

# Overlay Networks Topology Management

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- Peer-to-Peer (P2P) overlay networks
- Epidemic network topology
- Newscast algorithm
- Aggregation Protocol





#### Peer-to-Peer vs. Client-Server

- No distinction between Client and Server nodes
  - All nodes are peers (they both ask for and require services)
- High decentralization
  - No global knowledge or control
  - Each node has limited knowledge of the system and is able to interact with a (small) set of other nodes





### P2P Overlay Networks

Peer-to-Peer (P2P) networks are usually overlay

- Physical and communication level are assumed to exist
  - The Internet and TCP/IP protocol
  - Possibility to communicate between any pair of nodes
- Nodes: processes running at given locations
- Links: defined upon underlying physical topology
  - Logical label (i.e. IP,port pair)





### **Views and Topology**

- The set of nodes that a peer knows about is called its view
- Typically, views are a (very) small subset of all nodes
- Views define the overlay network topology on top of the physical network
  - Overlay network nodes and links between them
- Overlay network topology easy to manage
  - No physical modification needed
  - Views exchange





### **Topology Management**

- How to ensure that the overlay network topology satisfies certain properties:
  - has a desired structure (connected, random graph, ring, torus, binary tree, etc.)
  - maintains the desired structure in a dynamic setting (churn)
- Problem to be solved is topology management
- Solution inspired from the spreading of epidemics





- Some people (nodes) in the population (network) is infected (has some information)
- Each person knows a set of other people (view)
- Each infected person transmits the infection to a subset of the people he knows
- Infection (information) propagates throughout the population





#### Each node (say *i*) periodically:

- Chooses a neighbor (say j) through selectPeer() function
- Nodes i and j exchange their views
- Views are merged/updated using updateState() function
- Through the definition of selectPeer() and updateState() functions it is possible to define different emerging global topology

























Both peers apply updateState() thereby redefining topology





### **Using Epidemic Scheme**

- Epidemic spreading model define the way information is exchanged
- Some basic steps are not specified
  - SelectPeer() method
  - UpdateState() method
- Many different target topology can be achieved by defining these two methods





### Random Topology: Newscast

#### Node views

- Limited size (maximum view size=c)
- Node descriptors: timestamped links
- SelectPeer()
  - A random neighbor is chosen
- Each node adds a fresh descriptor of itself in its view
- Views are exchanged
- UpdateState()
  - Views are merged
  - Duplicated entries removed
  - To respect view size limit oldest descriptors are dropped



















Each node add its fresh descriptor to its view (suppose t=15)







Views are merged







Oldest descriptors discharged (suppose c=6)













- Continuously changing
- Random-like
  - Low clustering coefficient
  - Low average path length
    - Logarithmic growth with respect to network size
- Robust to node removal
  - Dinamicity
  - Aging of older and not refreshed descriptors



### Newscast Resulting Topology

Clustering coefficient (CC)

 Proportion between links between neighbors of a node and the number of possible links

Network CC

BISON

Average of nodes' CC





- Average Path Length (APL)
  - Average hop-distance between any pair of nodes



#### Newscast Robustness







## Aggregation





## **Collective Computation (Aggregation)**

- Each node has a (numeric) local state
- Compute (global) aggregate function over the initial values at *all* nodes
- The aggregate value has to be known by each node
- Examples of aggregate functions:
  - Average
  - Min-max
  - Geometric mean
  - Variance
  - Network size





- Each node periodically selects another (random) peer and exchanges local state information
  - Random peer sampling provided by Newscast
- Each node updates its local state based on the information exchanged
- System fully symmetric all nodes act identically
- Communication is symmetric "push-pull" epidemic





- Local value S<sub>p</sub> contains current estimate of the aggregate
- Suppose the (random) peer picked by node *p* is *q*
- Nodes p and q exchange current estimates
- Update local estimates depending on aggregation function
  - Average:  $S_p \leftarrow \frac{(S_p + S_q)}{2}$
  - Geometric mean:  $S_p \leftarrow \sqrt{(S_p S_q)}$
  - Maximum:  $S_p \leftarrow max(S_p, S_q)$
- Other, more complex functions can be built combining elementary functions





### **Properties of Epidemic-Based Aggregation**

- If peers are sampled randomly sampled, then the variance of nodes' estimates decreases exponentially
  - Aggregation goal: variance=0
    - Same estimated value throughout the network
- Global random peer sampling provided by Newscast!





- Through averaging it is possible to evaluate network size
- Aggregation initialized with one node at value 1 and every other nodes at value 0
- Average aggregation performed
- Each nodes reaches value 1/N
  - Network size is simply 1/(1/N)































1/0.166 = 6.02 ≈ N





#### **Exponential Convergence of Average**







#### **Robustness of Average**





Bibliography

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