of combination of allowed values;
 - $C_i \subseteq D(X_j) \times \dots \times D(X_k)$
- A **solution** to a CSP is an assignment of values to the variables which satisfies all the constraints simultaneously.

Constraint Optimization Problems

- CSP enhanced with an optimization criterion, e.g.:
 - minimum cost;
 - shortest distance;
 - fastest route;
 - maximum profit.
- Formalization of the optimization criterion as an objective function.

CSPs: A simple example

- **Variables**

$$X = \{X_1, X_2, X_3\}$$

- **Domains**

$$D(X_1) = \{1,2\}, D(X_2) = \{0,1,2,3\}, D(X_3) = \{2,3\}$$

- **Constraints**

$$X_1 > X_2 \text{ and } X_1 + X_2 = X_3 \text{ and } X_1 \neq X_2 \neq X_3 \neq X_1$$

- **Solution**

$$X_1 = 2, X_2 = 1, X_3 = 3$$

alldifferent([X₁, X₂, X₃])



Sudoku

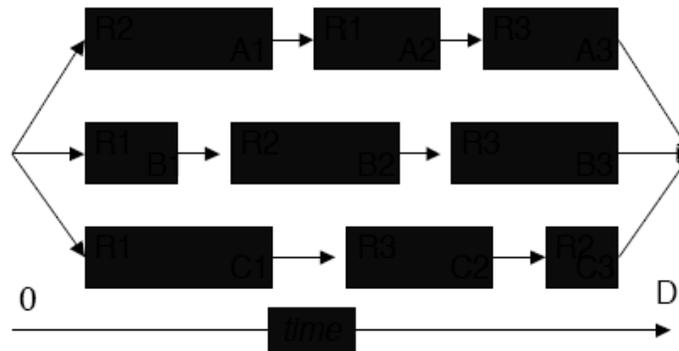
| | | | | | | | | |
|---|---|---|---|--|---|---|---|---|
| | 6 | | 1 | | 4 | | 5 | |
| | | 8 | 3 | | 5 | 6 | | |
| 2 | | | | | | | | 1 |
| 8 | | | 4 | | 7 | | | 6 |
| | | 6 | | | | 3 | | |
| 7 | | | 9 | | 1 | | | 4 |
| 5 | | | | | | | | 2 |
| | | 7 | 2 | | 6 | 9 | | |
| | 4 | | 5 | | 8 | | 7 | |

Sudoku

| | | | | | | |
|----------|---|---|---|---|--|----------|
| X_{11} | 6 | 1 | 4 | 5 | | X_{91} |
| . | 8 | 3 | 5 | 6 | | . |
| 2 | | | | | | 1 |
| . | | | | | | . |
| 8 | | 4 | 7 | | | 6 |
| . | 6 | | | 3 | | . |
| . | | | | | | . |
| 7 | | 9 | 1 | | | 4 |
| . | | | | | | . |
| 5 | | | | | | 2 |
| . | 7 | 2 | 6 | 9 | | . |
| X_{19} | 4 | 5 | 8 | 7 | | X_{99} |

- A simple CSP
 - 9x9 variables (X_{ij}) for each cell with domains $\{1, \dots, 9\}$
 - Not-equals constraints on the rows, columns, and 3x3 boxes. E.g.,
 - $\text{alldifferent}([X_{11}, X_{21}, X_{31}, \dots, X_{91}])$
 - $\text{alldifferent}([X_{11}, X_{12}, X_{13}, \dots, X_{19}])$
 - $\text{alldifferent}([X_{11}, X_{21}, X_{31}, X_{12}, X_{22}, X_{32}, X_{13}, X_{23}, X_{33}])$

Task Scheduling



- Schedule tasks on a machine, in time D , by obeying the temporal and precedence constraints:
 - each task i has a specific fixed processing time p_i ;
 - each task can be started after its release date r_i , and must be completed before its deadline d_i ;
 - tasks cannot overlap in time;
 - precedence relations must be respected.

Task Scheduling

- **Variables and Domains**
 - $Start_i$, representing the starting time of each task i , taking a value from $\{1, 2, \dots, D\}$.
 - This immediately ensures that each task starts at exactly one time point.
- **Constraints**
 - Respect release date and deadline:
 - for all i $r_i \leq Start_i \leq d_i - p_i$
 - Tasks cannot overlap in time:
 - for all $i < j$ $(Start_i + p_i < Start_j)$ OR $(Start_j + p_j < Start_i)$
 - Precedences:
 - $Start_i < Start_j$ for every pair of task i with precedence over task j

Variables and Domains

- Variable domains include the classical:
 - binary, integer, continuous.
- In addition, variables may take a value from *any* finite set.
 - e.g., X in $\{a,b,c,d,e\}$.
- There exist special “structured” variable types.
 - Set variables (take a set of elements as value).
 - Activities or interval variables (for scheduling applications).

Properties of Constraints

- Declarative (invariant) relations among objects.
 - $X > Y$
- The order of imposition does not matter.
 - $X + Y \leq Z, X + Y \geq Z$
- Non-directional.
 - A constraint between X and Y can be used to infer information on Y given information on X and vice versa.
- Rarely independent.
 - Shared variables as communication mechanism.
- Incremental operational behavior.
 - When new information available, the computation does not necessarily start from scratch.

Constraints – Examples

- Algebraic expressions.
 - $X_1 > X_2$
 - $X_1 + X_2 = X_3$
 - $X^3(Y^2 - Z) \geq 25 + X^2 \cdot \max(X, Y, Z)$
- Extensional constraints ('table' constraints).
 - $(X, Y, Z) \in \{(a, a, a), (b, b, b), (c, c, c)\}$
- Variables as subscripts ('element' constraints).
 - $Y = \text{cost}[X]$ (here Y and X are variables, 'cost' is an array of parameters)

Constraints – Examples

- Reasoning with meta-constraints.
 - $\sum_i (X_i > t_i) \leq 5$
- Logical relations in which (meta-)constraints can be mixed.
 - $((X < Y) \text{ OR } (Y < Z)) \Rightarrow (C = \min(X, Y))$
- Global constraints.
 - $\text{alldifferent}([X_1, X_2, X_3])$ instead of $X_1 \neq X_2, X_1 \neq X_3, X_2 \neq X_3$

Modeling is Critical!

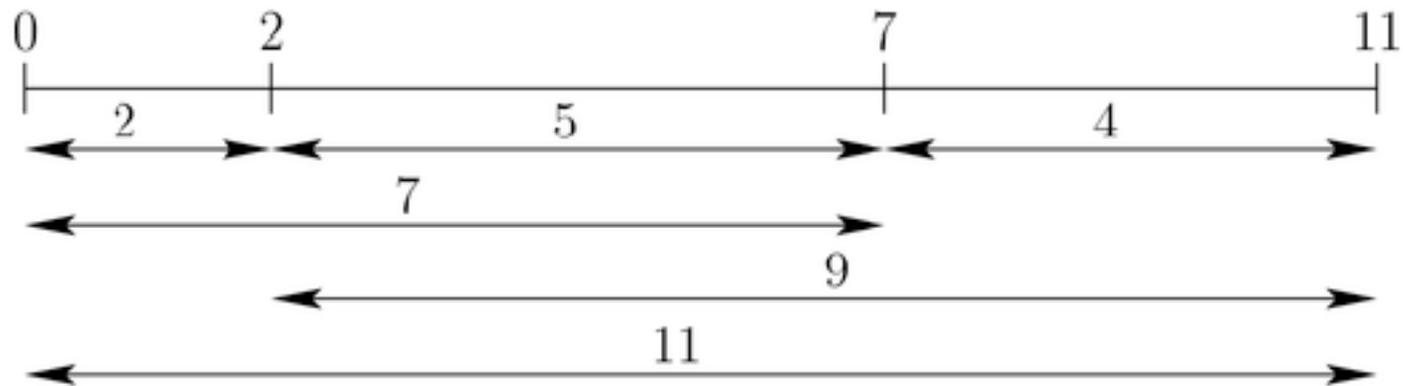
- Choice of variables defines search space.
 - $D(X_1) \times D(X_2) \times \dots \times D(X_n)$
- Choice of constraints defines:
 - how domains search space can be reduced;
 - how search can be guided.
- Need to go beyond the declarative specification!

Modeling is Critical

- Given the human understanding of a problem, we need to answer questions like:
 - which variables shall I choose?
 - which constraints shall I enforce?
 - can I exploit any global constraints?
 - so I need any auxiliary variables?
 - are some constraints redundant, therefore can be avoided?
 - are there any implied constraints?
 - can symmetry be eliminated?
 - are there any dual viewpoints?
 - among alternative models, which one shall I prefer?

Golomb Ruler

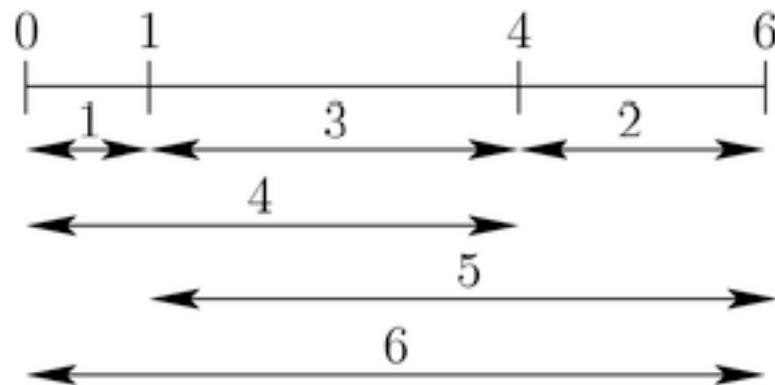
- Place m marks on a ruler such that:
 - distance between each pair of marks is different;
 - the length of the ruler is minimum.
- Applications in radio astronomy and information theory.
- Difficult to solve! Largest known ruler is of order 27.



A non optimal Golomb ruler of order 4.

Golomb Ruler

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An optimal Golomb ruler of order 4.

Naive Model

- **Variables and Domains**
 - $[X_1, X_2, \dots, X_m]$
 - X_i , representing the position of the i^{th} mark, taking a value from $\{0, 1, \dots, 2^{(m-1)}\}$

Naive Model

- **Variables and Domains**
 - $[X_1, X_2, \dots, X_m]$
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Naive Model

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 - $[X_1, X_2, \dots, X_m]$
 - X_i , representing the position of the i^{th} mark, taking a value from $\{0, 1, \dots, 2^{(m-1)}\}$
- **Constraints**
 - for all $i_1 < j_1, i_2 < j_2, i_1 \neq i_2$ or $j_1 \neq j_2$ $|X_{i_1} - X_{j_1}| \neq |X_{i_2} - X_{j_2}|$
- **Objective:** minimize $(\max([X_1, X_2, \dots, X_m]))$

Naive Model

- **Variables and Domains**
 - $[X_1, X_2, \dots, X_m]$
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- **Objective:** minimize ($\max([X_1, X_2, \dots, X_m])$)
- **Problematic model.**
 - $O(n^4)$ quaternary constraints.
 - Loose reduction in domains.

Better Model

- **Auxiliary Variables**

- Variables introduced into a model, because either:
 - it is difficult/impossible to express some constraints on the main variables;
 - or some constraints on the main variables do not lead to significant domain reductions.
- for all $i < j$ D_{ij} , representing the distance between i^{th} and the j^{th} marks.

- **Constraints**

- for all $i < j$ $D_{ij} = |X_i - X_j|$
- alldifferent($[D_{12}, D_{13}, \dots, D_{(m-1)m}]$)

Better Model

- **Constraints**

- for all $i < j$ $D_{ij} = |X_i - X_j|$
- alldifferent($[D_{12}, D_{13}, \dots, D_{(m-1)m}]$)
- alldifferent($[X_1, X_2, \dots, X_m]$)

- **Improvements:**

- $O(n^2)$ ternary constraints.
- **Implied constraint.**
 - Propagation is in general incomplete: inconsistent values are left in the domains.
 - Constraint implied by the constraints defining the problem which cannot be deduced by the incomplete solver.
 - No change in set of solutions, great reductions in search space!
- Tighter constraints and denser constraint graph.

Improved Model

- **Symmetry**
 - Positions can be permuted.
 - $X_1 = 0, X_2 = 1, X_3 = 4, X_4 = 6$
 - $X_1 = 0, X_2 = 1, X_3 = 6, X_4 = 4$
 - $X_1 = 0, X_2 = 4, X_3 = 1, X_4 = 6$
 - $X_1 = 0, X_2 = 4, X_3 = 6, X_4 = 1$
 - $X_1 = 0, X_2 = 6, X_3 = 1, X_4 = 4$
 - $X_1 = 0, X_2 = 6, X_3 = 4, X_4 = 1$
 - ...
 - $n!$ permutations

Improved Model

- **Symmetry**
 - Positions can be permuted.
 - A ruler can be reversed.

 - $X_1 = 0, X_2 = 1, X_3 = 4, X_2 = 6$
 - $X_1 = 0, X_2 = 2, X_3 = 5, X_2 = 6$

Improved Model

- **Symmetry**

- Positions can be permuted.
- A ruler can be reversed.

} Many symmetrically equivalent search states!

Improved Model

- **Symmetry**

- Positions can be permuted.
- A ruler can be reversed.
- Symmetry breaking constraints.

} Many symmetrically equivalent search states!

- Not implied by the constraints defining the problem.
- Reduce the set of solutions and search space!

- **Symmetry Breaking Constraints**

- $X_1 < X_2 < \dots < X_m$
- $X_1 = 0$
- $D_{12} < D_{(m-1)m}$

- **New objective**

- minimize (X_m)

Benefits of Remodeling

- The m marks must all be different in $1..d$.
 - Search space = number of m -permutations: $d! / (d-m)!$
- But they can also be totally ordered.
 - Search space = number of m -combinations: $d! / m!(d-m)!$
 - $m!$ times smaller!
- Instead of 0 being any mark on the ruler we can constrain it to be at position 1.
 - Reduction of d .
- Then we just have to minimise the position of the last mark.

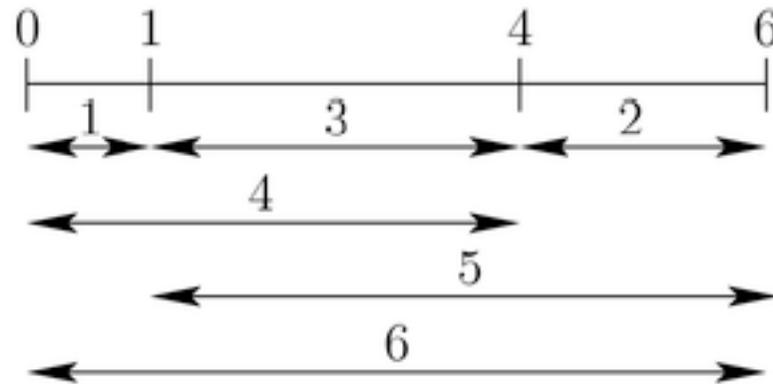
Improved Model

- Symmetry Breaking Constraints Enable Constraint Simplification
 - $X_1 < X_2 < \dots < X_m$
 - alldifferent on X_i becomes **redundant**.
 - for all $i < j$ $D_{ij} = |X_i - X_j| \rightarrow$ for all $i < j$ $D_{ij} = X_j - X_i$

Improved Model

- Symmetry Breaking Constraints Enable Additional Implied Constraints

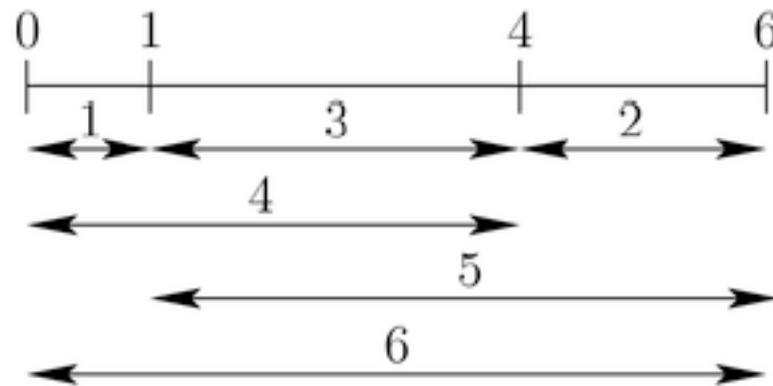
– for all $i < j < k$ $D_{ij} < D_{ik}$ and $D_{jk} < D_{ik}$
 $D_{ik} = D_{ij} + D_{jk}$



An optimal Golomb ruler of order 4.

Improved Model

- **Deducing information from Golomb Rulers of smaller order**
 - If you consider any k consecutive marks of a Golomb Ruler of order $n > k$, they form a Golomb Ruler of order k .



An optimal Golomb ruler of order 4.

Improved Model

- **Deducing information from Golomb Rulers of smaller order**
 - If you consider any k consecutive marks of a Golomb Ruler of order $n > k$, they form a Golomb Ruler of order k .
 - Therefore they must span over a distance at least as long as the optimal size of Rulers of order k .
 - for all $i < j$ $D_{ij} \geq \text{Ruler}(j-i+1)$

 - Can you find others?

Modeling is Critical!

- Given the human understanding of a problem, we need to answer questions like:
 - which variables shall I choose?
 - which constraints shall I enforce?
 - can I exploit any global constraints?
 - so I need any auxiliary variables?
 - are some constraints redundant, therefore can be avoided?
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 - can symmetry be eliminated?
 - are there any dual viewpoints?
 - among alternative models, which one shall I prefer?

Dual Viewpoint

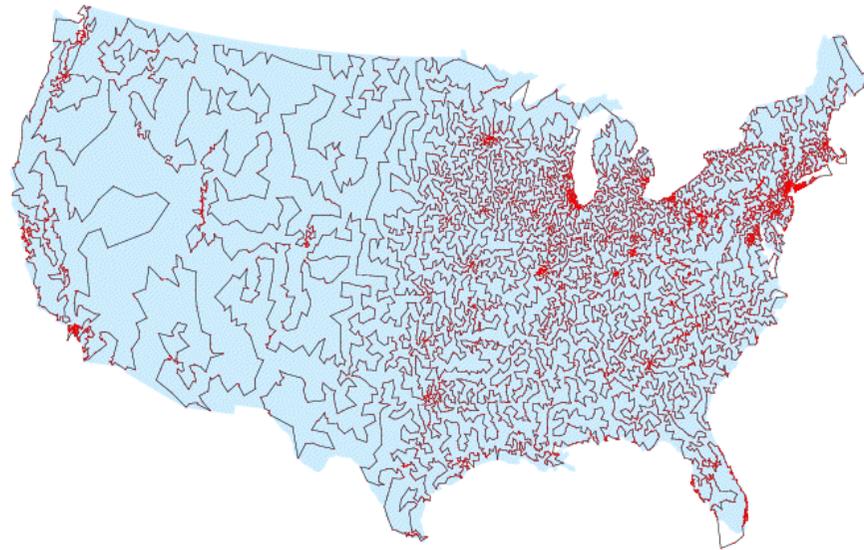
- Viewing a problem P from different perspectives may result in different models of P .
- Each models yields the same set of solutions.
- Each dual model exhibits in general a different representation of P .
 - Different variables.
 - Different domains.
 - Different constraints.
 - Different size of the search space!

Permutation Problems

- A CSP is a permutation problem if:
 - it has the same number of values as variables;
 - all variables have the same domain;
 - each variable must be assigned a different value → alldifferent constraint.
- Any solution assigns a permutation of the values to the variables.
- Other constraints determine which permutations are solutions.

Permutation Problems

- Many examples in:
 - scheduling;
 - assignment;
 - routing;
 - timetabling.



*TSP problem = find permutation of cities
which makes a tour of minimum length*

Dual Viewpoints in Permutations

- Each variable is assigned exactly one value.
- Each possible value is assigned to exactly one variable.
- The dual CSP:
 - Also a permutation problem.
 - Interchanges variables and values.
- E.g., if there are n people to do n tasks:
 - Tasks \rightarrow people
 - People \rightarrow tasks

Dual Viewpoints in Permutations

- Consider a permutation problem with n variables and values.
- **One model:**
 - Variables $[X_1, \dots, X_n]$ for n different positions.
 - Domains $\{1, \dots, n\}$ for n different values in the positions.
 - `alldifferent`($[X_1, X_2, \dots, X_n]$).
 - Other problem constraints.
- **Dual model:**
 - Variables $[Y_1, \dots, Y_n]$ for n different values in the positions.
 - Domains $\{1, \dots, n\}$ for n different positions.
 - `alldifferent`($[Y_1, Y_2, \dots, Y_n]$).
 - Other problem constraints.
- The two models may yield the same or totally different CSPs.

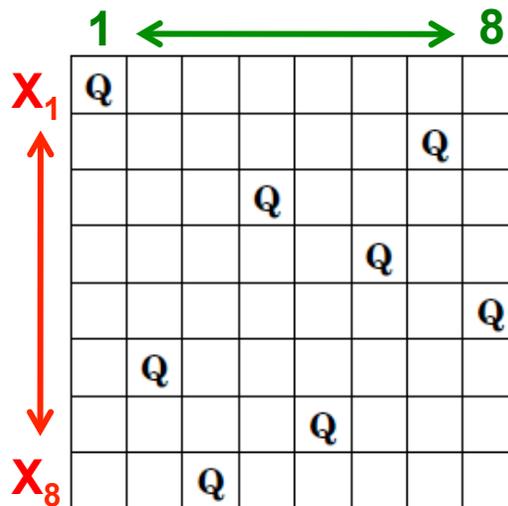
N-Queens

- Place n queens in an $n \times n$ board so that no two queens can attack each other.

| | | | | | | | |
|---|---|---|---|---|---|---|---|
| Q | | | | | | | |
| | | | | | | Q | |
| | | | Q | | | | |
| | | | | | Q | | |
| | | | | | | | Q |
| | Q | | | | | | |
| | | | | Q | | | |
| | | Q | | | | | |

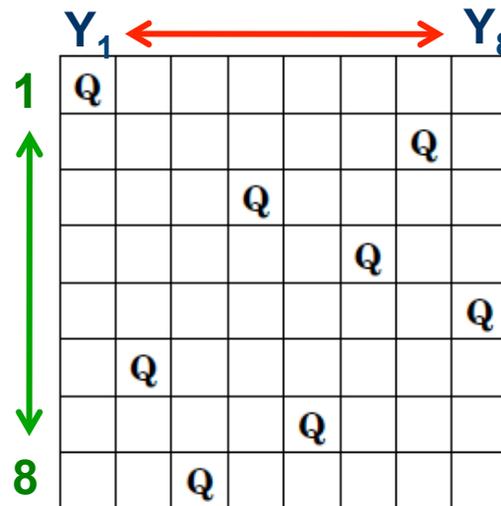
8 queens on an 8×8 chessboard, no queen attacking any other

N-Queens



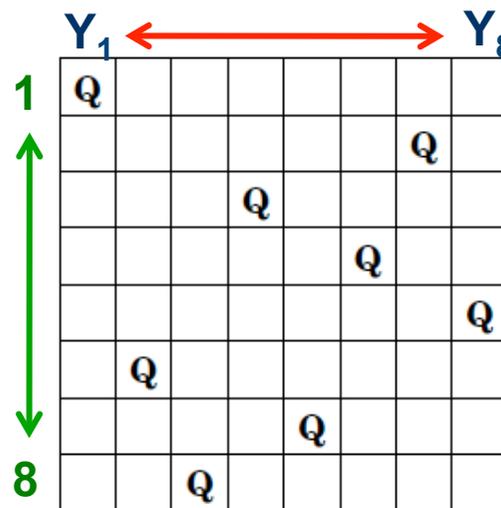
- One model
 - A **variable for each row** : $[X_1, X_2, \dots, X_n] \rightarrow$ no row attack
 - **Domain values represent the columns**: $\{1, \dots, n\}$
 - $X_i = j$ means that the queen in row i is in column j
 - Constraints:
 - $\text{alldifferent}([X_1, X_2, \dots, X_n]) \rightarrow$ no column attack
 - for all $i < j$ $|X_i - X_j| \neq |i - j| \rightarrow$ no diagonal attack

N-Queens



- Dual model
 - A **variable for each column** : $[Y_1, Y_2, \dots, Y_n] \rightarrow$ no column attack
 - **Domain values represent the rows**: $\{1, \dots, n\}$
 - $Y_i = j$ means that the queen in column i is in row j
 - Constraints:
 - $\text{alldifferent}([Y_1, Y_2, \dots, Y_n]) \rightarrow$ no row attack
 - for all $i < j$ $|Y_i - Y_j| \neq |i - j| \rightarrow$ no diagonal attack

N-Queens



Both viewpoints yield the same CSP!

- Dual model
 - A **variable for each column** : $[Y_1, Y_2, \dots, Y_n] \rightarrow$ no column attack
 - **Domain values represent the rows**: $\{1, \dots, n\}$
 - $Y_i = j$ means that the queen in column i is in row j
 - Constraints:
 - $\text{alldifferent}([Y_1, Y_2, \dots, Y_n]) \rightarrow$ no row attack
 - for all $i < j$ $|Y_i - Y_j| \neq |i - j| \rightarrow$ no diagonal attack

Langford's Problem: $L(k,n)$

- Find a sequence of length $n*k$ of numbers $[1,n]$ s.t. for all i in $[1,n]$:
 - i appears k times in the sequence;
 - there are i other numbers between the k successive occurrences of i .

$L(2,4)$



- Due to the mathematician Dudley Langford.
 - Watched his son build a tower which solved $L(2,3)$.

First Basic Model

- **Variables and Domains**
 - Find a position for each number.
 - How many numbers?
 - How many positions?

First Basic Model

- **Variables and Domains**
 - Find a position for each number.
 - How many numbers?
 - How many positions?
 - **$n \cdot k$ numbers, $n \cdot k$ positions**

First Basic Model

- **Variables and Domains**

- Find a position for each number.
- $[N_{11}, N_{21}, \dots, N_{n1}, \dots, N_{1k}, N_{2k}, \dots, N_{nk}]$
- N_{ij} for the j^{th} occurrence of i in $[1, n]$:
 - $N_{11} \rightarrow 1^{\text{st}}$ occurrence of 1
 - $N_{21} \rightarrow 1^{\text{st}}$ occurrence of 2
 - ...
 - $N_{12} \rightarrow 2^{\text{nd}}$ occurrence of 1
 - $N_{22} \rightarrow 2^{\text{nd}}$ occurrence of 2
 - ...
- Domain values $\{1, \dots, n \cdot k\}$ represent the positions in the sequence:
 - $N_{ij} = p$ means the j^{th} occurrence of i is in position p .

 i appears k times in the sequence

First Basic Model

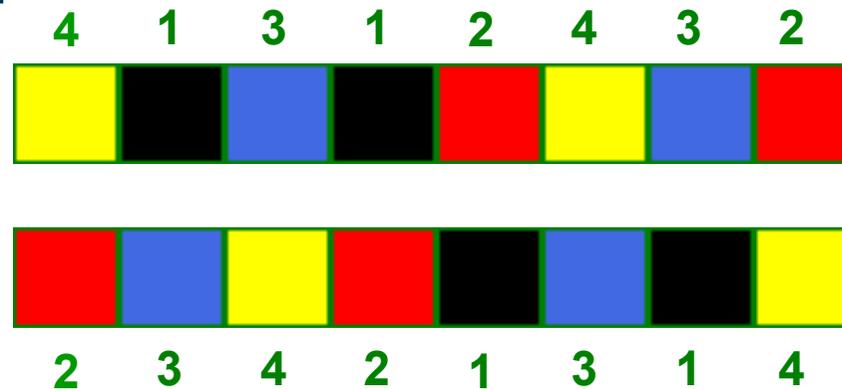
- **Constraints**

- alldifferent($[N_{11}, N_{21}, \dots, N_{n1}, \dots, N_{1k}, N_{2k}, \dots, N_{nk}]$)
- for all i, j , $N_{i(j+1)} = (i+1) + N_{ij}$

First Basic Model

- **Symmetry**

- Each sequence can be inverted.



- **Symmetry Breaking Constraints**

- $L(3,9)$:

$N_{92} < 14$ (2nd occurrence of 9 is in the 1st half)

OR $N_{92}=14$ and $N_{82} < 14$ (2nd occurrence of 8 is in the 1st half)

Dual Model

- **Variables and Domains**
 - Find a number for each position.
 - $[P_1, P_2, \dots, P_{n \times k}]$
 - Domain values?

Dual Model

- **Variables and Domains**
 - Find a number for each position.
 - $[P_1, P_2, \dots, P_{n \cdot k}]$
 - Domain values?
 - If use the numbers to place, we cannot use an alldifferent constraint.
 - Each number occurs not once but k times.

Dual Model

- **Variables and Domains**

- Find a number for each position.
- $[P_1, P_2, \dots, P_{n \cdot k}]$
- Domain values
 - Solution 1: use $\{1, \dots, n \cdot k\}$ with the value $i + (j-1) \cdot n$ standing for the j^{th} occurrence of i .

- **Constraints**

- alldifferent($[P_1, P_2, \dots, P_{n \cdot k}]$)
- distance constraints?

Dual Model

- **Variables and Domains**

- Find a number for each position.
- $[P_1, P_2, \dots, P_{n \cdot k}]$
- Domain values
 - Solution 2: use $\{1, \dots, n\}$

- **Constraints**

- Each number must occur exactly k times!
 - Global cardinality constraint
 $\text{gcc}([X_1, X_2, \dots, X_n], [v_1, \dots, v_m], [O_1, \dots, O_m])$ iff
for all $j \in \{1, \dots, m\}$ $O_j = |\{X_i \mid X_i = v_j, 1 \leq i \leq n\}|$
 - $\text{gcc}([P_1, P_2, \dots, P_{n \cdot k}], [1, \dots, n], [k, \dots, k])$

Dual Model

- **Variables and Domains**
 - Find a number for each place.
 - $[P_1, P_2, \dots, P_{n \cdot k}]$
 - Domain values
 - Solution 2: use $\{1, \dots, n\}$
- **Constraints**
 - Each number must occur exactly k times!
 - $\text{gcc}([P_1, P_2, \dots, P_{n \cdot k}], [1, \dots, n], [k, \dots, k])$
 - Distance constraints?

Which Model?

- **First Model**
 - alldifferent constraint.
 - Distance constraints.
 - Easily expressible.
 - Good propagation.
- **Dual model**
 - More natural to find numbers for positions.
 - Searching on P variables might be beneficial.
 - alldifferent/gcc constraint.
 - Distance constraints?

Combined Models

- Often different views allow different expression of the constraints and different implied constraints:
 - can be hard to decide which is better!
- We can then use multiple models and combine them via **channelling constraints** to keep consistency between the variables.
- Benefits:
 - Enhanced constraint propagation.
 - Facilitation of the expression of constraints.
 - More options for search variables.

Combined Permutation Models

- **Model 1**
 - Variables $[X_1, \dots, X_n]$ for n different positions.
 - Domains $\{1, \dots, n\}$ for n different values in the positions.
 - `alldifferent`($[X_1, X_2, \dots, X_n]$).
 - Other problem constraints PC_1 .
- **Model 2**
 - Variables $[Y_1, \dots, Y_n]$ for n different values in the positions.
 - Domains $\{1, \dots, n\}$ for n different positions.
 - `alldifferent`($[Y_1, Y_2, \dots, Y_n]$).
 - Other problem constraints PC_2 .
- **Combined Model**
 - Model 1 + Model 2
 - **Channelling constraints**

Combined Permutation Models

- **Model 1**
 - Variables $[X_1, \dots, X_n]$ for n different positions.
 - Domains $\{1, \dots, n\}$ for n different values in the positions.
 - `alldifferent`($[X_1, X_2, \dots, X_n]$).
 - Other problem constraints PC_1 .
- **Model 2**
 - Variables $[Y_1, \dots, Y_n]$ for n different values in the positions.
 - Domains $\{1, \dots, n\}$ for n different positions.
 - `alldifferent`($[Y_1, Y_2, \dots, Y_n]$).
 - Other problem constraints PC_2 .
- **Combined Model**
 - Model 1 + Model 2
 - **Channelling constraints**
 - for all i, j $X_i = j \leftrightarrow Y_j = i$

Combined Permutation Models

- Channelling Constraints
 - N-Queens
 - for all i, j $X_i = j \leftrightarrow Y_j = i$

Combined Permutation Models

- **Channelling Constraints**
 - **N-Queens**
 - for all i, j $X_i = j \leftrightarrow Y_j = i$
 - **Langford**
 - for all i, j, p $N_{ij} = p \leftrightarrow P_p = i + (j-1) * n$
 - for all i, j, p $N_{ij} = p \rightarrow P_p = i$

Combined Permutation Models

- **Benefits of Channelling**
 - **N-Queens**
 - Improved propagation.
 - Search on both sets of variables.
 - **Langford**
 - Improved propagation.
 - Search on both sets of variables.
 - Facilitates the expression of messy constraints.

Combined Permutation Models

- **Combined Model 1**
 - Model 1 + Model 2
 - **Channelling constraints**
 - for all i, j $X_i = j \leftrightarrow Y_j = i$
- **Combined Model 2**
 - Model 1 + $[Y_1, Y_2, \dots, Y_n]$, or
 - $[X_1, X_2, \dots, X_n]$ + Model 2
 - **Channelling constraints**
 - for all i, j $X_i = j \leftrightarrow Y_j = i$

Combined Permutation Models

- **Observation**
 - alldifferent is **redundant**.
- **Combined Model 3**
 - $[X_1, X_2, \dots, X_n] + PC_1 + [Y_1, Y_2, \dots, Y_n] + PC_2$
 - **Channelling constraints**
 - for all i, j $X_i = j \leftrightarrow Y_j = i$
- **Combined Model 4**
 - $[X_1, X_2, \dots, X_n] + [Y_1, Y_2, \dots, Y_n] + PC$ on suitable variables
 - **Channelling constraints**
 - for all i, j $X_i = j \leftrightarrow Y_j = i$

Combined Permutation Models

- **Best N-Queens Problem Model**

- **Variables**

- $[X_1, X_2, \dots, X_n], [Y_1, Y_2, \dots, Y_n]$

- **Constraints**

- for all i, j $X_i = j \leftrightarrow Y_j = i \rightarrow$ Channelling constraints
- for all $i < j$ $|X_i - X_j| \neq |i - j| \rightarrow PC_1$

- **Search variables**

- $[X_1, X_2, \dots, X_n, Y_1, Y_2, \dots, Y_n]$

Combined Permutation Models

- **Best Langford's Problem Model**

- **Variables**

- $[N_{11}, N_{21}, \dots, N_{n1}, \dots, N_{1k}, N_{2k}, \dots, N_{nk}], [P_1, P_2, \dots, P_{n \cdot k}]$

- **Constraints**

- for all i, j, p $N_{ij} = p \leftrightarrow P_p = i + (j-1) \cdot n \rightarrow$ **Channelling constraints**
- for all i, j , $N_{i(j+1)} = (i+1) + N_{ij} \rightarrow$ **PC₁**

- **Search variables**

- $[P_1, P_2, \dots, P_{n \cdot k}]$
- $[N_{11}, N_{21}, \dots, N_{n1}, \dots, N_{1k}, N_{2k}, \dots, N_{nk}, P_1, P_2, \dots, P_{n \cdot k}]$

Combined Models in General

- **Best Golomb Ruler Problem Model**

- **Variables:**

- $[X_1, X_2, \dots, X_m], [D_{12}, D_{13}, \dots, D_{(m-1)m}]$

- **Constraints:**

- for all $i < j$ $D_{ij} = X_j - X_i$
- $\text{alldifferent}([D_{12}, D_{13}, \dots, D_{(m-1)m}])$
- $X_1 < X_2 < \dots < X_m$
- $X_1 = 0$
- for all $i < j < k$ $D_{ij} < D_{ik}$ and $D_{jk} < D_{ik}$
 $D_{ik} = D_{ij} + D_{jk}$
- for all $i < j$ $D_{ij} \geq \text{Ruler}(j-i+1)$

→ Channelling constraints

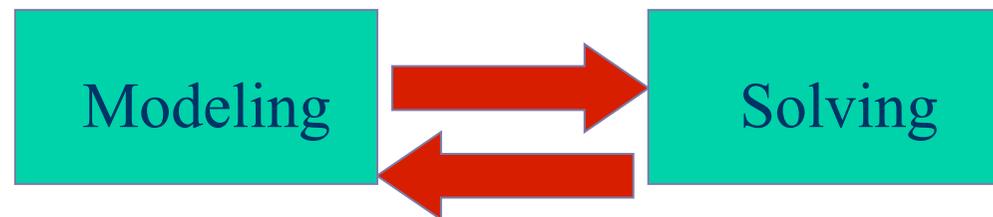
Problem/IMPLIED/symmetry
breaking constraints
on suitable variables

- **Search Variables:**

- $[X_1, X_2, \dots, X_m]$

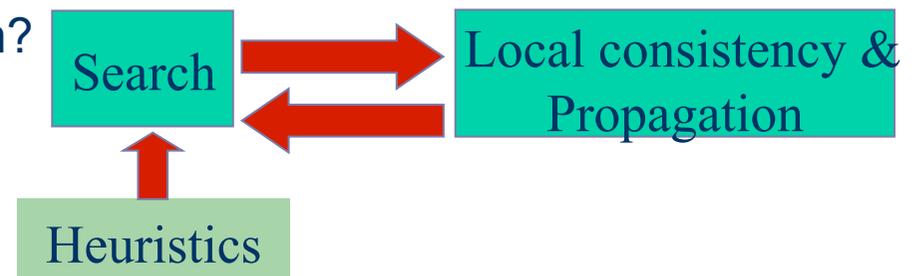
CP Machinery

- Modeling and solving are strongly interconnected.



Solving

- The user lets the CP technology solve the CSP:
 - choose a search algorithm:
 - usually backtracking search performing a depth-first traversal of a search tree.
 - integrate local consistency and propagation.
 - choose heuristics for branching on the search tree:
 - which variable to branch on?
 - which value to branch on?



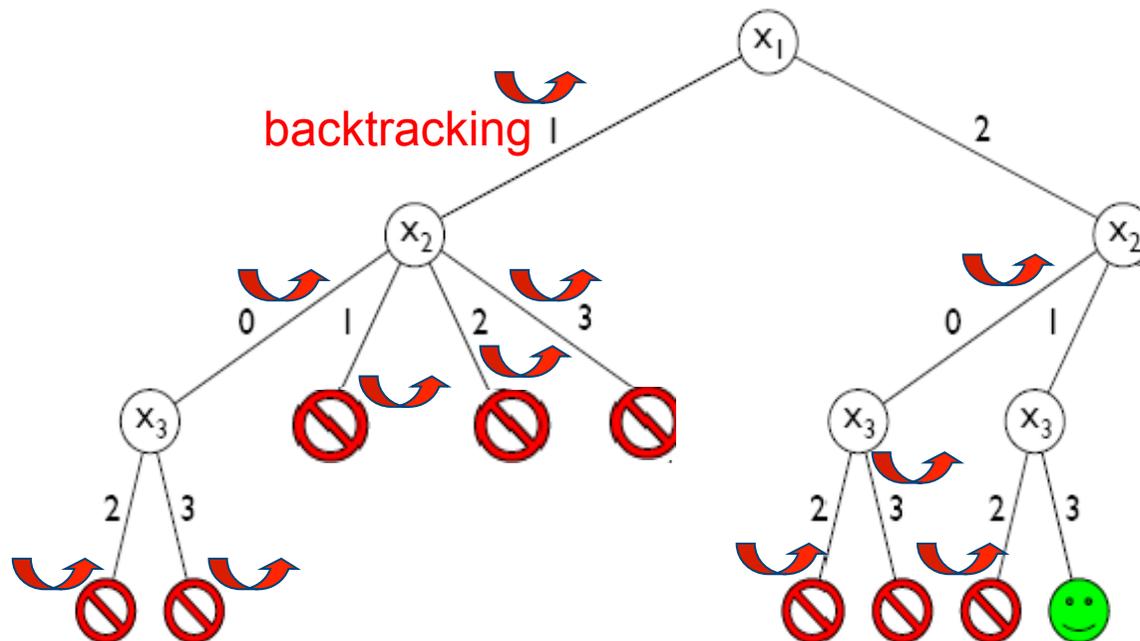
Backtracking Search

- A possible efficient and simple method.
- Variables are instantiated sequentially.
- Whenever all the variables of a constraint is instantiated, the validity of the constraint is checked.
- If a (partial) instantiation violates a constraint, backtracking is performed to the most recently instantiated variable that still has alternative values.
- Backtracking eliminates a subspace from the cartesian product of all variable domains.
- Essentially performs a depth-first search.

Backtracking Search

- $X_1 \in \{1,2\}$ $X_2 \in \{0,1,2,3\}$ $X_3 \in \{2,3\}$
- $X_1 > X_2$ and $X_1 + X_2 = X_3$ and $\text{alldifferent}([X_1, X_2, X_3])$

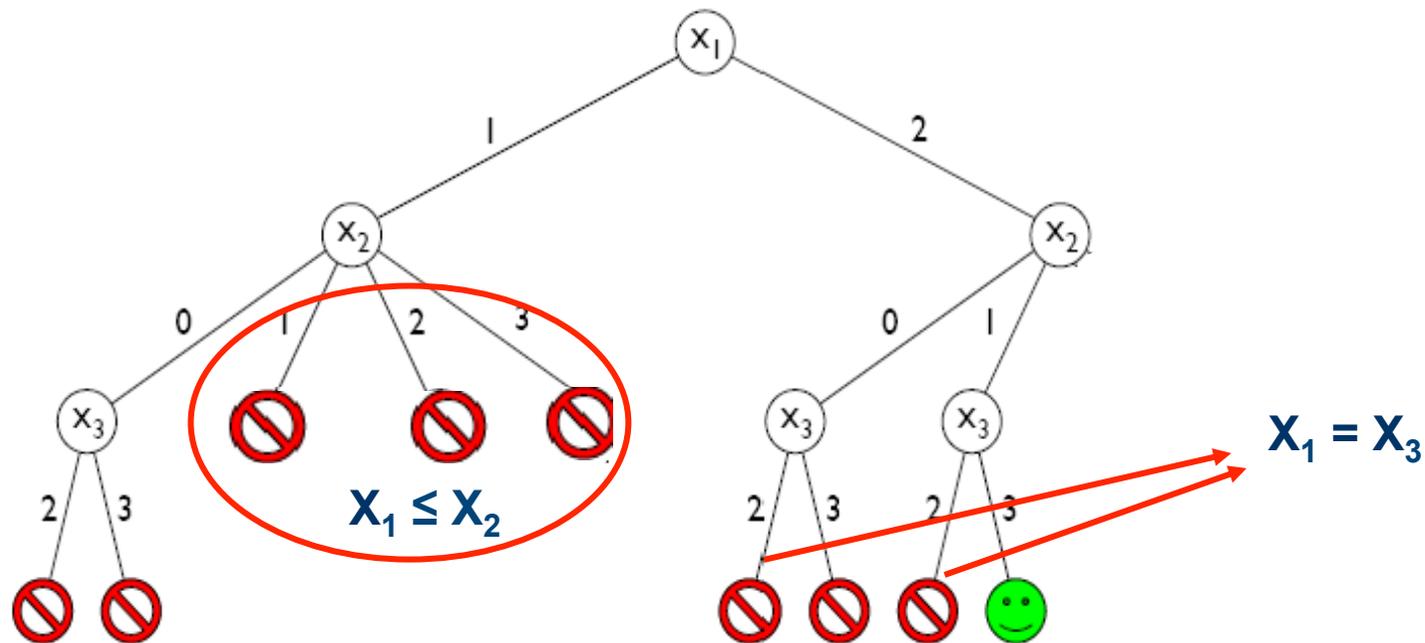
Backtracking search



Fails 8 times!

Backtracking Search

- Backtracking suffers from thrashing ☹️ :
 - performs checks only with the current and past variables;
 - search keeps failing for the same reasons.

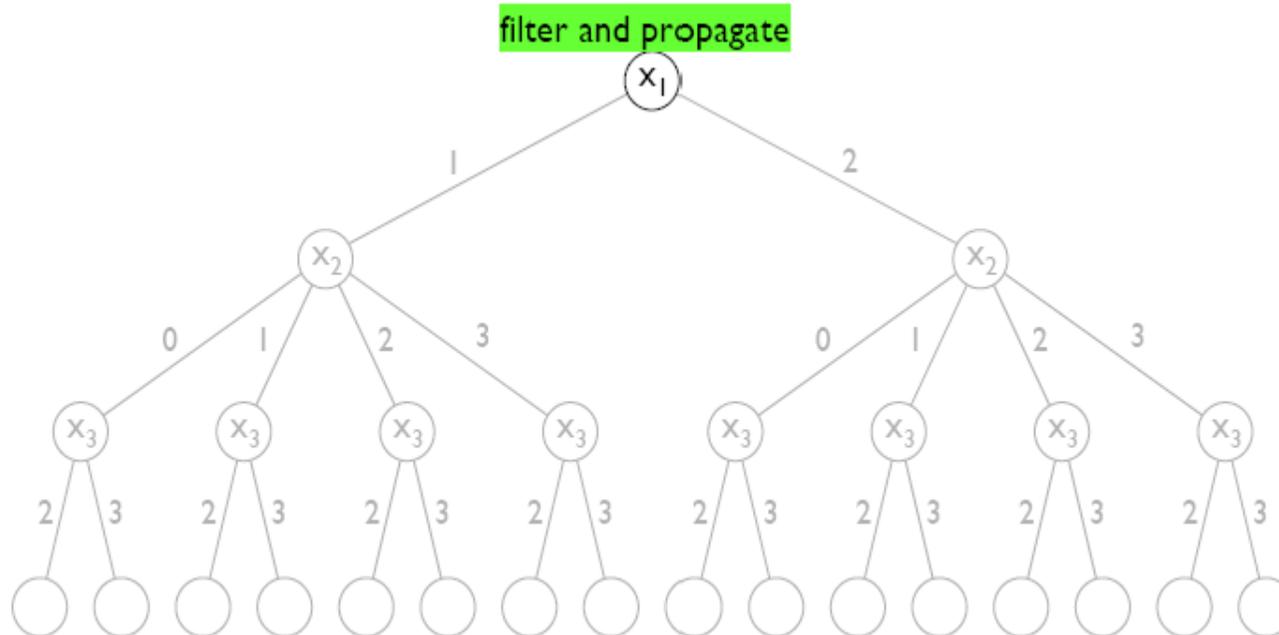


Local Consistency & Propagation

- We can reason about the properties of constraints and their effect on their variables.
- Some values can be **filtered** from some domains, reducing the backtracking search space significantly!

BS + Local Consistency & Propagation

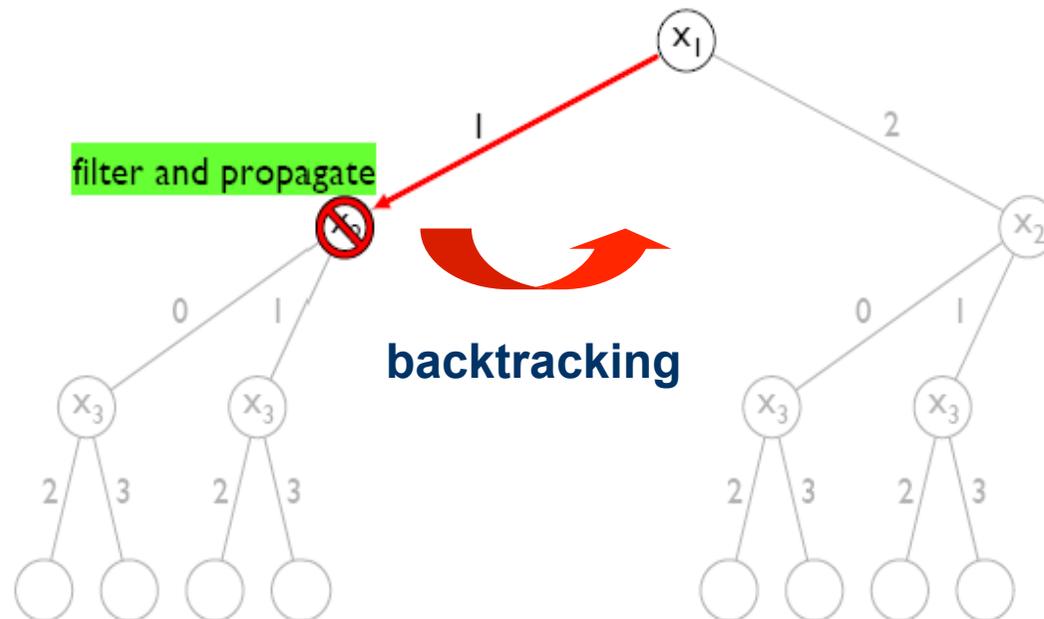
- $X_1 \in \{1,2\}$ $X_2 \in \{0,1,2,3\}$ $X_3 \in \{2,3\}$
- $X_1 > X_2$ and $X_1 + X_2 = X_3$ and alldifferent($[X_1, X_2, X_3]$)



BS + Local Consistency & Propagation

- $X_1 \in \{1, \cancel{2}\}$ $X_2 \in \{0, \cancel{1}\}$ $X_3 \in \{\cancel{2}, \cancel{3}\}$
- $X_1 > X_2$ and $X_1 + X_2 = X_3$ and $\text{alldifferent}([X_1, X_2, X_3])$

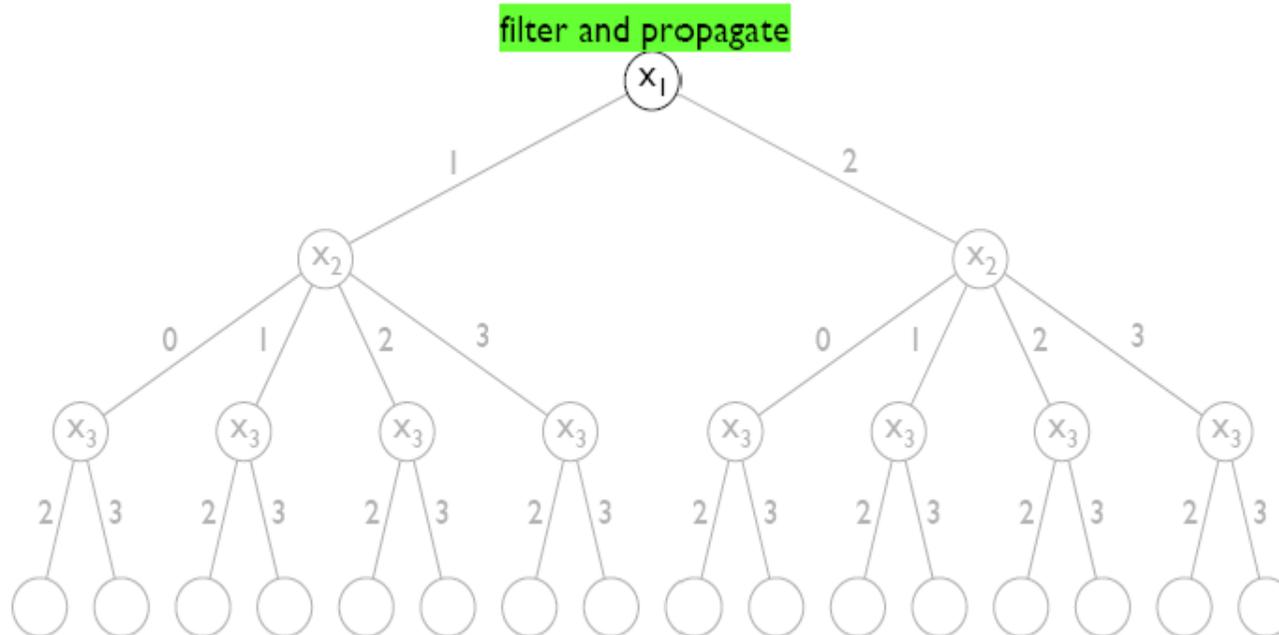
Backtracking search + local consistency/propagation



BS + Local Consistency & Propagation

- $X_1 \in \{1,2\}$ $X_2 \in \{0,1,2,3\}$ $X_3 \in \{2,3\}$
- $X_1 > X_2$ and $X_1 + X_2 = X_3$ and $\text{alldifferent}([X_1, X_2, X_3])$

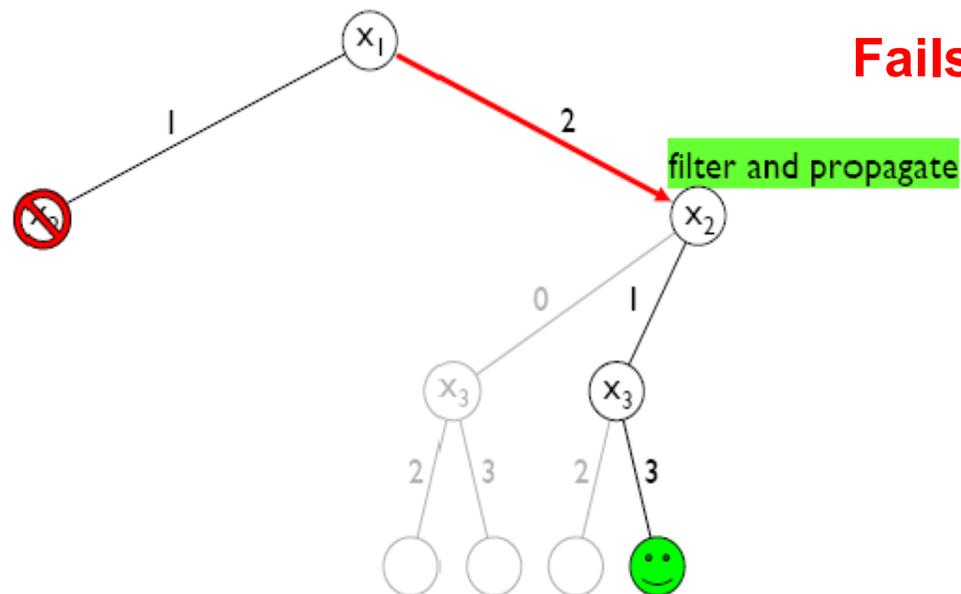
Backtracking search + local consistency/propagation



BS + Local Consistency & Propagation

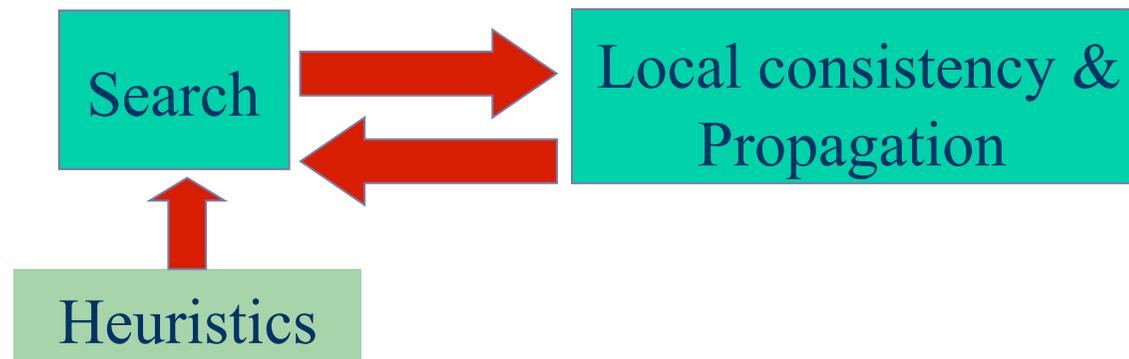
- $X_1 \in \{1,2\}$ $X_2 \in \{0,1\}$ $X_3 \in \{2,3\}$
- $X_1 > X_2$ and $X_1 + X_2 = X_3$ and alldifferent($[X_1, X_2, X_3]$)

Backtracking search + local consistency/propagation



Local Consistency & Propagation

- Central to the process of solving CSPs which are inherently intractable.



PART II: Local Consistency & Constraint Propagation



AIMMS

- Modeling language with an interface to CP and MIP solvers (<http://www.aimms.com/cp>)
- Student license
- GUI support available only in the Windows version
 - Create a virtual machine with Windows OS in your computer
 - Virtual Box (<https://www.virtualbox.org/>)
- Extensive documentation
 - One-hour tutorial (general introduction)
 - Chapter 21 – “Constraint Programming” of Language Reference